



Advanced Unit Testing in Hedron

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What's Hedron?

Microhypervisor with capability-based security model written in C++

Focus on:





Why (Unit) Test?

A healthy software project

- is easy to change by multiple people
- with confidence that it doesn't break.

Good unit test coverage helps:

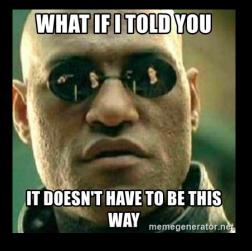
- Tests can run anywhere,
- developer feedback in seconds,
- sanitizers (UBSAN, ASAN, ...)!

Kernels Are Not Doing Well

OS kernels are particularly hostile:

- strange programming environments
- interaction with hardware
- mindset / lack of education

Result: Usually extremely poor (unit) tests.



Let's spread some testing ideas!

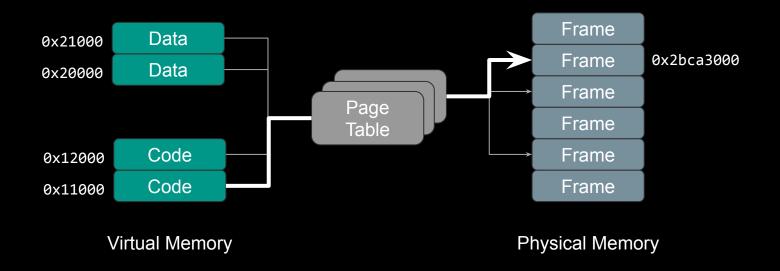
Example: Page Table Manipulation

Needed to modify Hedron's page table. No existing tests!

- Important piece of code in a microkernel.
- Bugs are extremely hard to debug.

How to get good test coverage? We decided to redesign it.

Recap: Address Spaces



Recap: Page Tables (x86-64)



Unit Testing Challenges

Code uses kernel-internal APIs.

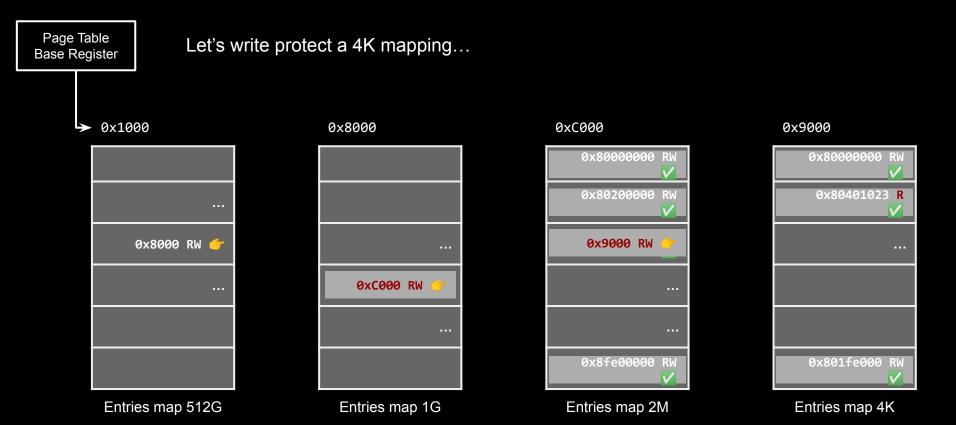
Prevents us from compiling the code for Linux.

2. CPU reads page tables while they are being modified.

Checking result when everything is done: Insufficient!



Updating Page Tables (x86-64)



Idea: Record Observable Side Effects

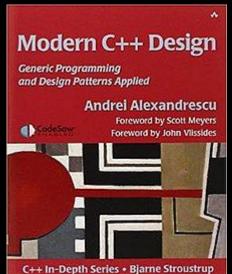
In unit tests, we want to:

- Record all visible side-effects while the code is running.
- Check all transient states for validity.



Applying Policy-Based Design

```
template <int BITS_PER_LEVEL, typename ENTRY, typename MEMORY, typename CACHE_FLUSH,
         typename PAGE_ALLOC, typename DEFERRED_CL
                                                       IP, typename ATTR>
class Generic_page_table
   private:
       MEMORY
                    memory:
       // Use a superpage from the given level to fill out a new page table one
       // hierarchy deeper with the same mappings.
       void fill from superpage(pte_pointer_t new_table, pte_t superpage_pte, level_t cur_level)
            assert_slow (is_superpage(cur_level, superpage_pte));
           pte_t attr_mask {cur_level == 1 ? static_cast<pte_t>(ATTR::PTE_S) : 0};
            for (size_t i {0}; i < static_cast<size_t>(1) << BITS_PER_LEVEL; i++) {</pre>
                pte_t offset
                                   (PAGE_BITS + (cur_level - 1) * BITS_PER_LEVEL)};
               memory_.write (new_table + i, (superpage_pte & ~attr_mask) | offset);
            flush_cache_page (new_table);
```



```
static entry read (pointer ptr) { return Atomic::load(*ptr); }
static void write (pointer ptr, entry e) { Atomic::store(*ptr, e); }
```

class Atomic_access_policy

public:

```
class Fake_memory
       using location = std::pair<pointer, entry>;
       using memory_list = std::forward_list<location>;
       // Keep a list of address/value pairs that can be prepended with new
       // content.
       memory_list memory_;
    public:
       void write(pointer ptr, entry e)
            memory_.emplace_front (ptr, e);
       entry read(pointer ptr) const
            auto it {std::find_if (memory_.cbegin(), memory_.cend(),
                                  [ptr] (auto const &pair) { return pair.first == ptr; } )};
           if (it == memory_.cend()) {
               // Reading unwritten memory. Fake_page_alloc relies on this being
                // zero.
                return 0;
            return it->second;
```

```
TEST CASE("Superpage splitting is correct") {
    Fake_hpt hpt {4, 3};
   // What we initially have in the page table.
    Fake_hpt::Mapping const initial_1gb_mapping {0, 0x80000000, Fake_attr::PTE_P, onegb_order};
   // Allowed transient mappings.
   Fake_hpt::Mapping const trans_2mb_mapping {0, 0x80000000, Fake_attr::PTE_P, twomb_order};
    Fake_hpt::Mapping const trans_4k_mapping {0, 0x80000000, Fake_attr::PTE P, PAGE BITS};
   // Final mapping.
   Fake hpt::Mapping const page mapping {0, 0xC0000000, Fake attr::PTE P | Fake attr::PTE W, PAGE BITS};
    // Create initial state.
   hpt.update (initial_1gb_mapping);
    auto before mapping {hpt.memory().now()};
   // Map 4K over 1GB page.
   hpt.update(page_mapping);
   // Look through all intermediate states of the memory.
    for (auto it {hpt.memory().now()}; it != before mapping; ++it) {
        auto const cur map {rewind (it, hpt).lookup (0)};
        REQUIRE((cur_map == initial_1gb_mapping ||
                 cur_map == trans_2mb_mapping ||
                 cur_map == trans_4k_mapping ||
                cur_map == page_mapping));
```

Taking It Further

We can use this technique to test other properties:

- Are page tables disconnected before they are deallocated?
- Do we handle atomic-compare-exchange failures?
- Do we read/write memory exactly as often as needed?
- ...

We only started to tap the potential.

C++ 20 Concepts: Better Error Messages

Incompatible policy classes lead to hard to read C++ error messages...

```
template <typename T, typename P, typename V>
concept MemoryAccess = requires(T v, P p, V v)
{
     {v.read(p)} -> std::convertible_to<uint64_t>;
     {v.write(p, v)};

     // ...
};
```

Somewhat similar to traits in Rust.

Summary

We re-wrote Hedron's page table code:

- after deciding what we want to test,
- using policy-based design,
- to unit test otherwise hard-to-test properties,
- by recording all observable side effects of operations,
- with a concept that's applicable to C++ and Rust.

Let's improve kernel testing: Please share your test stories!

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