

Yandex



Caveats of replication in HA clusters and CDC systems

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Yandex and PostgreSQL

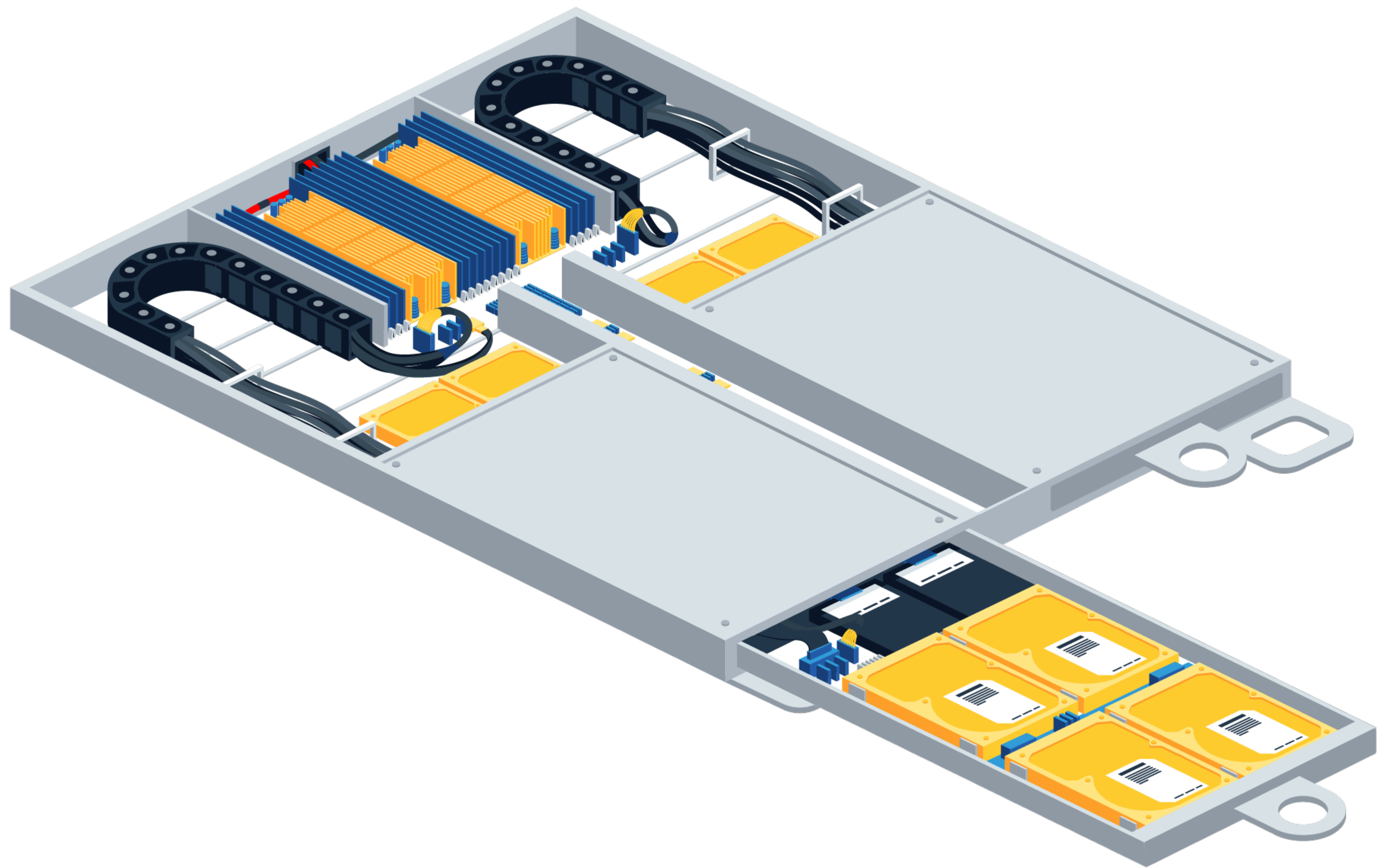
Yandex.Cloud

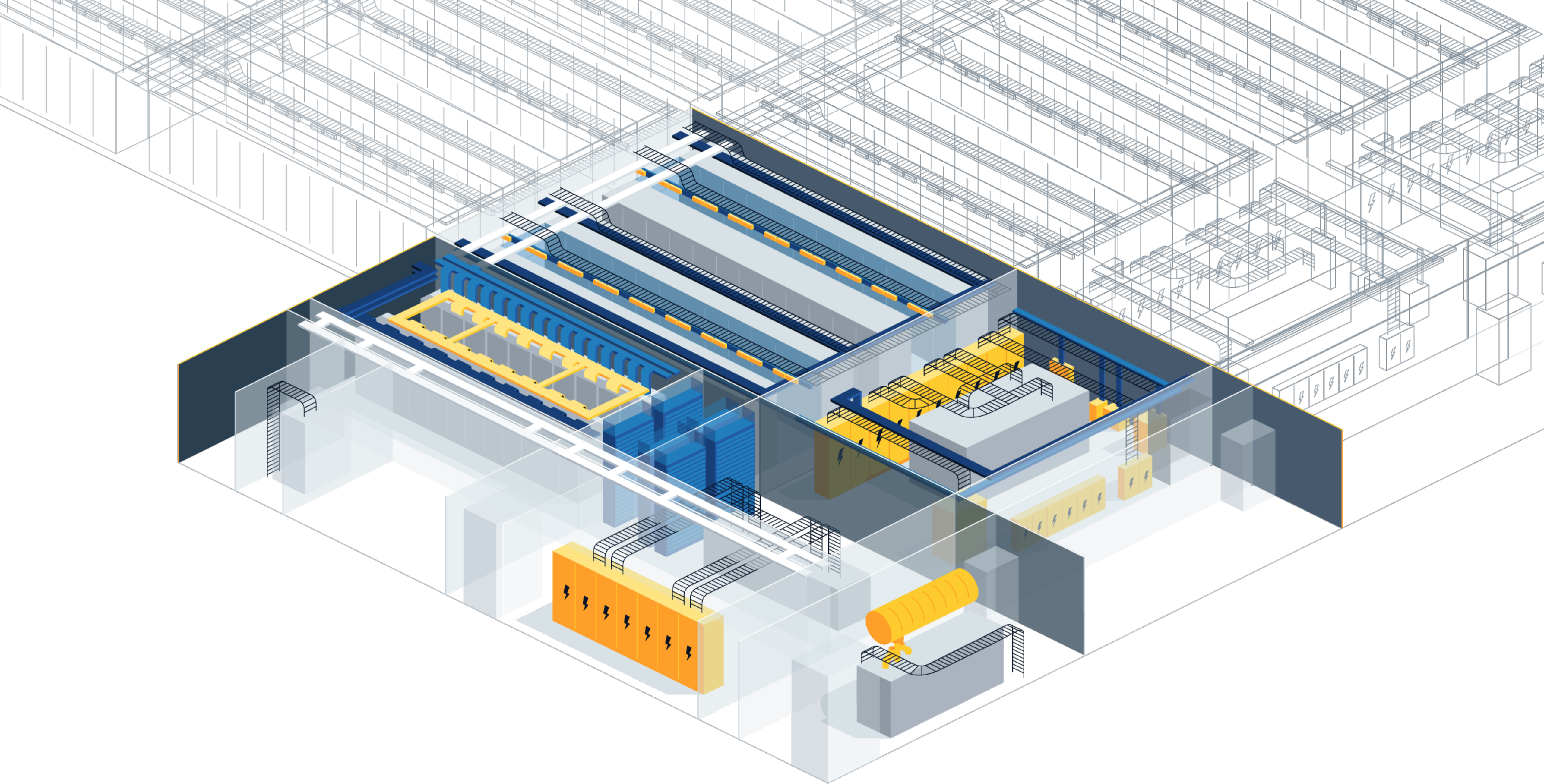
- › 2+ petabytes of Postgres
- › ~3+ million requests per second

And many other services like Yandex.Mail, Yandex.Taxi, Yandex.Maps, weather forecast, carsharing, food delivery etc.

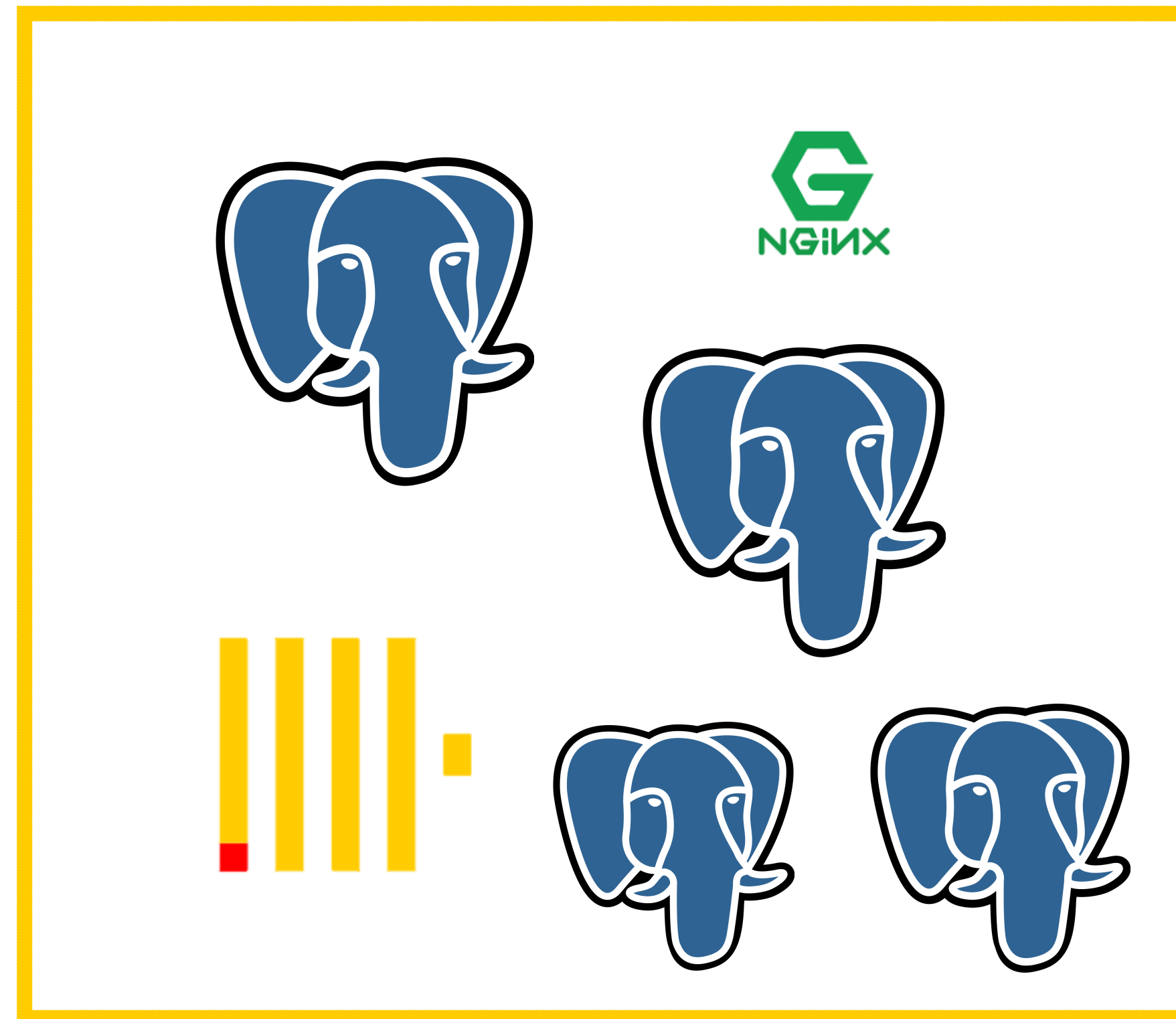
Essential expectations

- › 0.9999 read availability
- › 0.9995 write availability
- › Scalable multi-AZ deployments
- › Most up-to date copy of operational data in analytical system

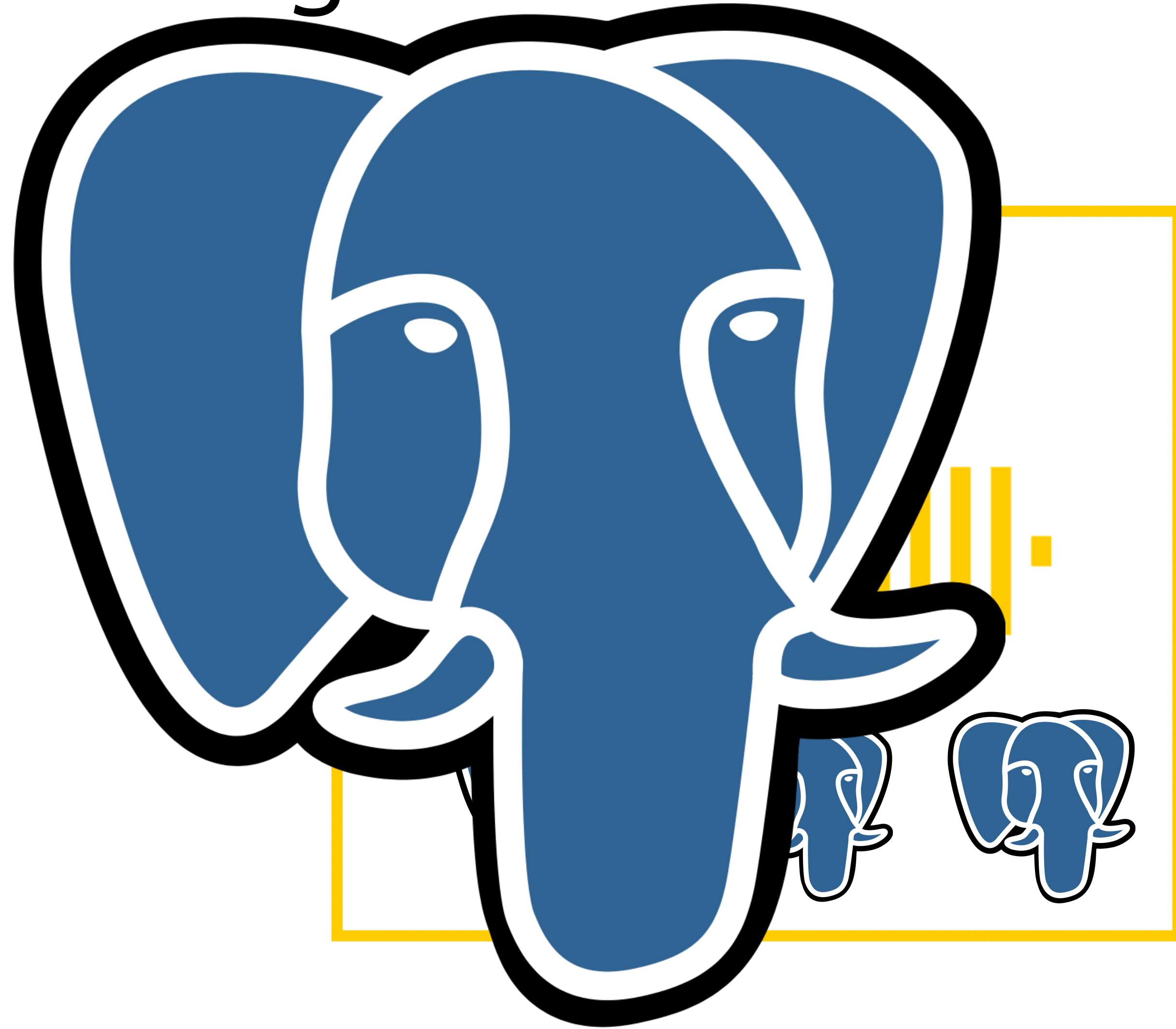




Virtualization utilize resources efficiently



Vertical scaling is kind of failure too



| Redundancy

Network block storage

NBS is a kind of redundancy



But databases are working better on local drives
This will not be discussed in this talk

Incrementally rebuilding copy of the DB

- › WAL archive
- › Streaming replication

WAL archive

- Primary node calls `archive_command` for every segment (typically 16Mb)
 - › Synchronous interface, but `wal-g` and `pg_backrest` try to solve it
- Standby nodes recover segments calling `restore_command`

WAL archive

- Primary node calls `archive_command` for every segment (typically 16Mb)

- › Synchronous interface, but `wal-g` and `pg_backrest` try to solve it

- Standby nodes recover segments calling `restore_command`

- › It's preferred to have archive even if you use streaming replication

Streaming replication

■ walsender\walreceiver processes work in pair sending WAL with granularity up to one WAL record

- › Standby startup process can shuttle between archive and replication

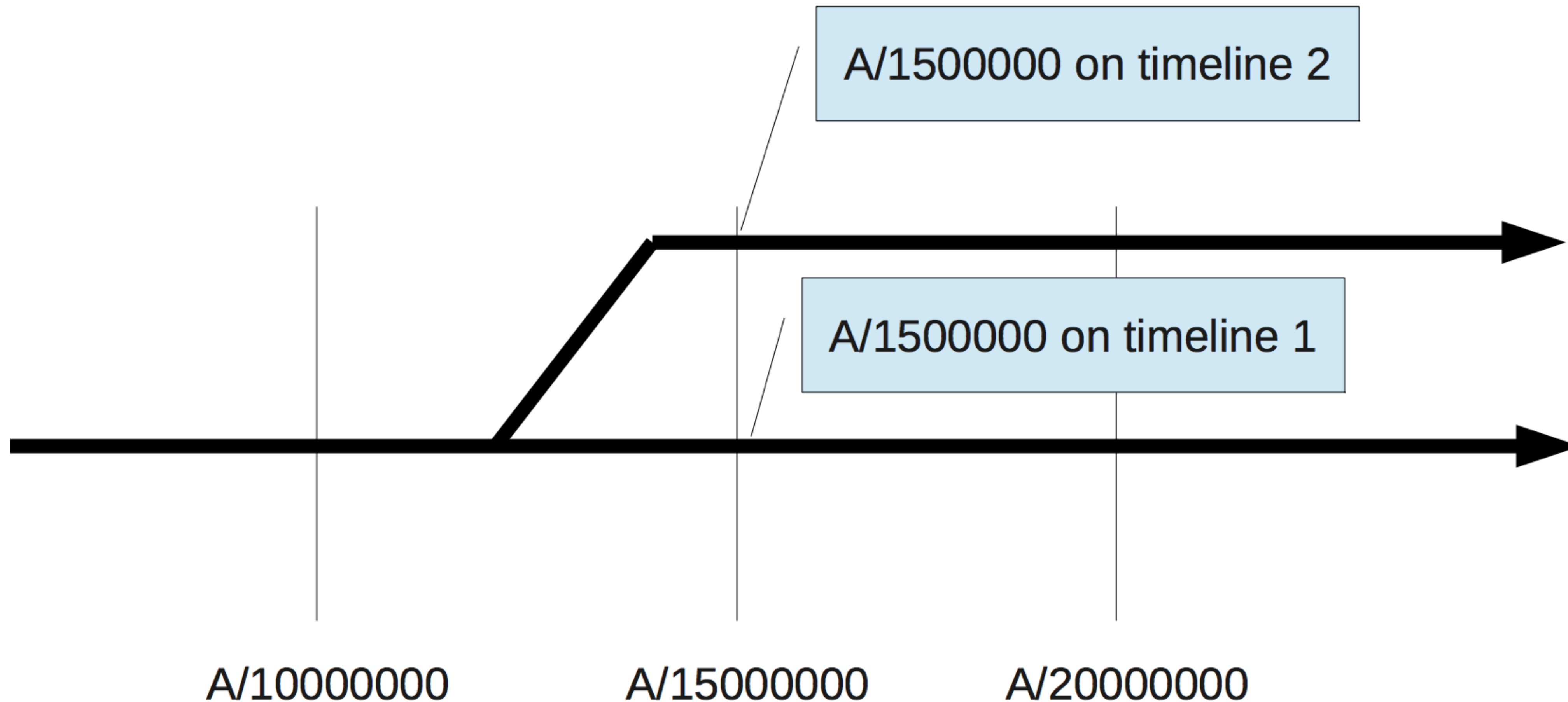
■ **Cascade is possible**

- › Replication slots have information what was sended and appied by replication target

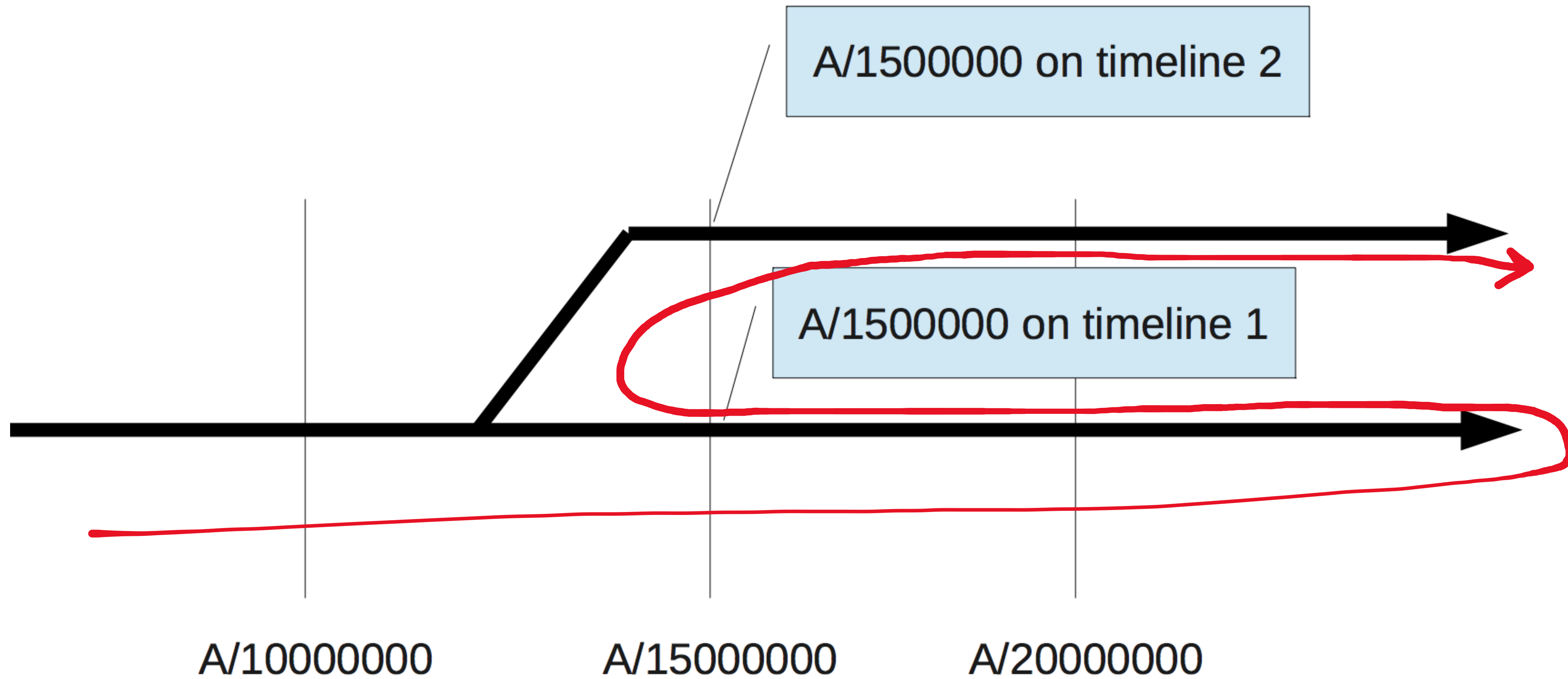
Synchronous streaming replication

- › Do not acknowledge commit to client until replica has all data wrt current transaction

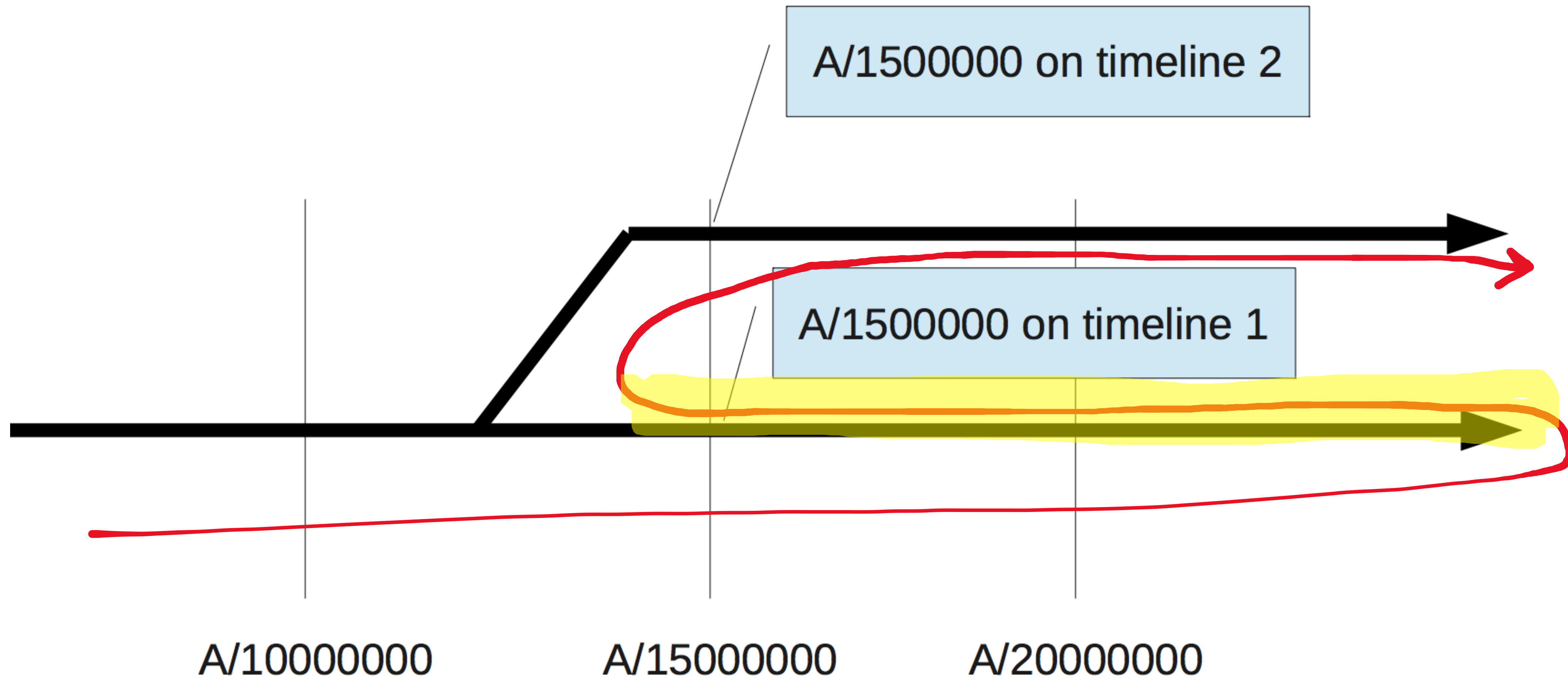
PostgreSQL Timeline



pg_rewind



pg_rewind

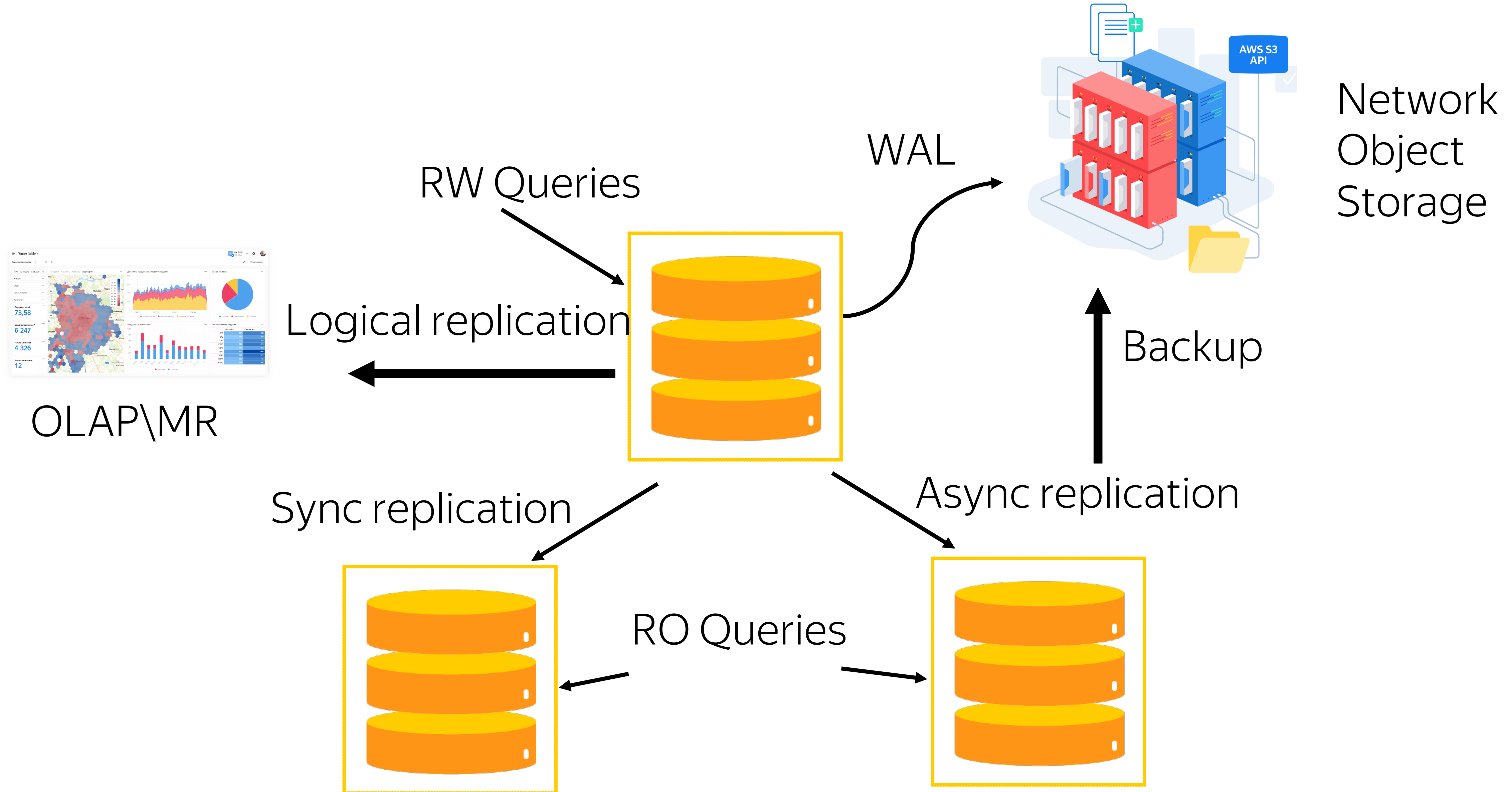


Logical replication

Good for replicating:

- › only parts of a database
- › to other systems (e.g. OLAP)
- › Between different versions

Cluster in the cloud



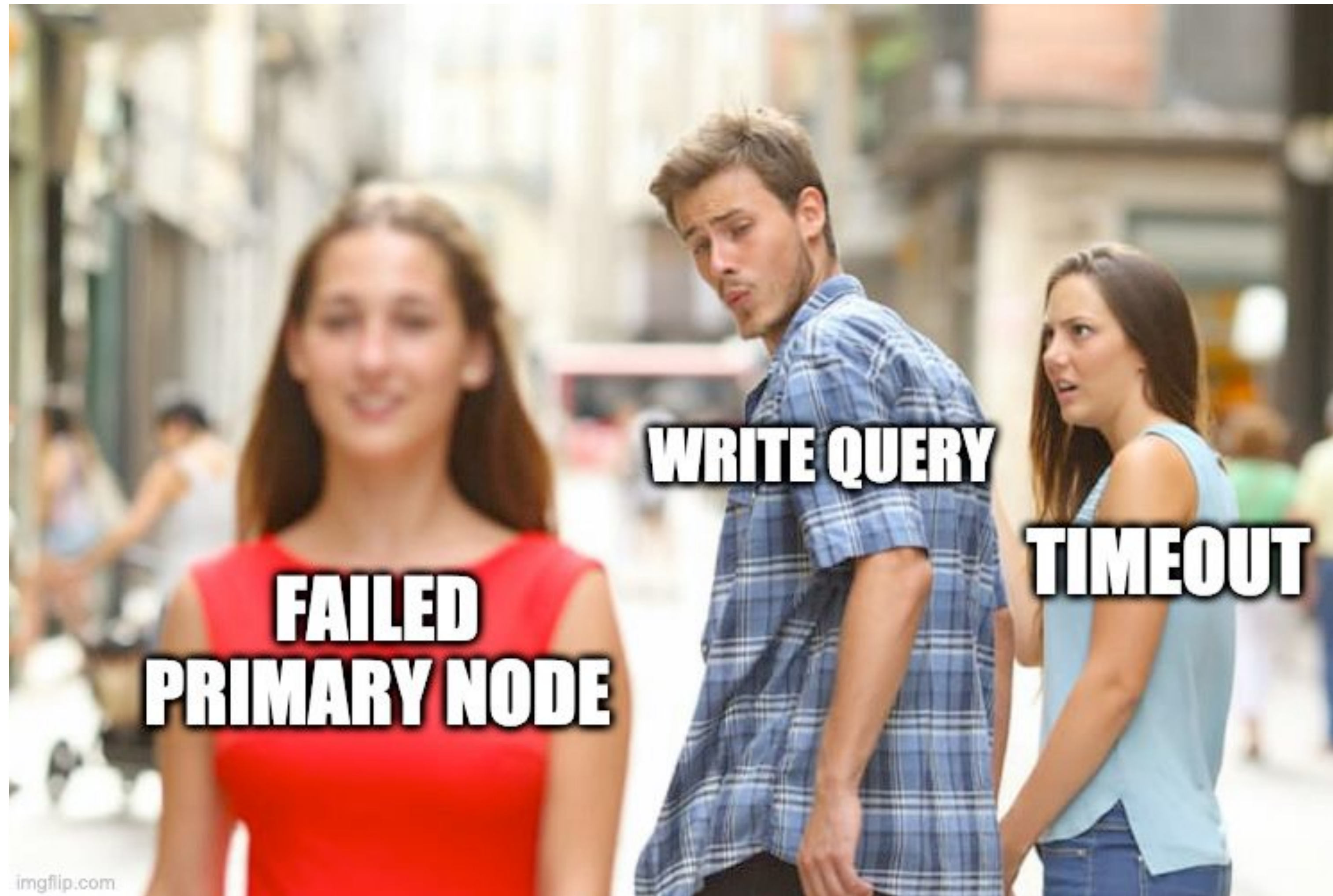
How to find primary node?

```
psql "host=<host 1 FQDN>,<host 2 FQDN>,<host 3 FQDN> \  
port=6432 \  
sslmode=verify-full \  
dbname=<DB name> \  
user= \  
target_session_attrs=read-write"
```

How to find primary node?

```
psql "host=<host 1 FQDN>,<host 2 FQDN>,<host 3 FQDN> \  
port=6432 \  
sslmode=verify-full \  
dbname=<DB name> \  
user= \  
target_session_attrs=read-write"
```

But how do you know the node had failed?



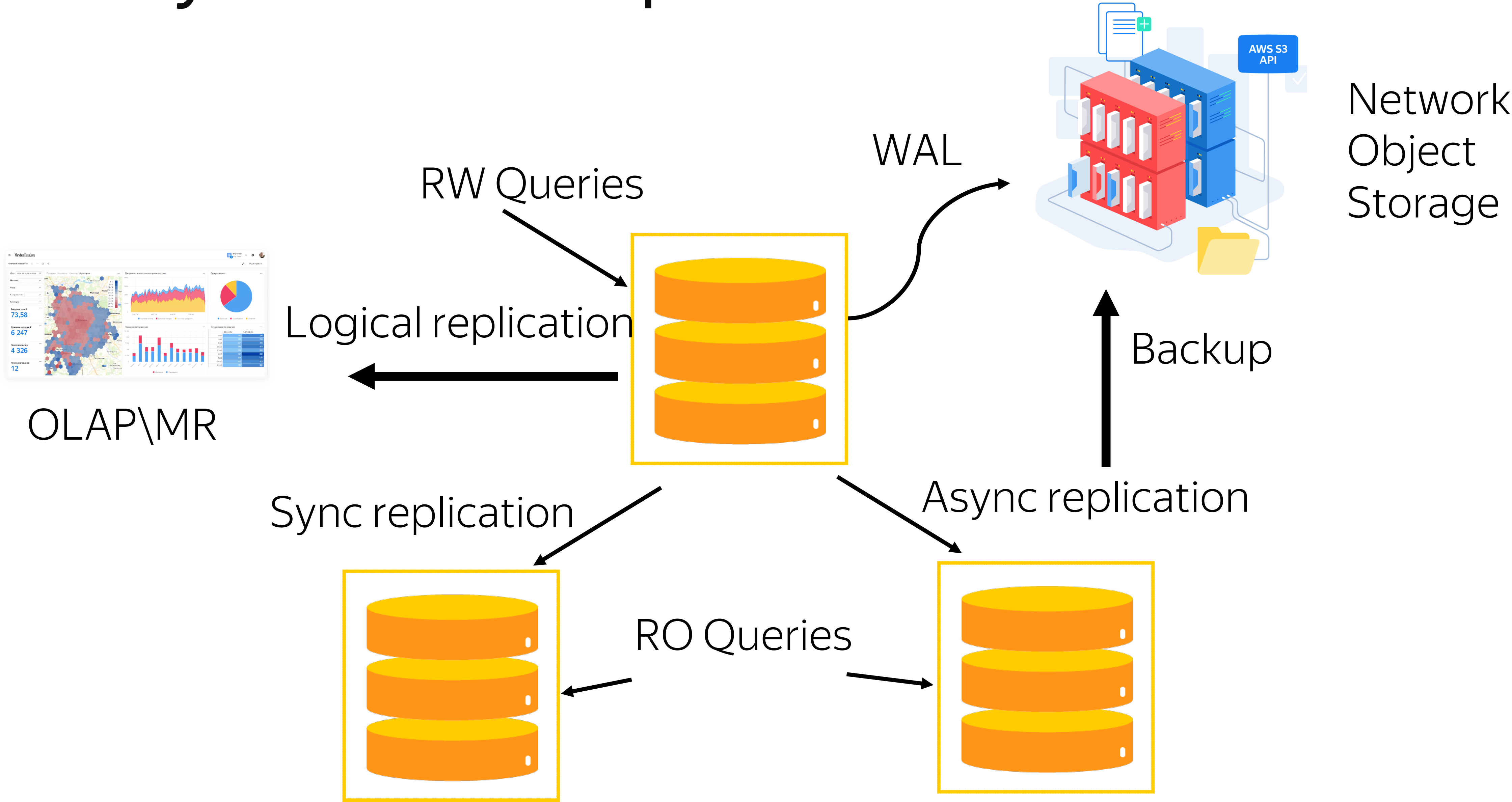
Make sure network timeouts work

- `tcp_user_timeout`

Libpq have some infinite timeouts relying on keepalives

- `keepalives_count`, `keepalives_interval`, `keepalives_idle`

Maybe automate promotion?



HA orchestration

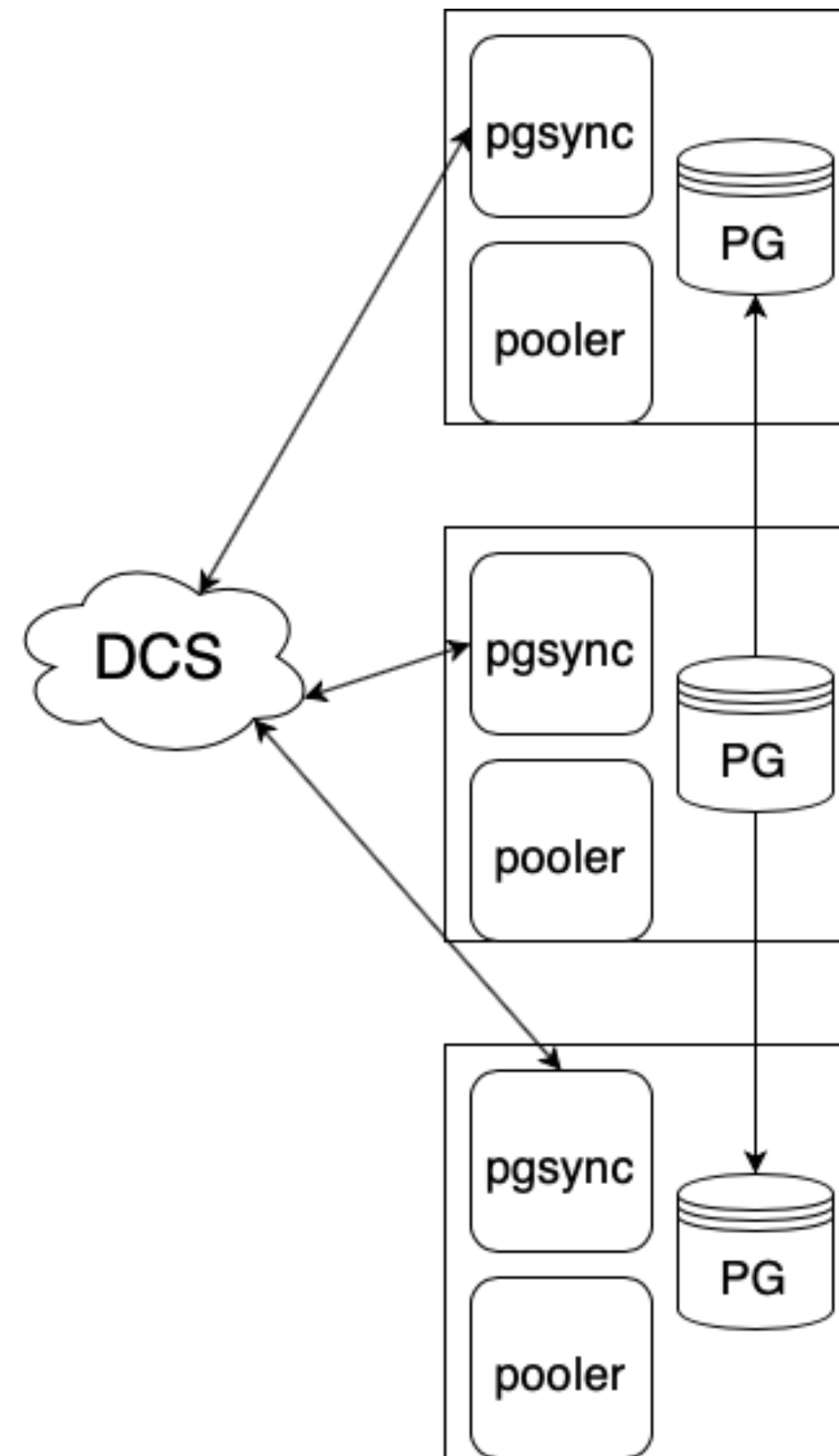
- › Patroni
- › Stolon
- › Repmgr

Yet another HA solution

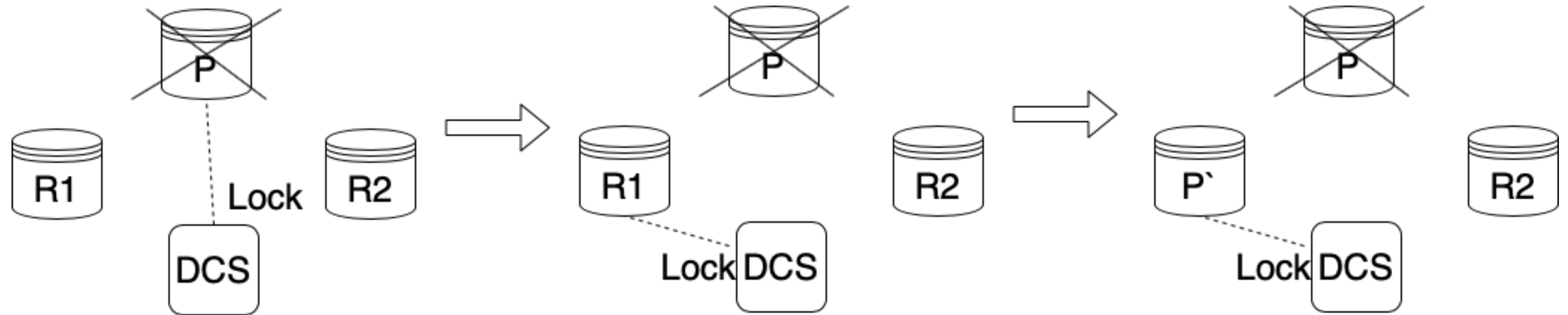


Bird's-eye view on pgsync

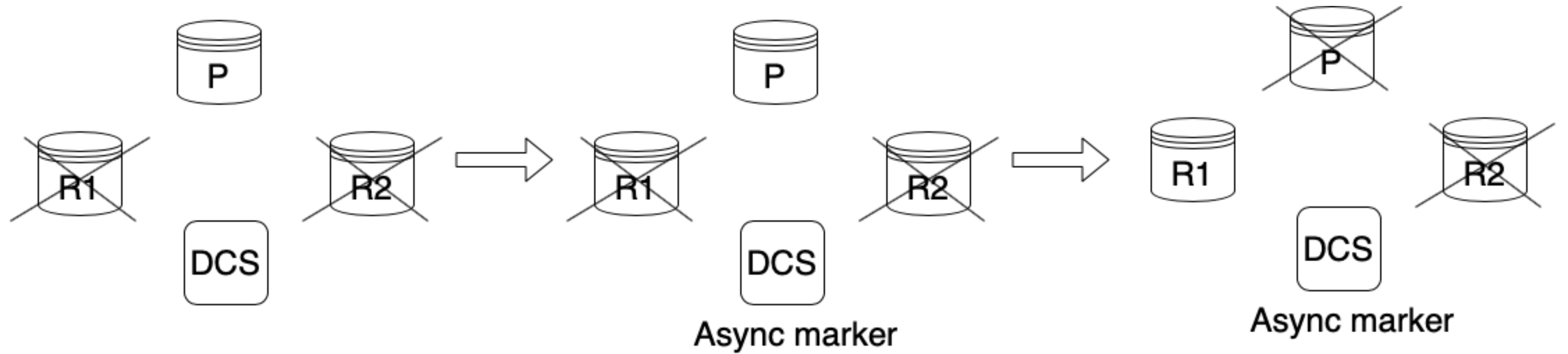
- › Shared DCS (2-3k pg clusters per 1 DCS cluster)
- › Agent on VM with PostgreSQL
- › Pooler for fencing



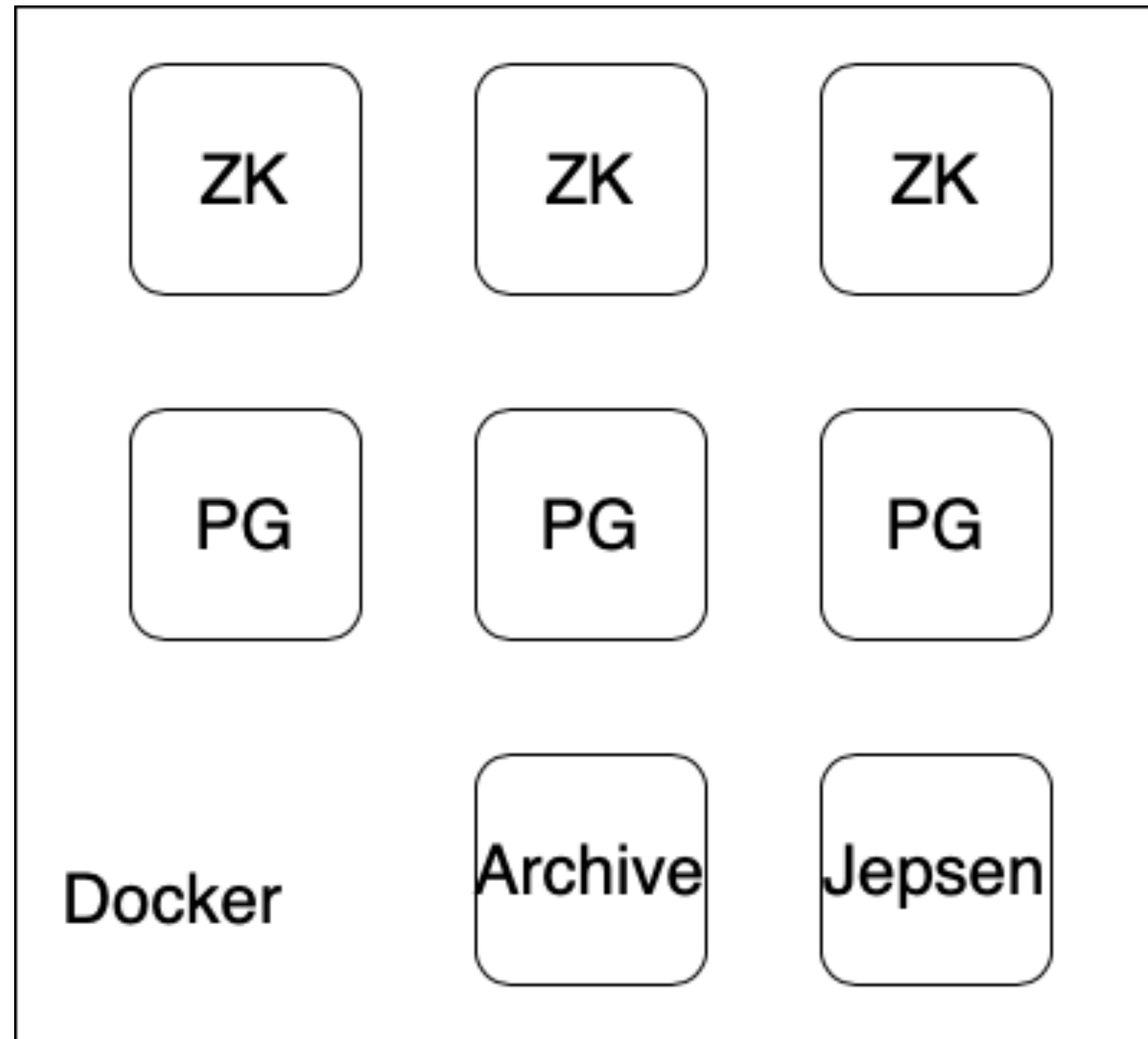
Primary failure



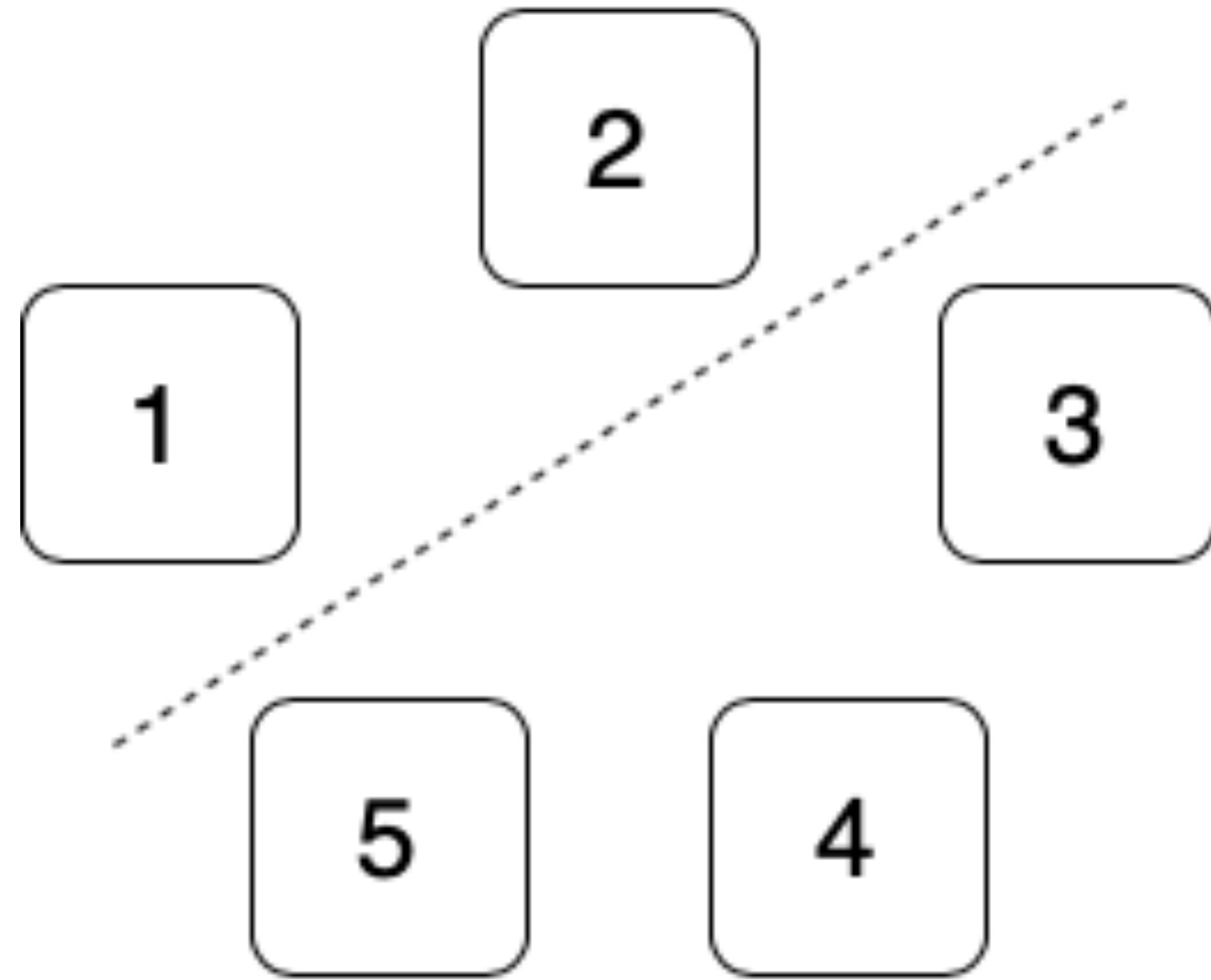
No failover?



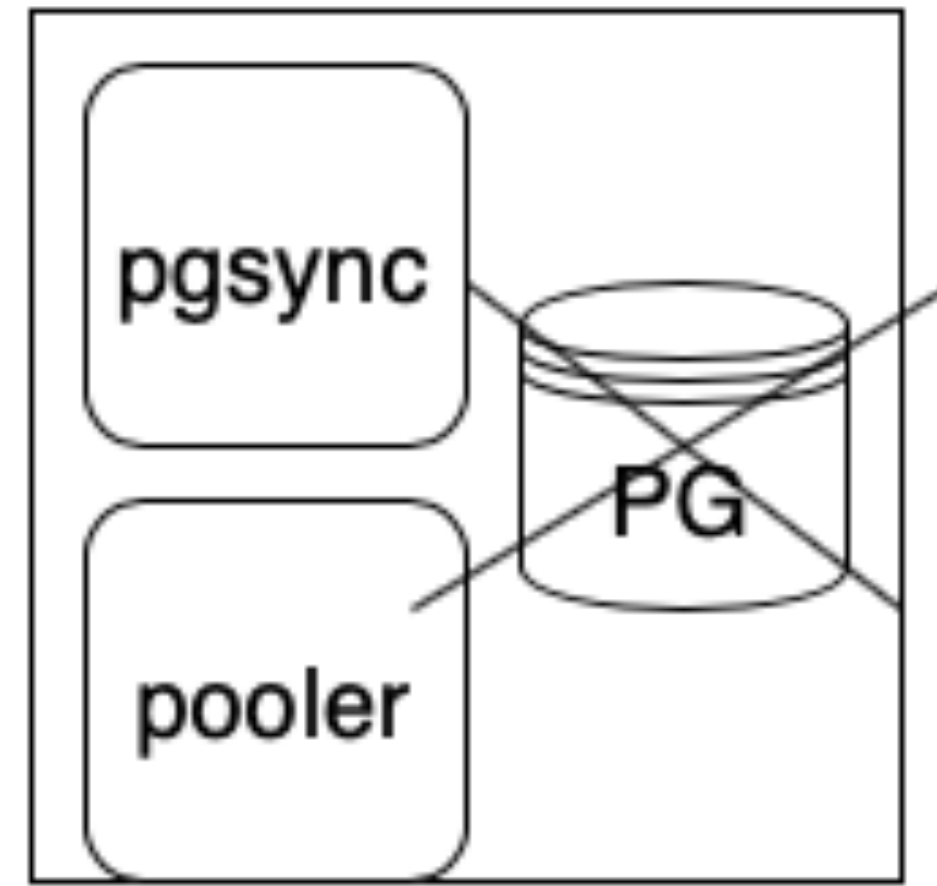
“Distributed fuzzing”



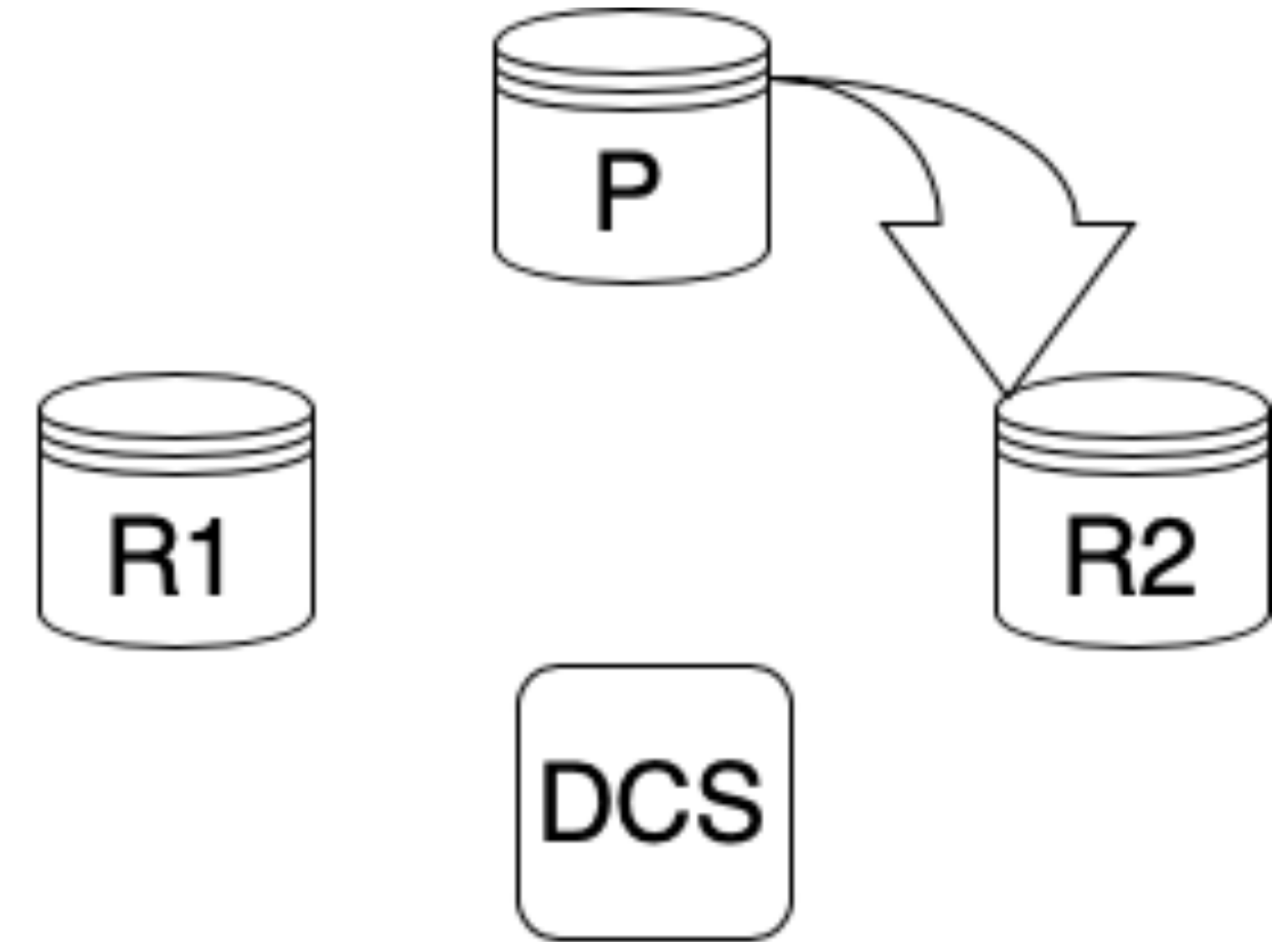
Nemesis



Partition



OOM



Switchover

Test workload

› CREATE TABLE test (value bigint PRIMARY KEY);

try

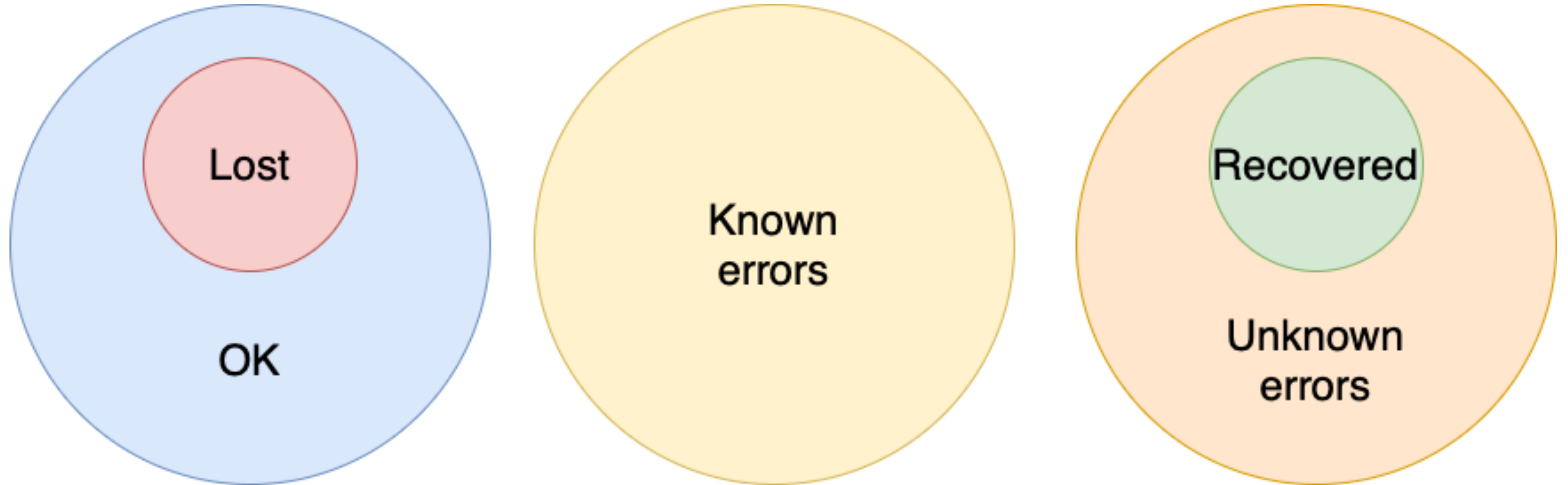
execute “INSERT INTO test VALUES (:counter);”

remember “ok”

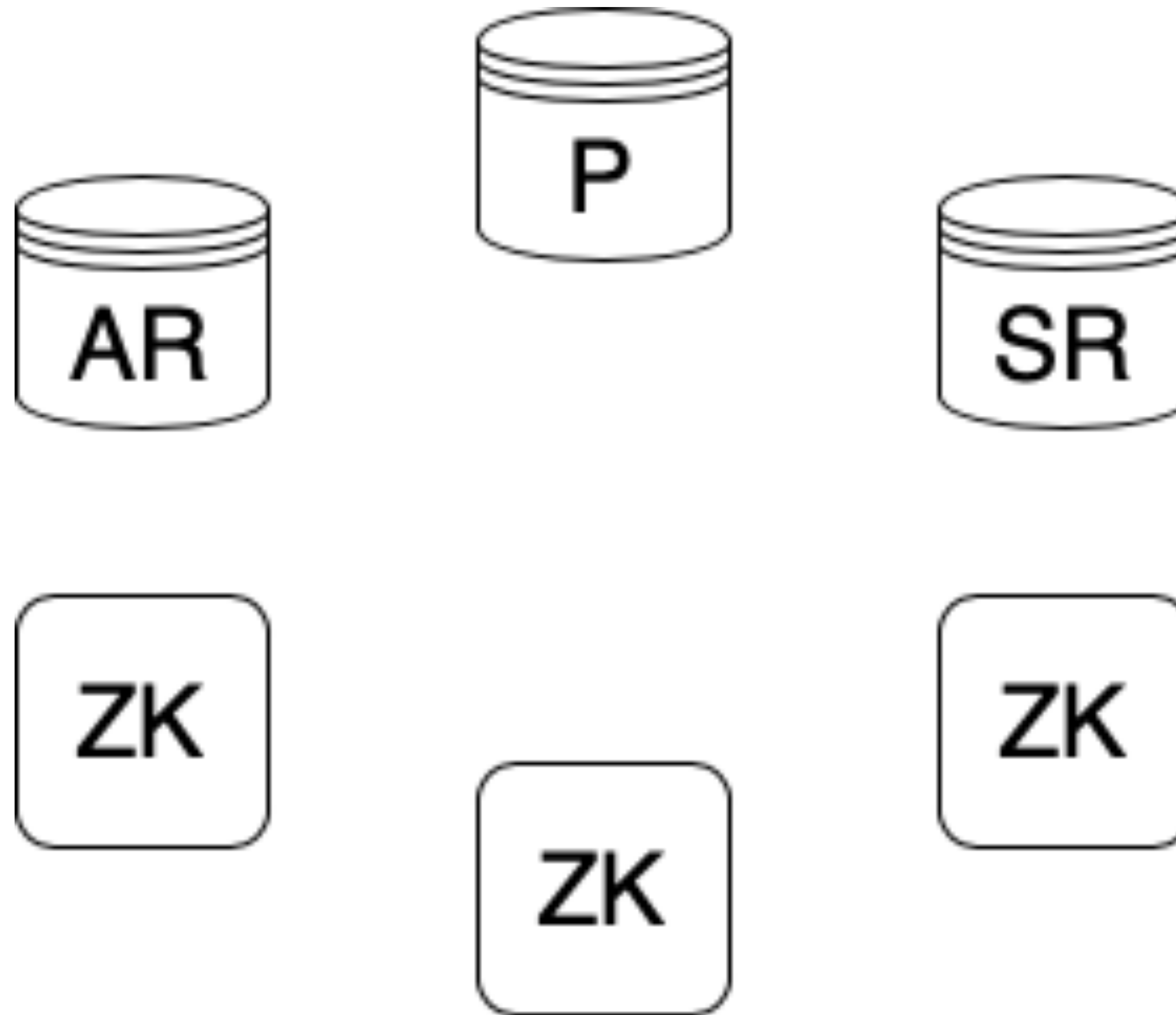
catch Throwable

remember “error”

Result sets



Why '*' is not always the best value for synchronous_standby_names



Primary/replica loops

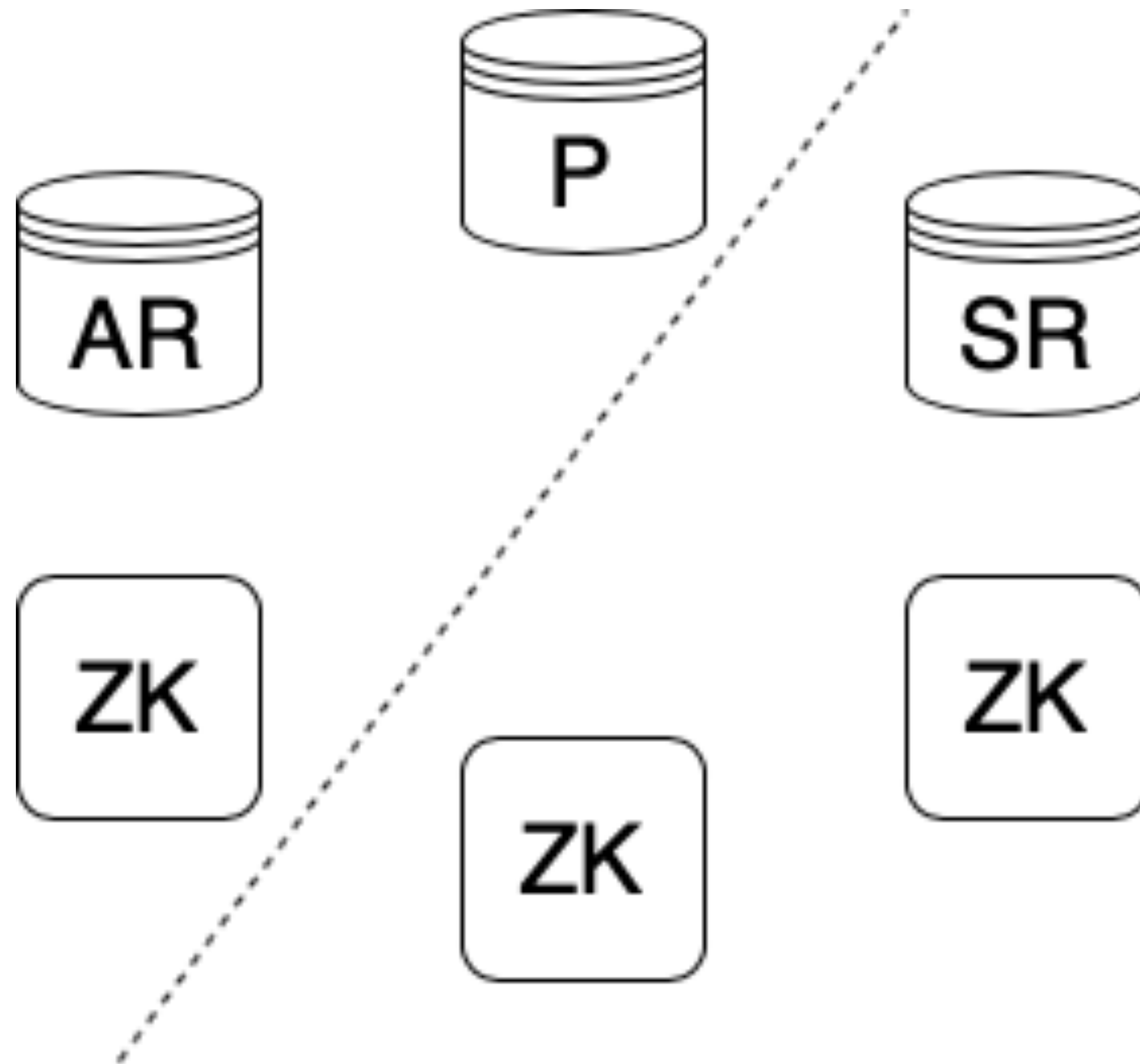
| Primary

```
if has_leader_lock():  
    replicas_state = execute("SELECT * FROM pg_stat_replication")  
    save(replicas_state)
```

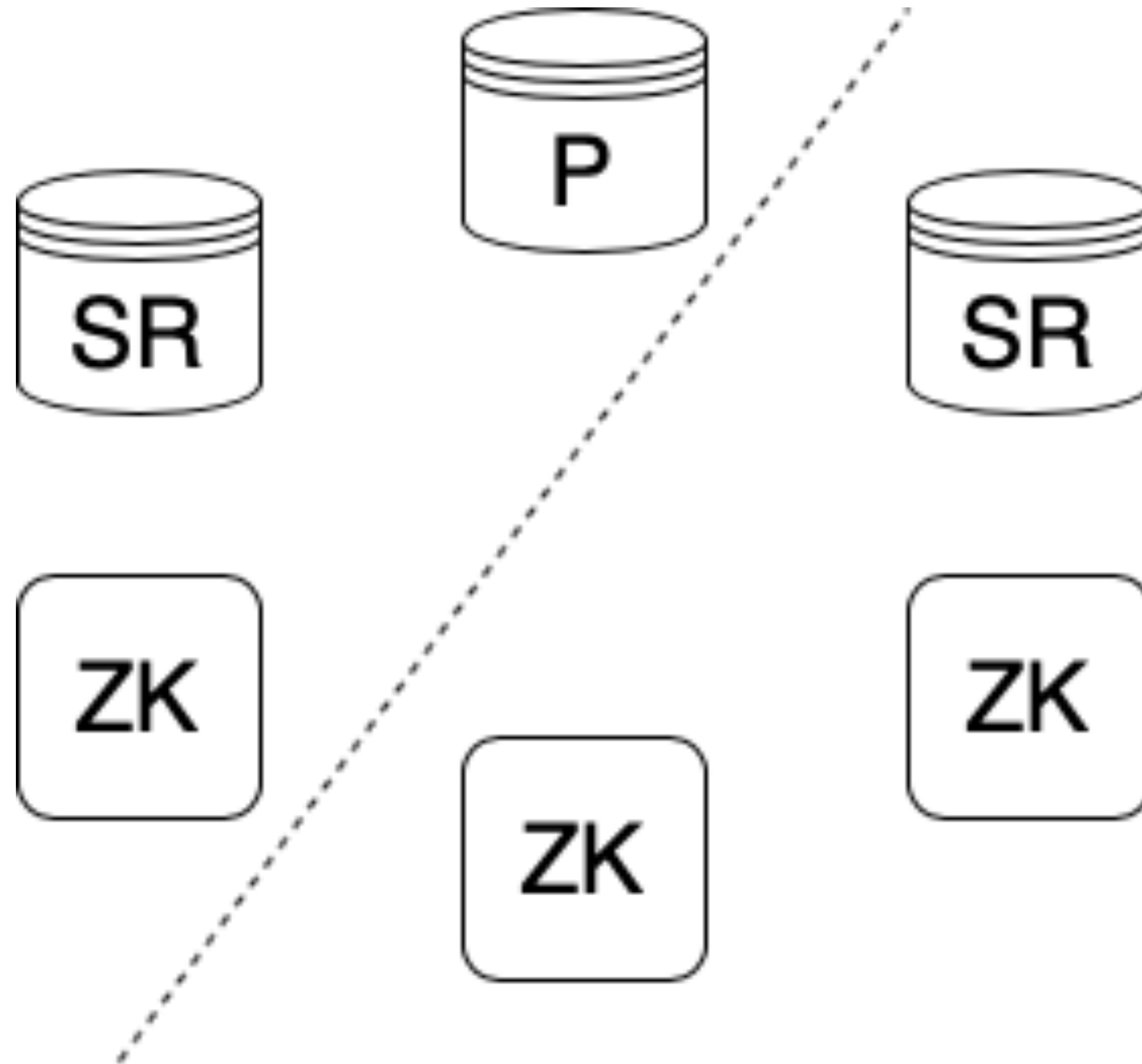
| Replica

```
if not leader_lock_holder():  
    saved_state = get_replicas_state()  
    if saved_state.get(my_application_name)['sync_state'] == 'sync':  
        take_lock_and_promote()
```

Partition



Partition



Fixed replica loop

```
try_take_sync_lock()
```

```
if not leader_lock_holder():
```

```
    saved_state = get_replicas_state()
```

```
    if saved_state.get(my_application_name)['sync_state'] == 'sync':
```

```
        if has_sync_lock():
```

```
            take_lock_and_promote()
```

Fixed primary loop

```
if has_leader_lock():
```

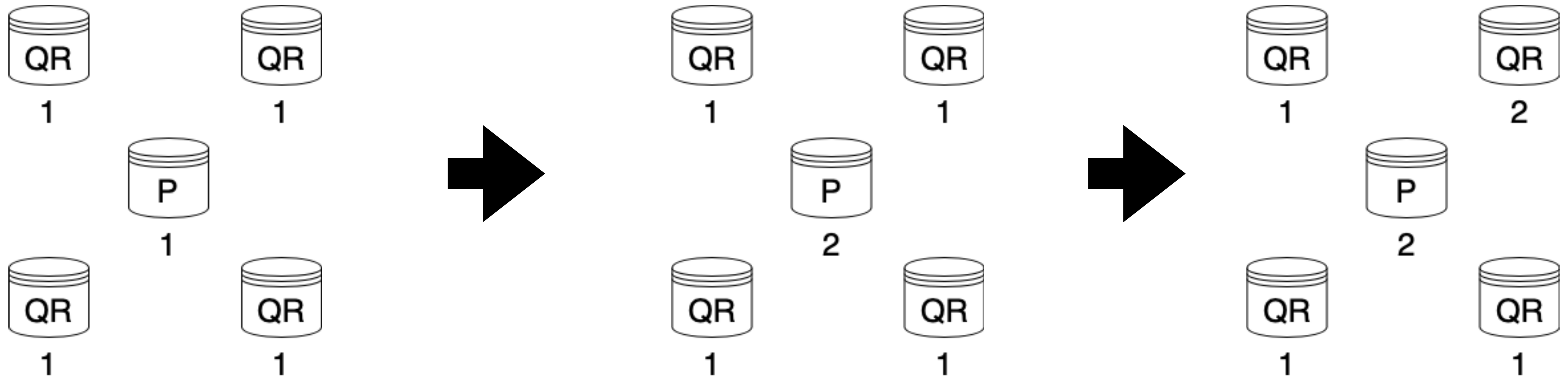
```
    sync_holder = get_sync_lock_holder()
```

```
    fix_standby_names(sync_holder)
```

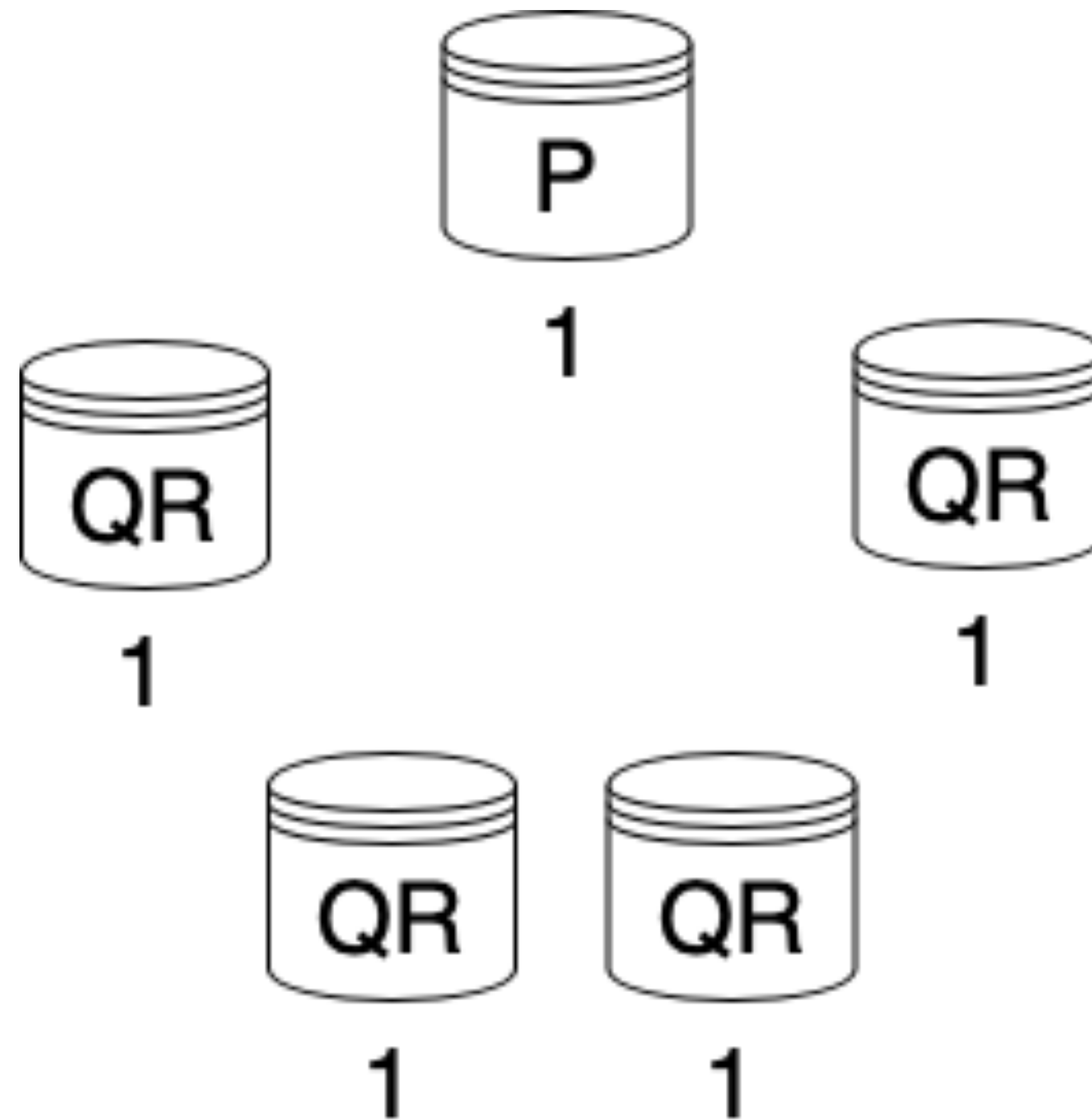
```
    replicas_state = execute("SELECT * FROM pg_stat_replication")
```

```
    save(replicas_state)
```

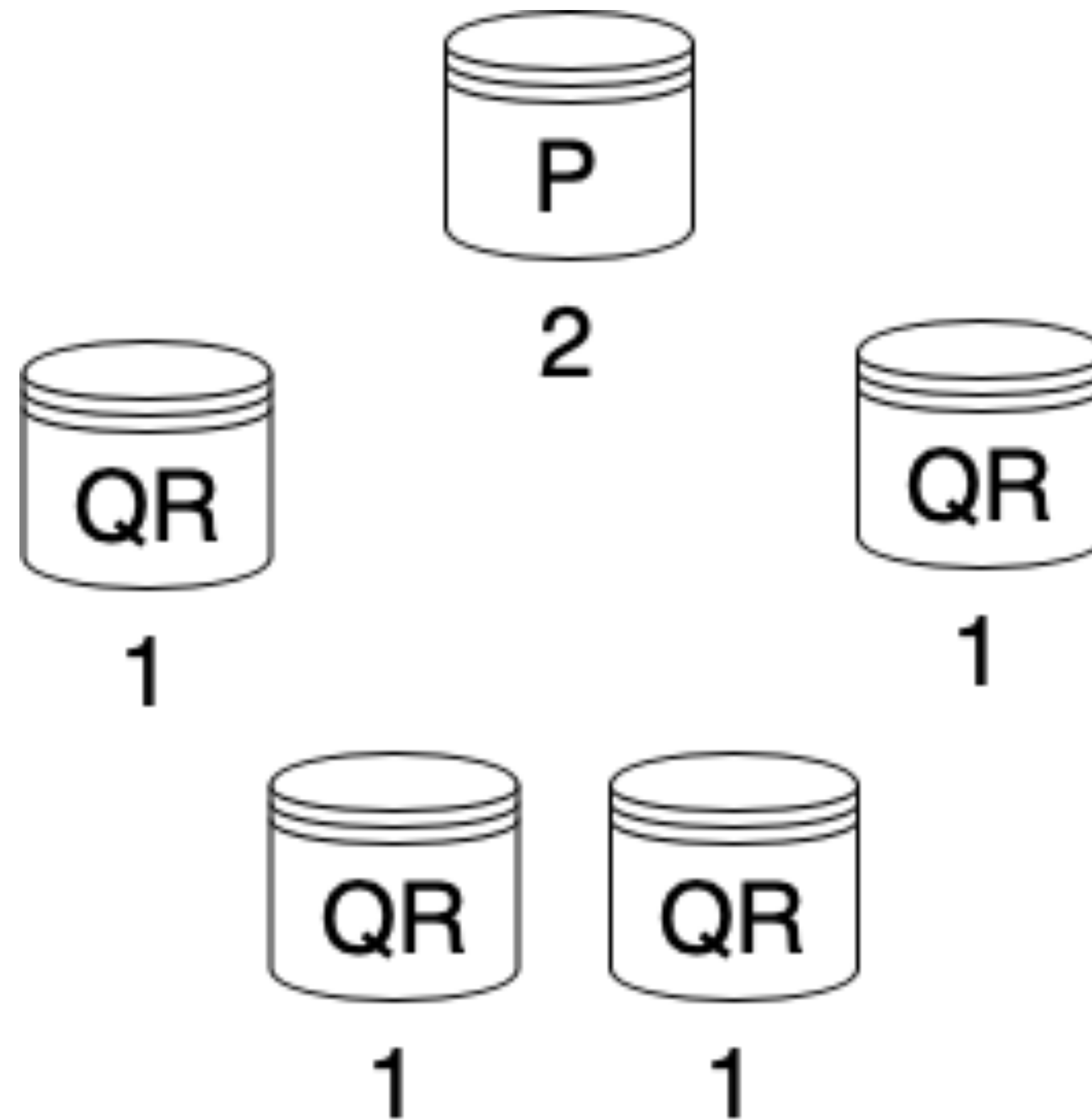
Quorum replication overview



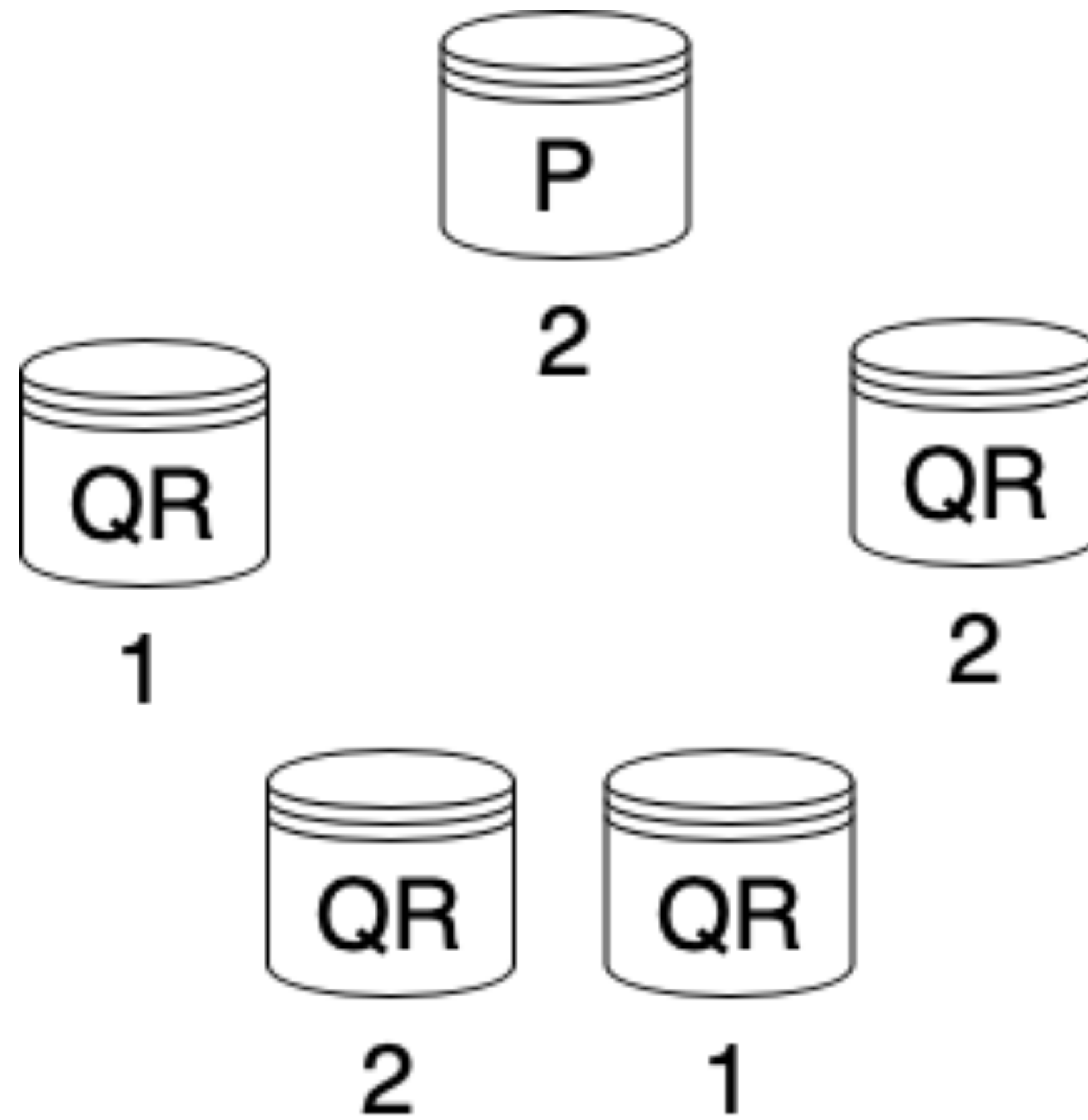
Quorum replication failover



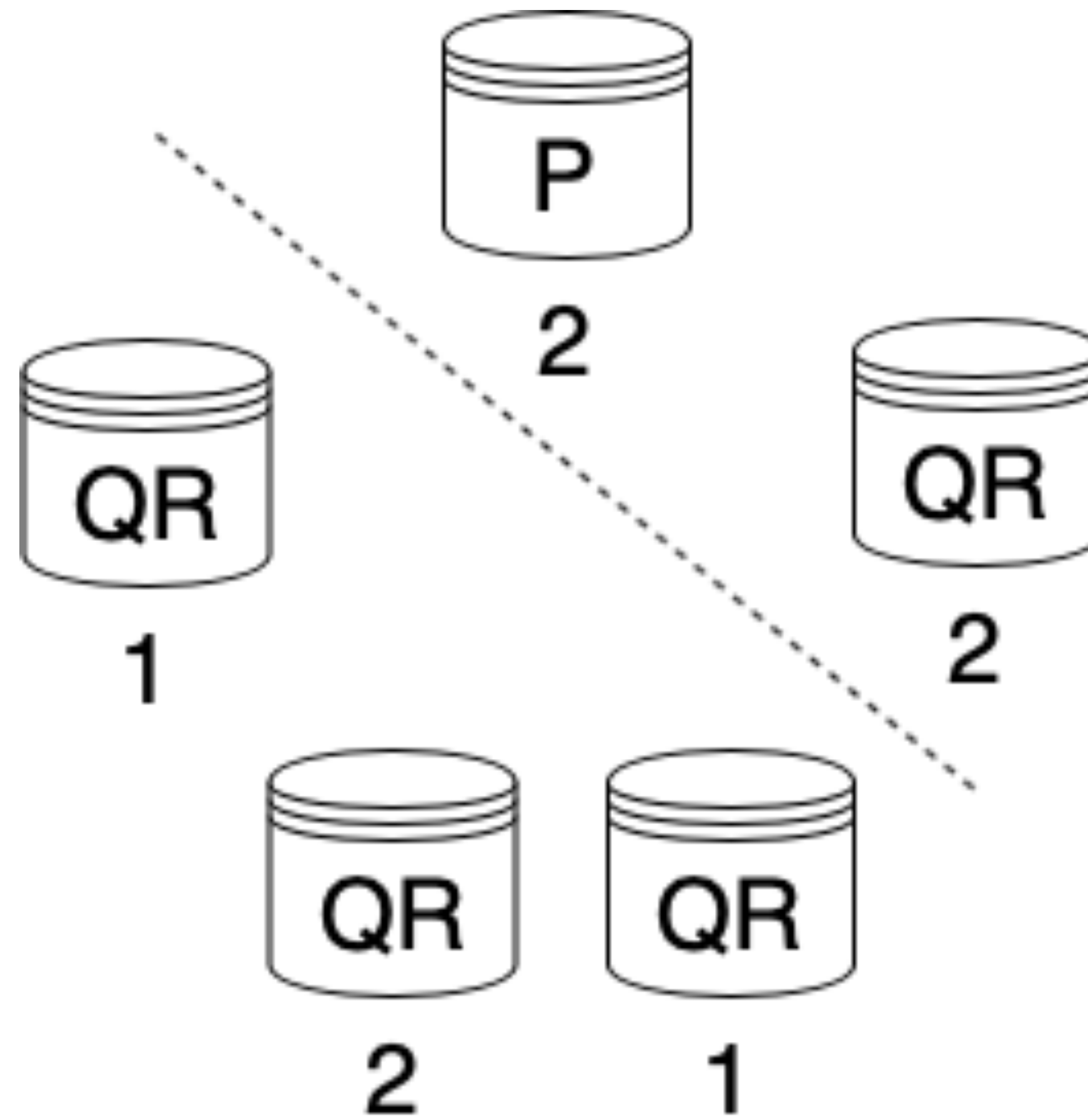
Quorum replication failover



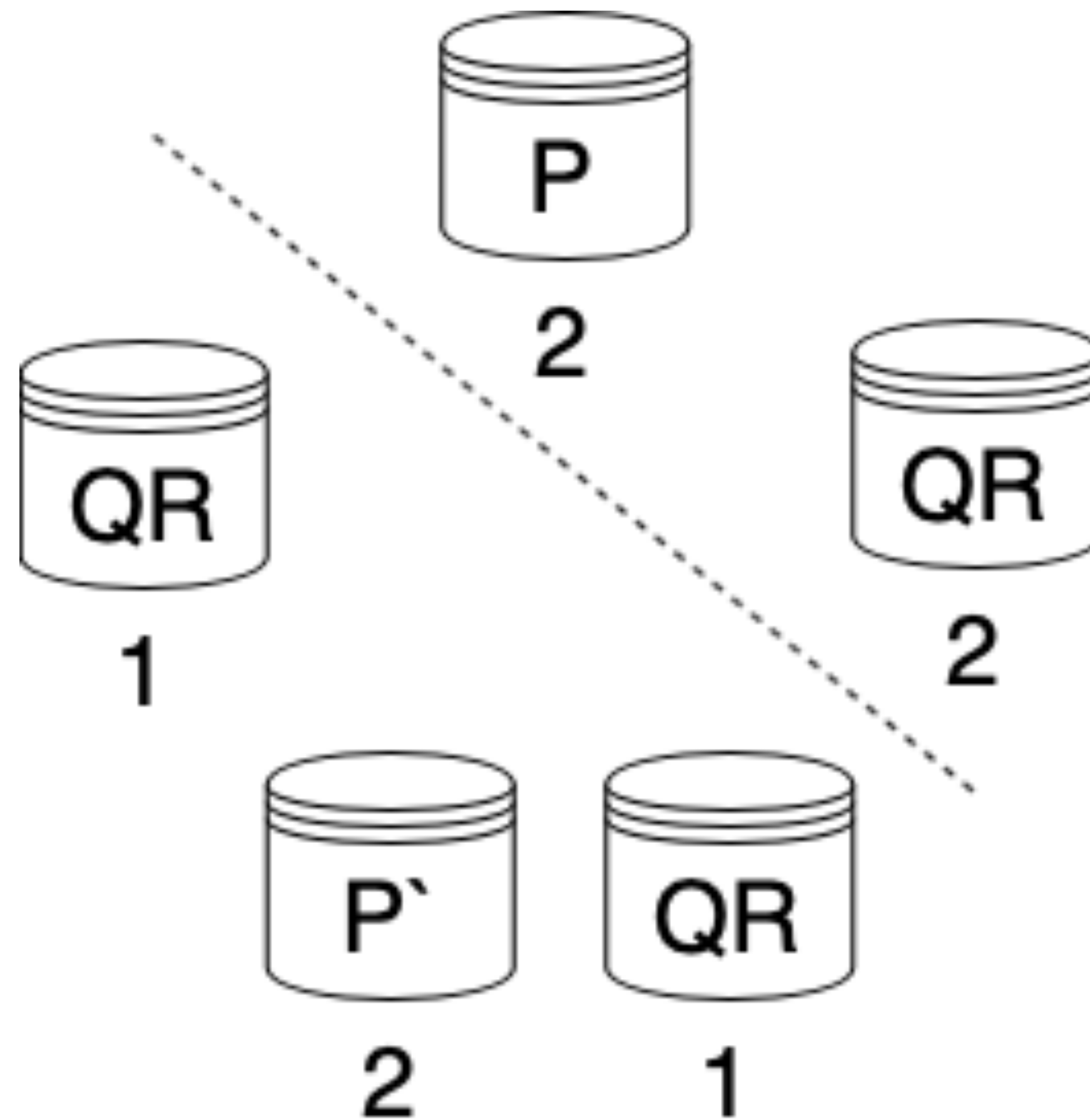
Quorum replication failover



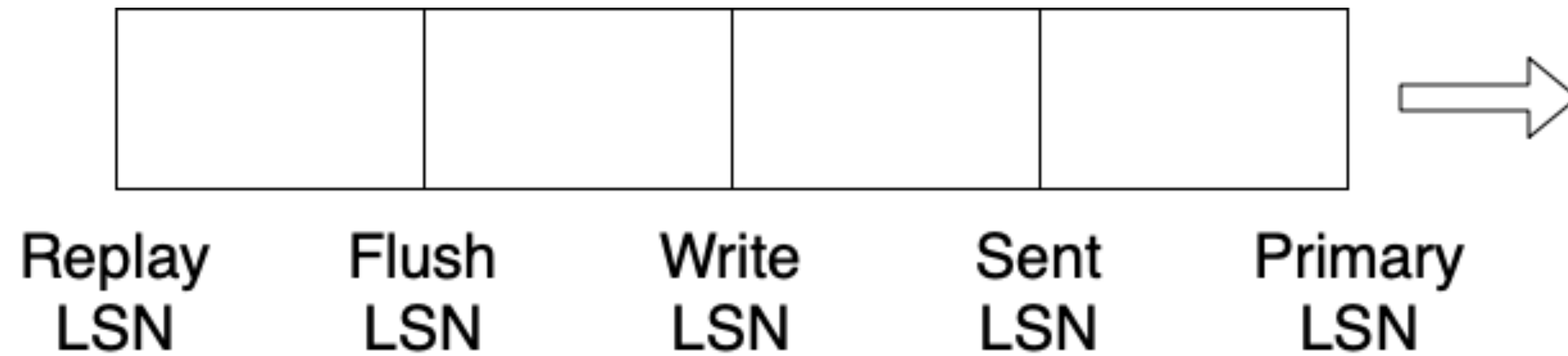
Quorum replication failover



Quorum replication failover



So many LSNs



Unforeseen consequences

- › `SELECT pg_last_wal_receive_lsn() -> 0/403F482`
- › `SELECT pg_last_wal_replay_lsn() -> 0/403F482`

—— Restart PostgreSQL ——

- › `SELECT pg_last_wal_receive_lsn() -> 0/40000000`
- › `SELECT pg_last_wal_replay_lsn() -> 0/403F482`

Just read the WAL

```
start = get_replay_lsn() # GetXLogReplayRecPtr
```

```
state['lsn'] = start
```

```
while read_wal(state): # XLogReadRecord
```

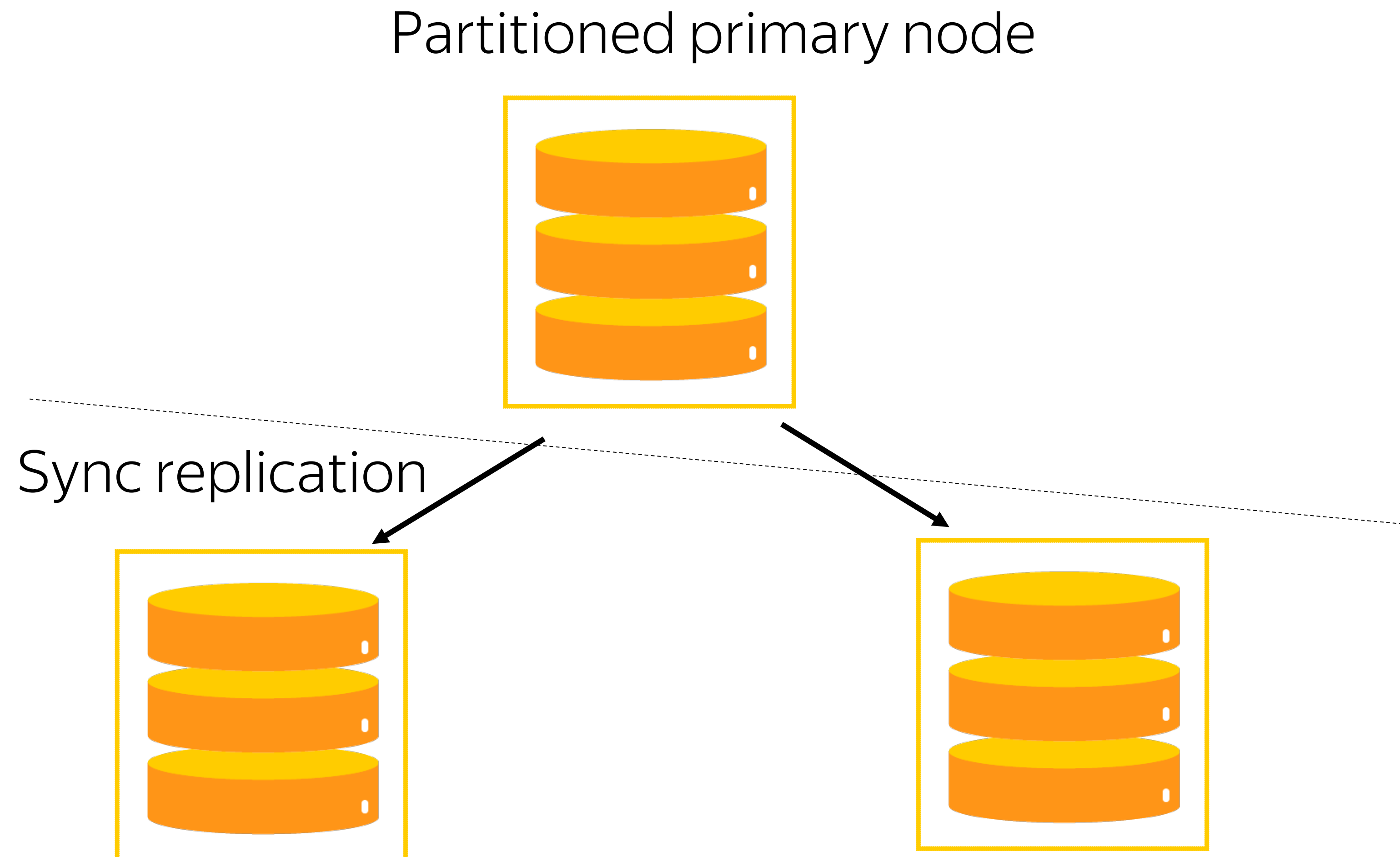
```
    pass
```

```
return state['lsn']
```

Canceling running query



Synchronous commit guarantees



Transaction pseudocode

Lock data

Modify data locally

Wait for

- › WAL flush locally
- › WAL shipment remotely

Unlock data for observer

Replicas can be inconsistent

Lock data

Modify data locally

Wait for

- › WAL flush locally
- › WAL shipment remotely

Client read data on Standby that is not observed on Primary

Unlock data for observer

Replicas can be inconsistent

Lock data

Modify data locally

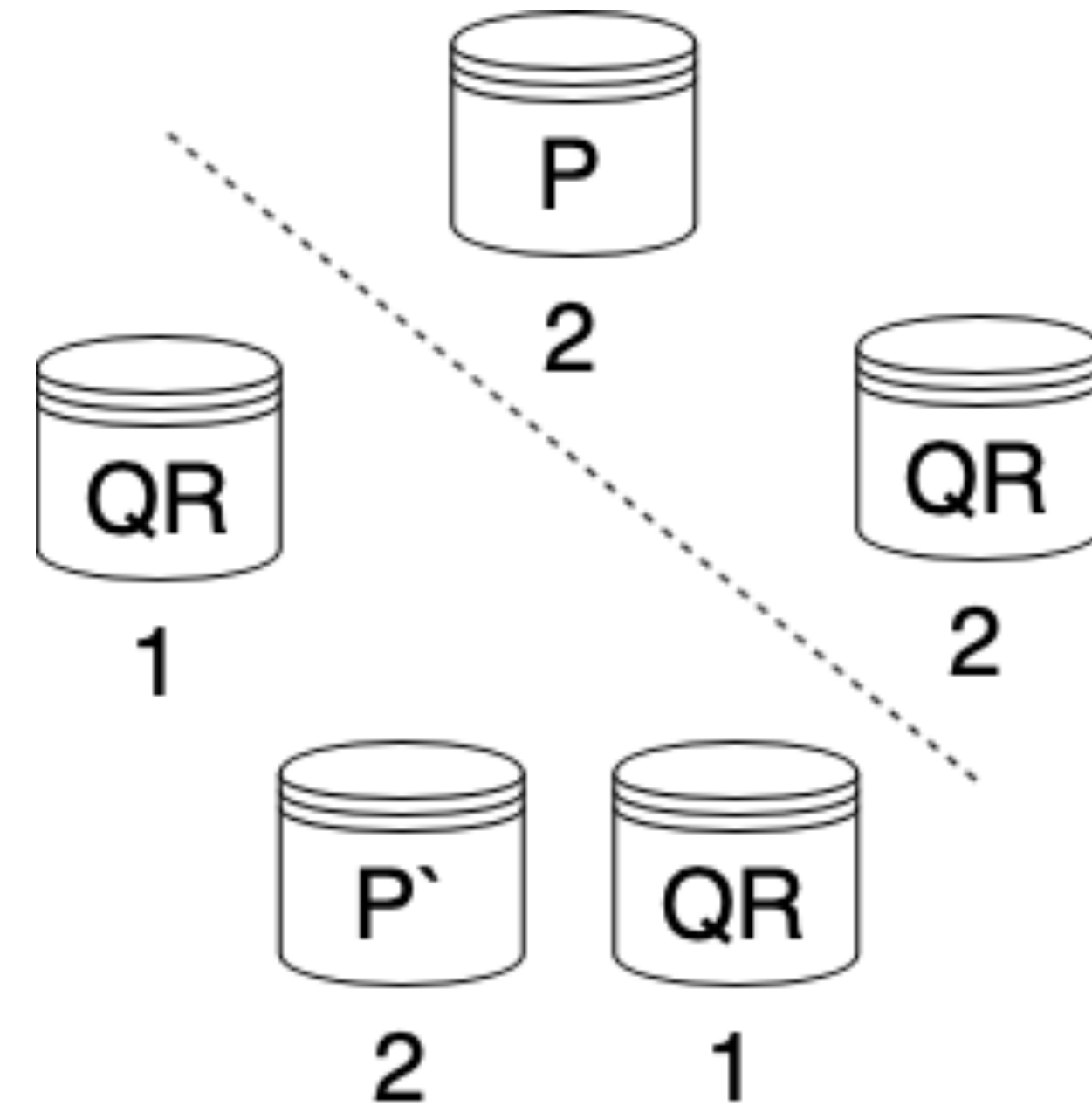
Wait for

- › WAL flush locally
- › WAL shipment remotely

Client read data on Standby that is not observed on Primary

Even on failed part of quorum

Unlock data for observer



Acknowledge not replicated data

Primary is partitioned

Lock data

Modify data locally

Wait for

- › WAL flush locally
- › WAL shipment remotely (hangs due to network partition or standby promotion)
- › Client cancels query, but it's committed locally

Unlock data for observer

User issues INSERT ON CONFLICT DO NOTHING

No WAL is written => commit is acknowledged

Disable cancellation of executed locally query

 ALTER SYSTEM SET synchronous_commit_cancelation to off;

<https://commitfest.postgresql.org/31/2402/>

Only partial solution

- Primary restart still makes not-replicated data visible

Some additional information

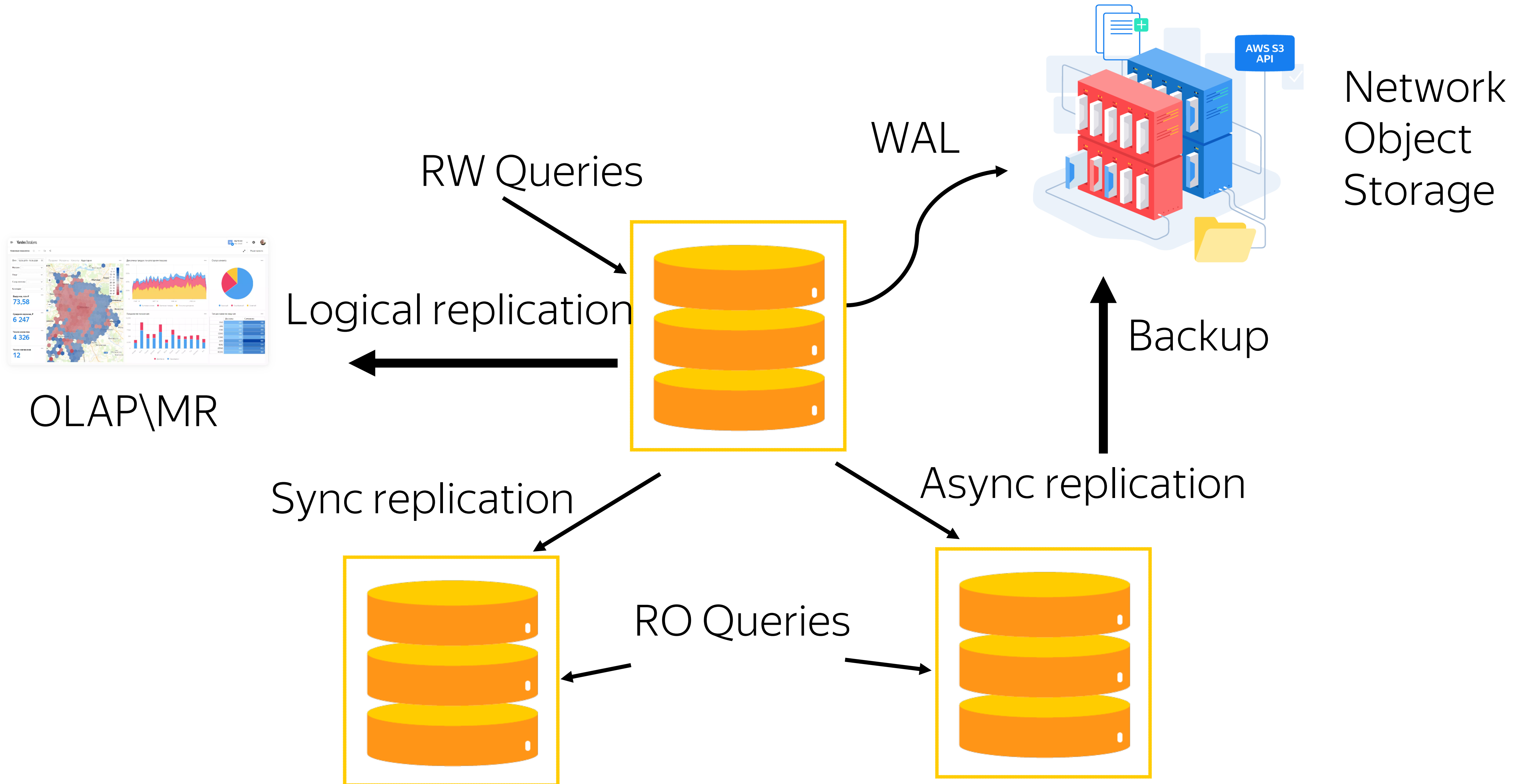
The topic was discussed at PGCon unconference 2020

https://wiki.postgresql.org/wiki/PgCon_2020_Developer_Unconference/Edge_cases_of_synchronous_replication_in_HA_solutions

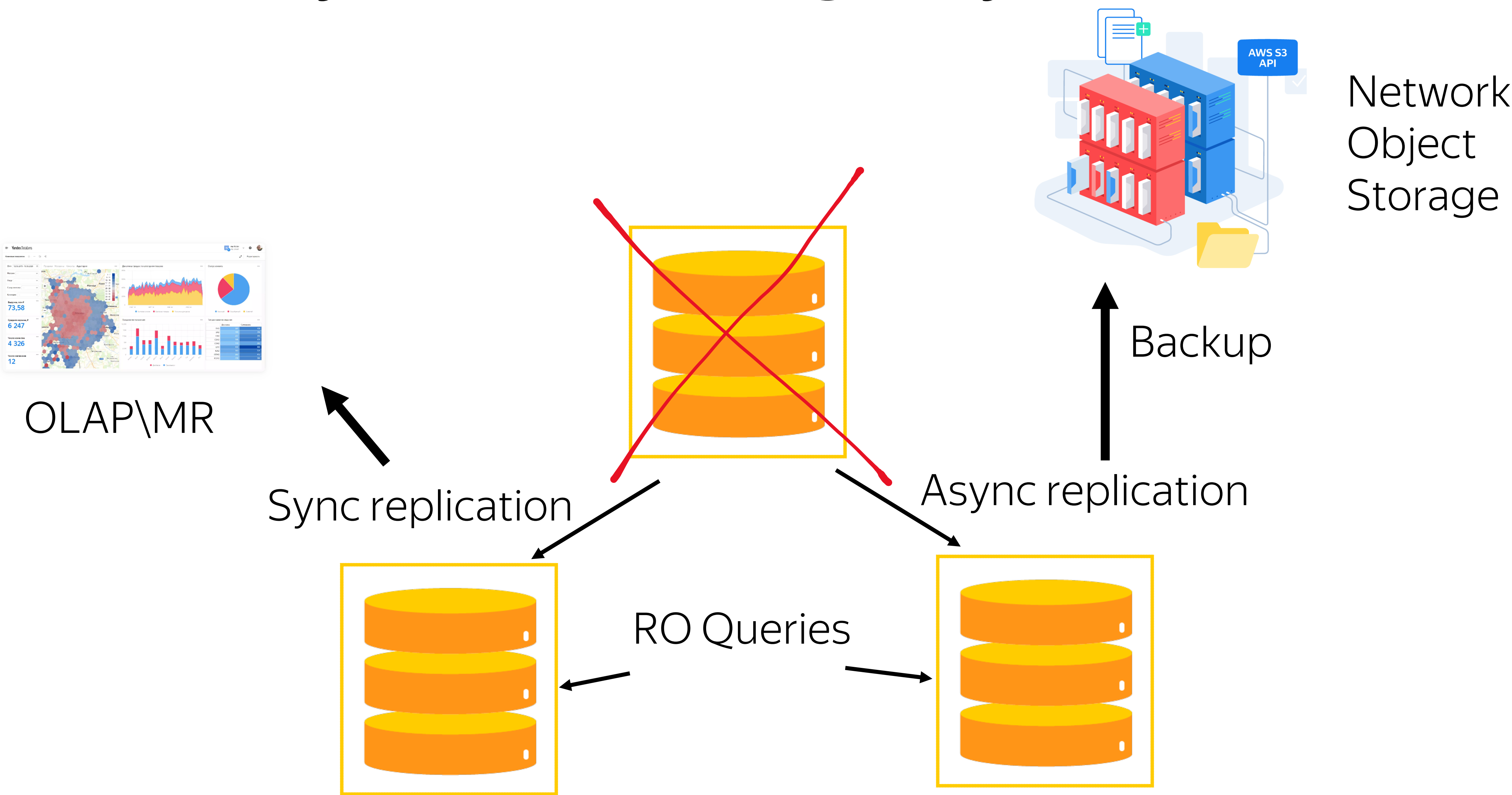
Changed data capture



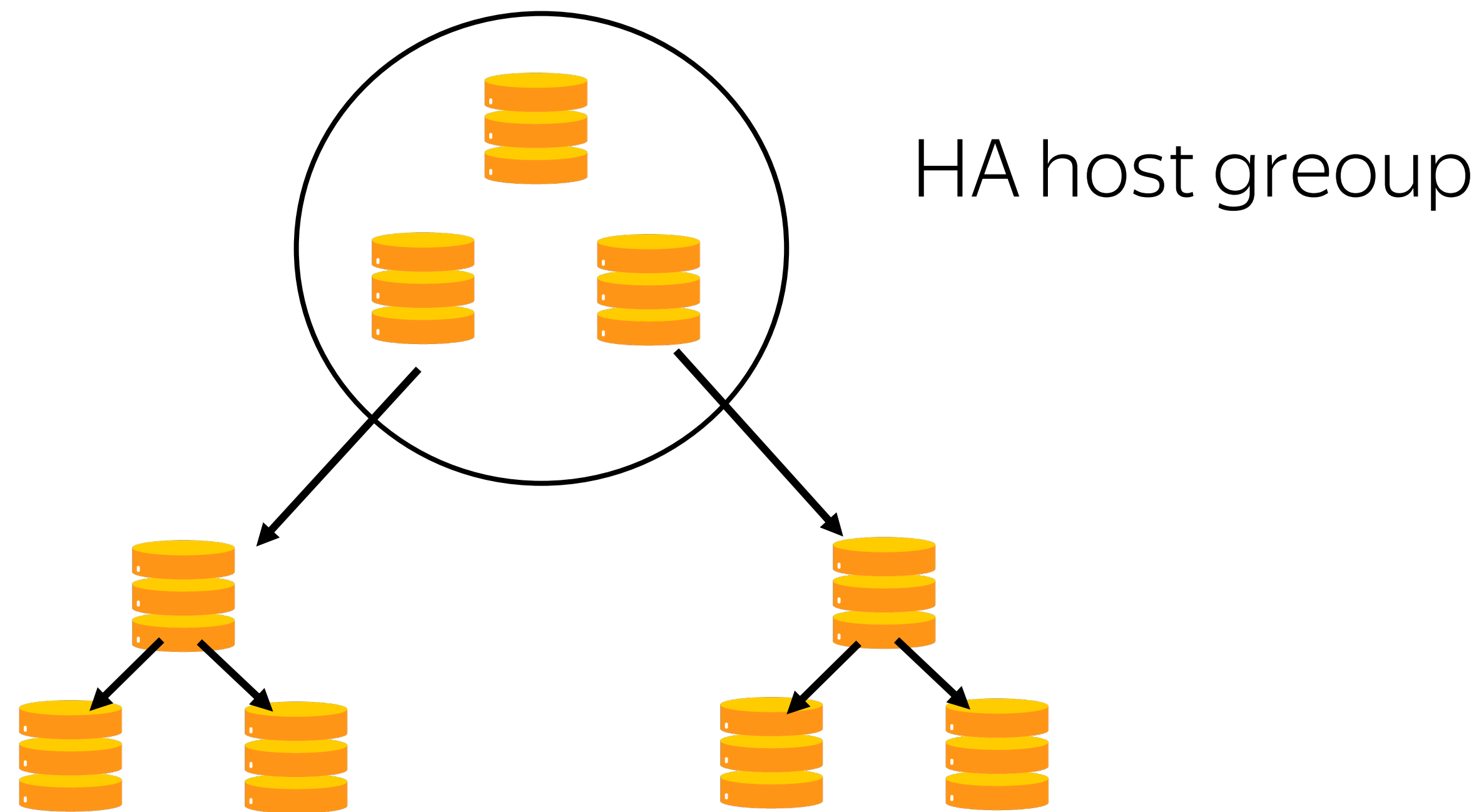
So you want to stream logically?



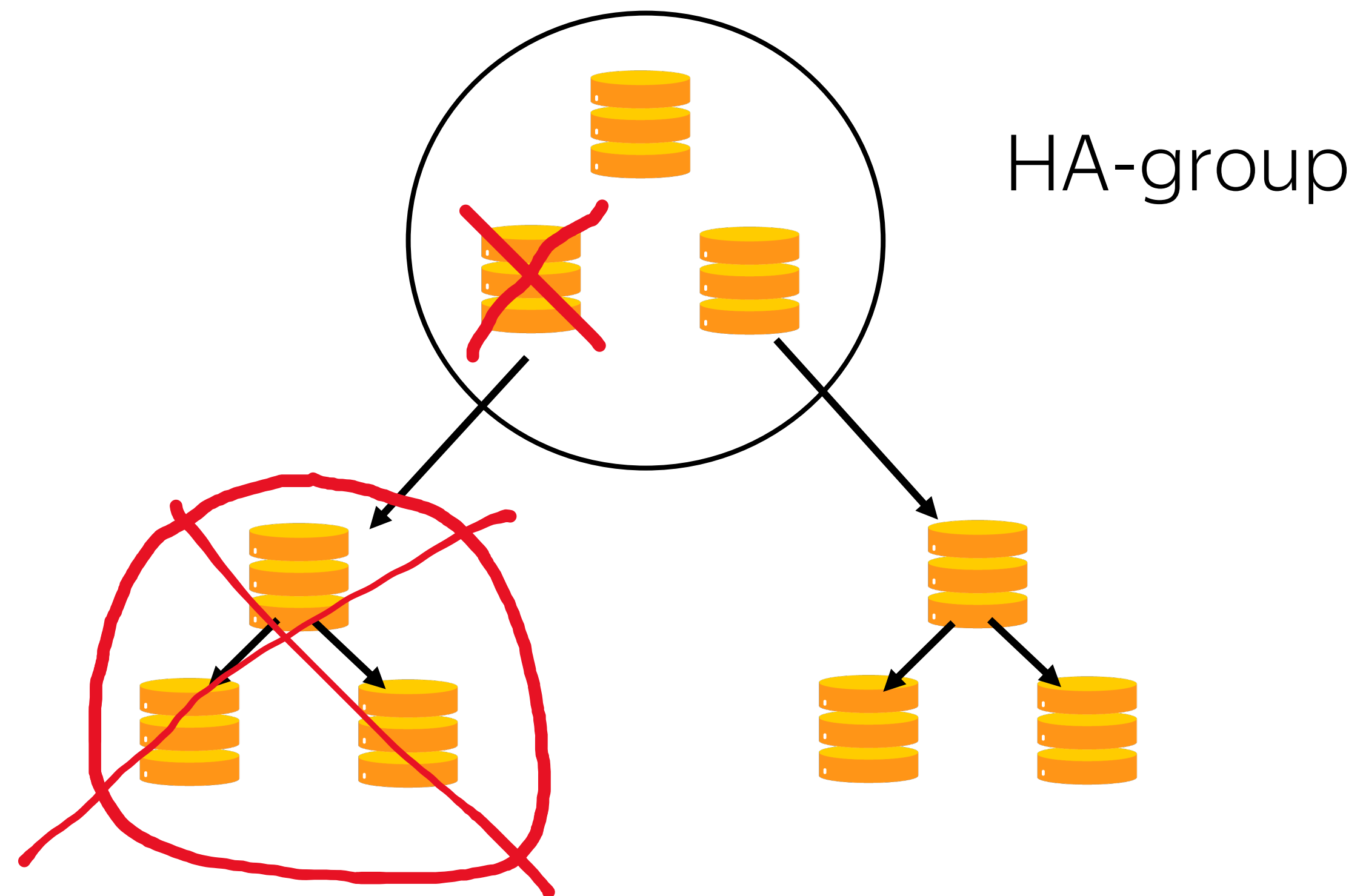
Be ready to failover. Logically.



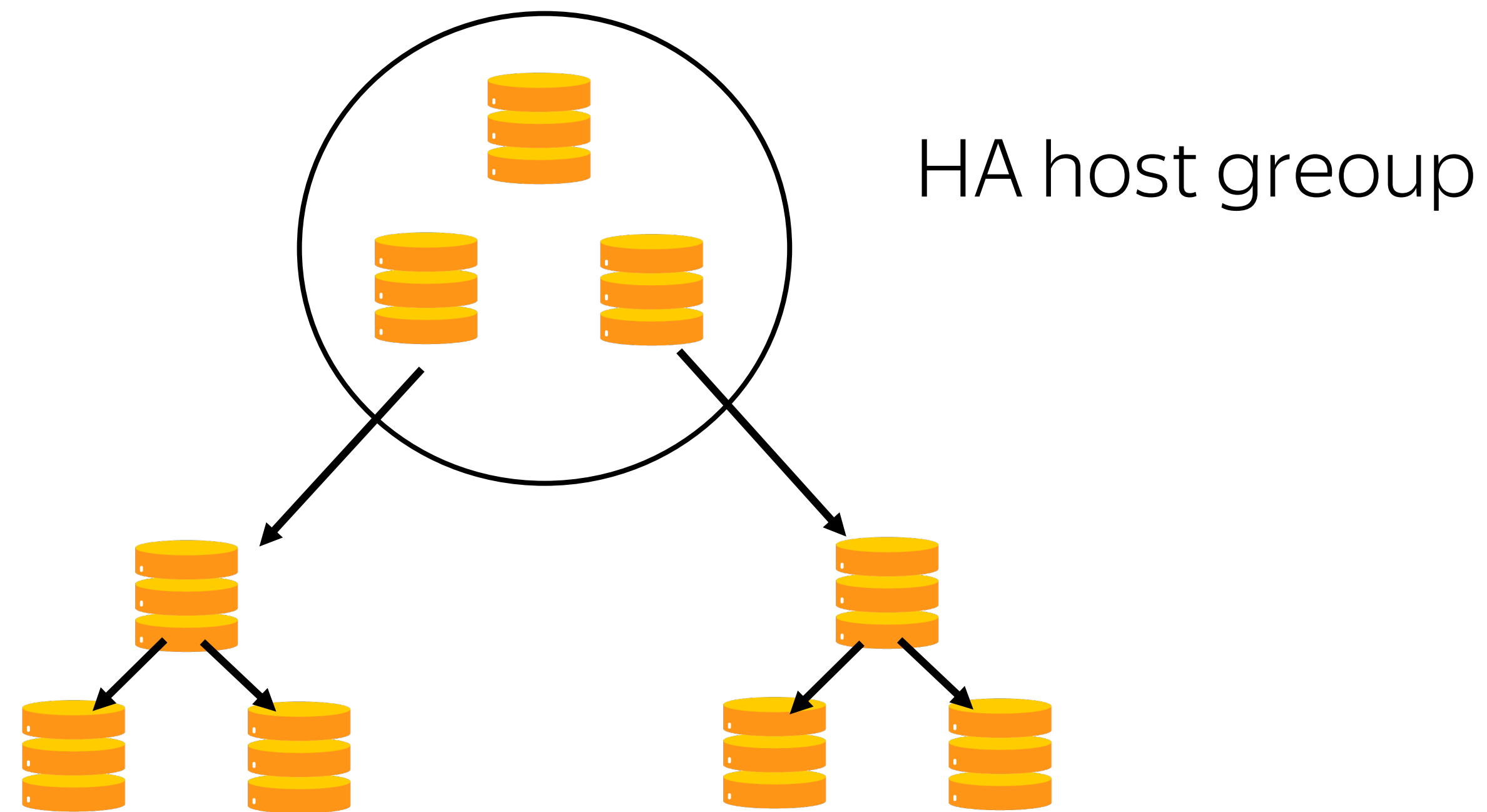
Non-HA standbys



Rebuilding topology



Non-HA standbys



Useless for CDC before logical decoding is allowed on Standby

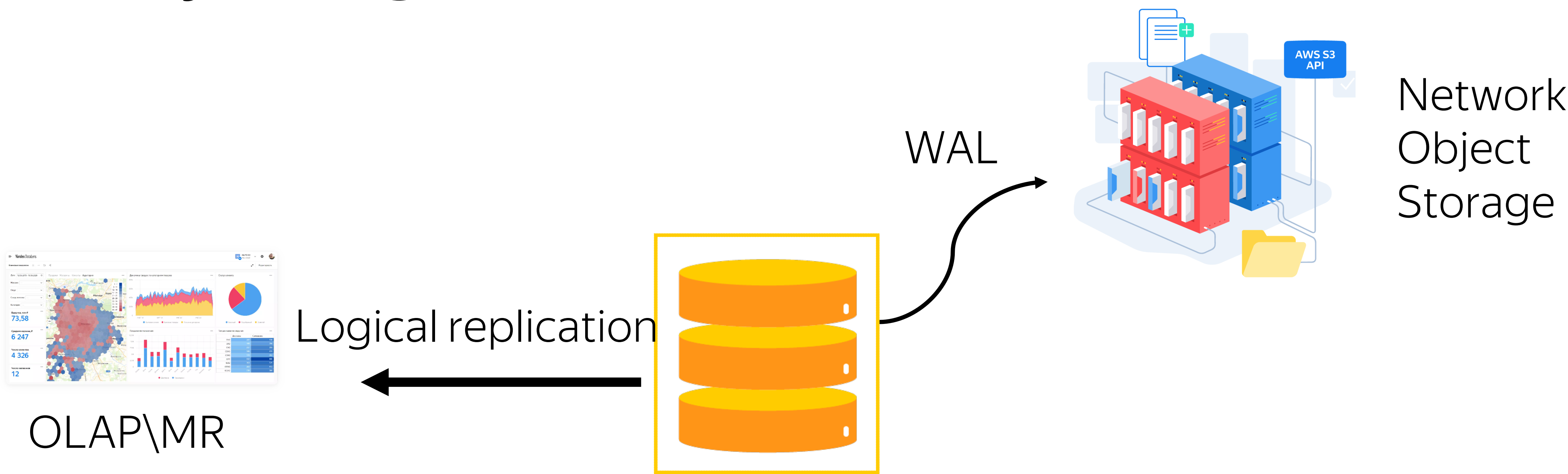
Slots are necessary

Logical streaming start from slot's restart position

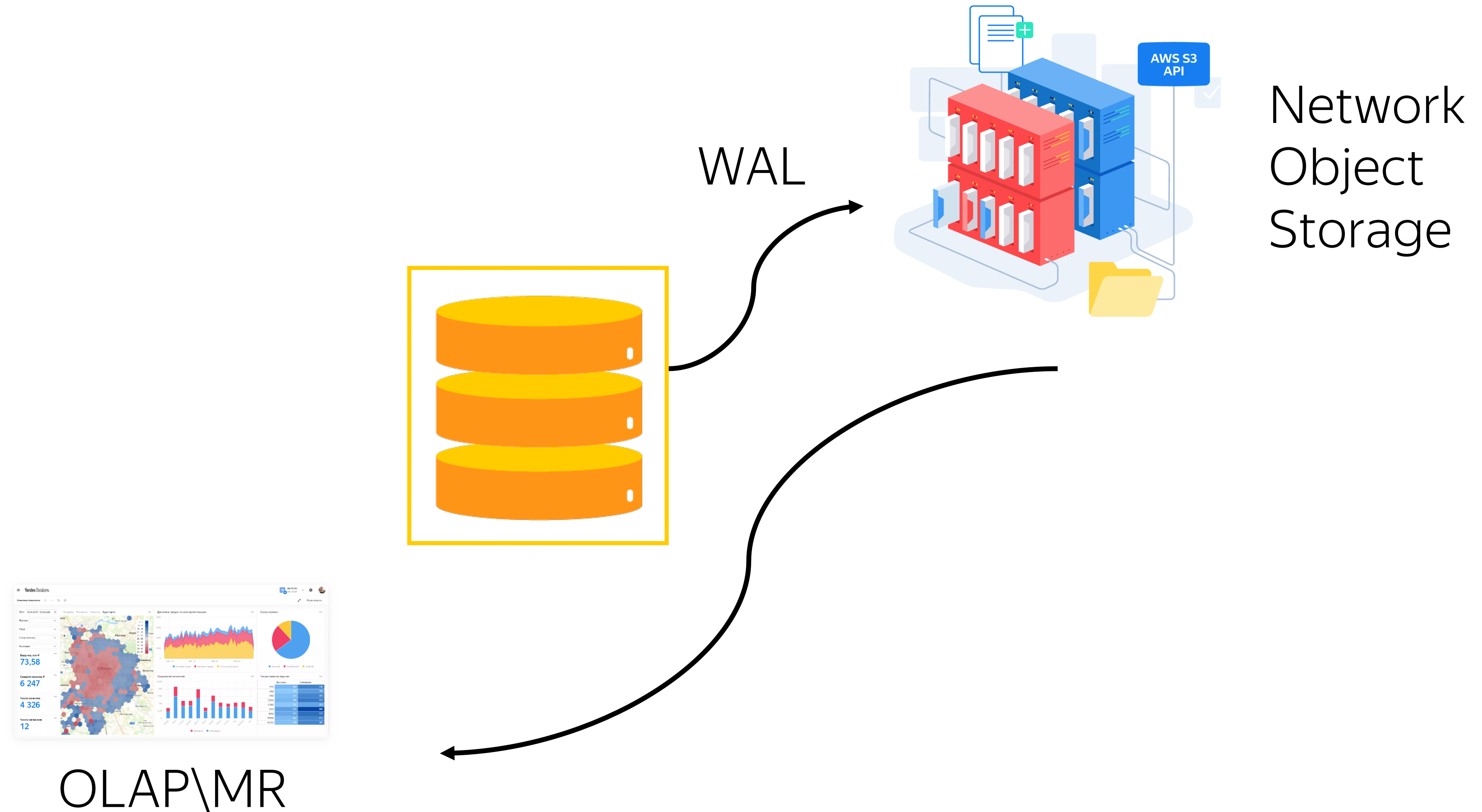
 But you can only create slot on latest WAL insert pointer

No chances to have LSN on failed primary == LSN of created slot

Maybe logical archive?



Maybe logical decoding in WAL-G?



WAL-G already parses WAL to make backups and replay faster

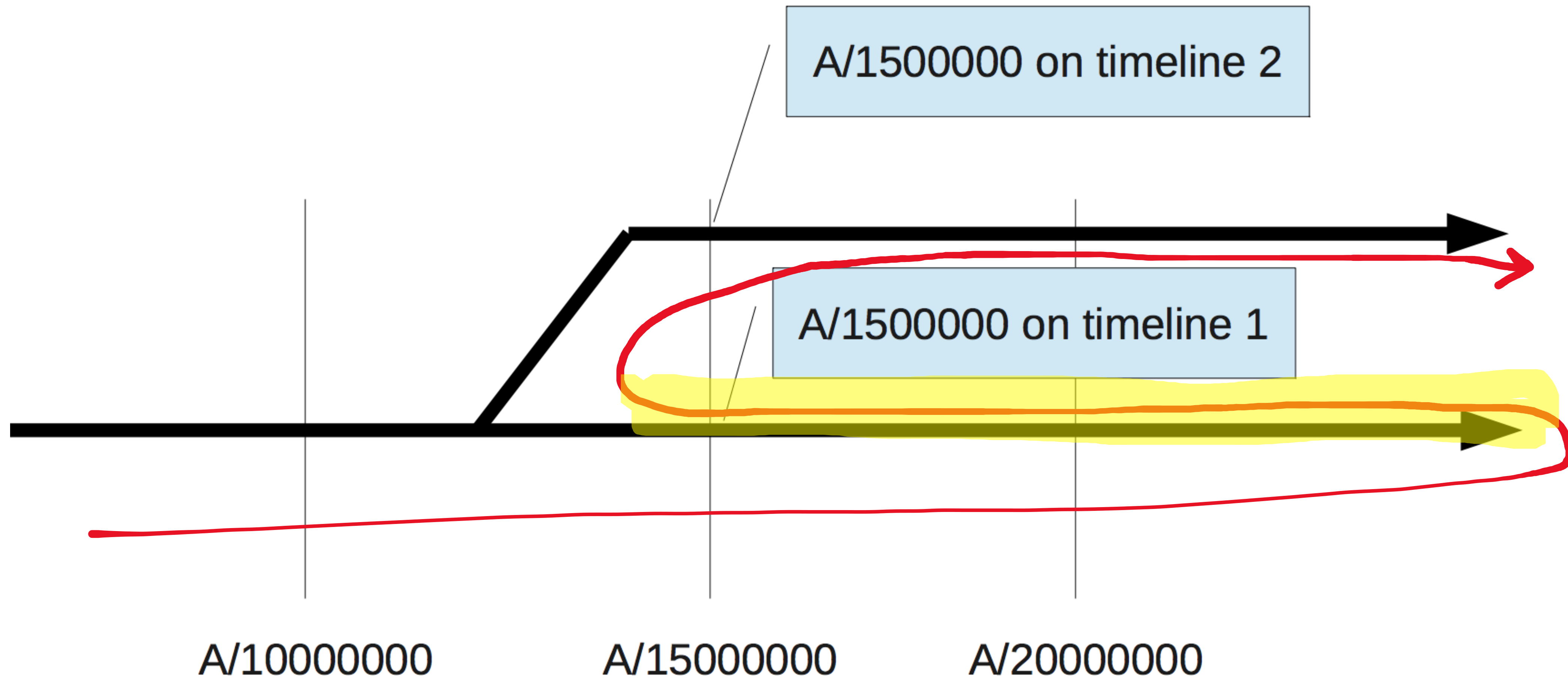
Synchronous standby names

Logical streaming can be ahead of

- › Synchronous standby
- › Quorum

Maybe add post_synchronous_standby_names?

Maybe logical pg_rewind?



So we have to hack the PostgreSQL

Currently we just implemented an extension to create slot in the past

- › We accept the risk of catalog vacuum after promotion
- › Anyway we need to stream data from PG 10,11,12,13,14 where things won't change much
- › But we are working to make it better

Some word about MySQL



**ARCHIVE_COMMAND
IS NOT SYNCHRONOUS**

**IF THERE IS
NOT ARCHIVE COMMAND**

No timelines: separate binlogs on each host



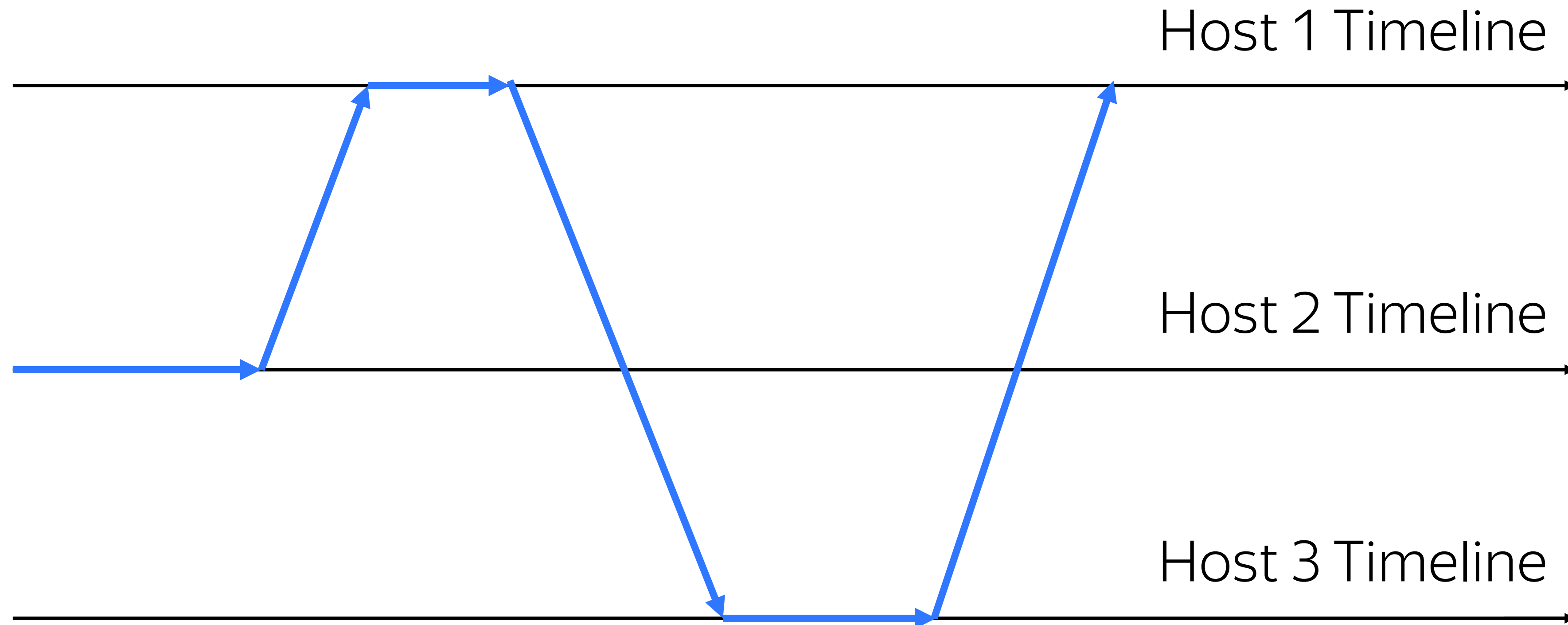
```
[mysql> show binary logs;
```

Log_name
mysql-bin-log-mst-134vfhk1sc1awst1r-d8-pandora-net.000256
mysql-bin-log-mst-134vfhk1sc1awst1r-d8-pandora-net.000257
mysql-bin-log-mst-134vfhk1sc1awst1r-d8-pandora-net.000258

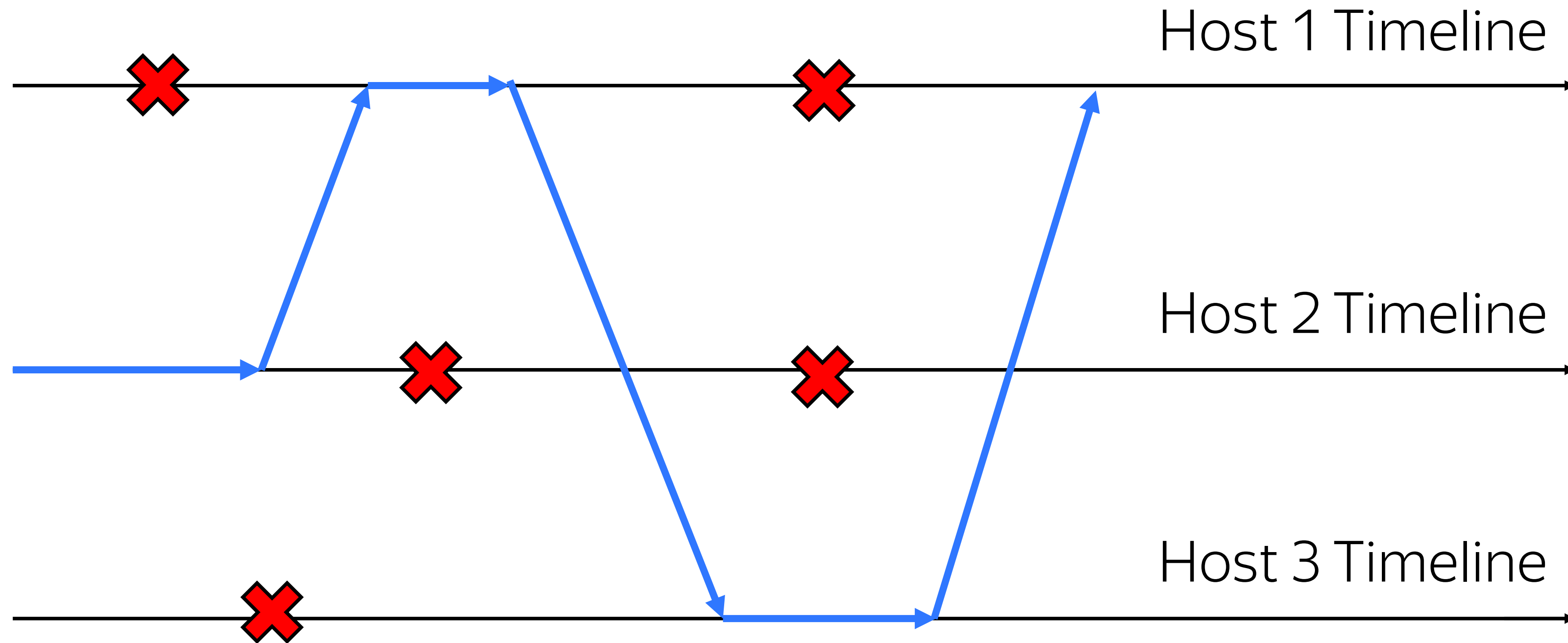
```
[mysql> show binary logs;
```

Log_name
mysql-bin-log-lve-ws1awst1r1k1sc1awst1r-d8-pandora-net.000249
mysql-bin-log-lve-ws1awst1r1k1sc1awst1r-d8-pandora-net.000250
mysql-bin-log-lve-ws1awst1r1k1sc1awst1r-d8-pandora-net.000251

Switchover/Failover



Nonlinear history == difficult PITR



MySQL logical replication

Some really nice concepts

- › GTID sets are beautiful
- › Automatic repositioning if fantastic

Tools



Disable cancellation of executed locally query

 ALTER SYSTEM SET synchronous_commit_cancelation to off;

<https://commitfest.postgresql.org/31/2402/>

lwaldump

Usage on primary: `CREATE EXTENSION lwaldump;`
Usage on standby: `SELECT lwaldump();`

<https://github.com/g0djan/lwaldump>

Create logical slot in the past

 `SELECT pg_create_logical_replication_slot(Isn(name, .., restart_Isn);`

https://github.com/x4m/pg_tm_aux

HArd to get it right

but not impossible



HArd to get it right

but not impossible

and corner cases are not that frequent actually



Waiting for questions 😊

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