Making MySQL-8.0 XA transaction processing crash safe

zettadb.com

David Zhao

About the author

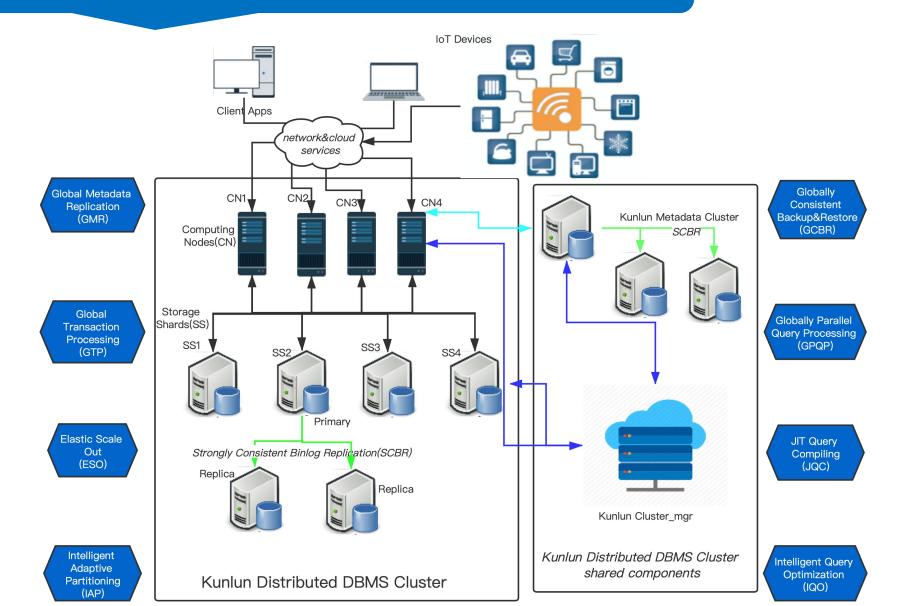
- Zhao Wei (David Zhao) twitter/linkedin/wechat: david.zhao.cn@gmail.com
- Database kernel developer in Oracle 2007-2015
 - Berkeley DB
 - > MySQL
- Database kernel developer in Tencent 2015-2019
 - > TDSQL --- most popular distributed RDBMS inside Tencent and Tencent Public/Private Cloud
 - > Lead the effort to evolve TDSQL from a table-sharding solution to a distributed RDBMS
- Left Tencent and started Kunlun project in Aug 2019
 - > To build an epochal distributed RDBMS based on widely adopted technologies(MySQL & PostgreSQL).
 - For more info about kunlun project visit www.zettadb.com
 - Kunlun project is open source at https://github.com/zettadb/kunlun



About Kunlun Distributed RDBMS

- Goals&objectives
 - A distributed RDBMS for PostgreSQL and MySQL users
 - with knowledge&lessons learned from work on TDSQL
 - compatible with PostgreSQL in DML grammar and most DDL grammar
 - compatible with MySQL in client protocol and DML grammar
 - > to manage multi-terabyte data for both OLTP and inplace OLAP workloads
 - Full-fledged NewSQL capabilities
 - highly available&fully consistent, crash safe&fault tolerant
 - highly scalable: table sharding & elastic horizontal scale-out
 - Cloud native & DBaaS
- Stands on the shoulders of giants
 - Kunlun computing node derived from PostgreSQL-11.5
 - Kunlun storage nodes derived from Percona-MySQL-8.0
 - ➤ keeps up with upper streams
 - ➢ Finished over 70% kernel development, released 0.7 in Sep. 2020

Kunlun Architecture



Kunlun Advantages & highlights

- Share nothing & adaptive sharding (IAP)
 - distribute tables/tablets intelligently to multiple storage shards
 - > multiple nodes for read & write queries, premium scalability
- Distributed transaction mgmt & distributed query processing
 - > works transparently, no burden for users, no client code changes needed
 - > appears like a standalone DB server for apps and app developers
- Crash safety&fault tolerance, high availability&strong consistency (GTP)
 - > any node/network fault at any time won't break transaction ACID guarantees
- Elastic horizontal scale-out (ESO) (*)
 - Automatic & elastic & continuous & on demand
 - unnoticible by application or end users
- Full-fledged distributed query processing (IQO) (on-going)
 - > cross shard join, subqueries, prepared statement, jit
 - OLAP analysis, stored procedures and more
- Global parallel query processing (GPQP)
 - Inside computing nodes(CN)
 - > CN send queries to target storage nodes(SN) asyncly so SS work in parallel
 - Inside a SN (*)

Kunlun Basics

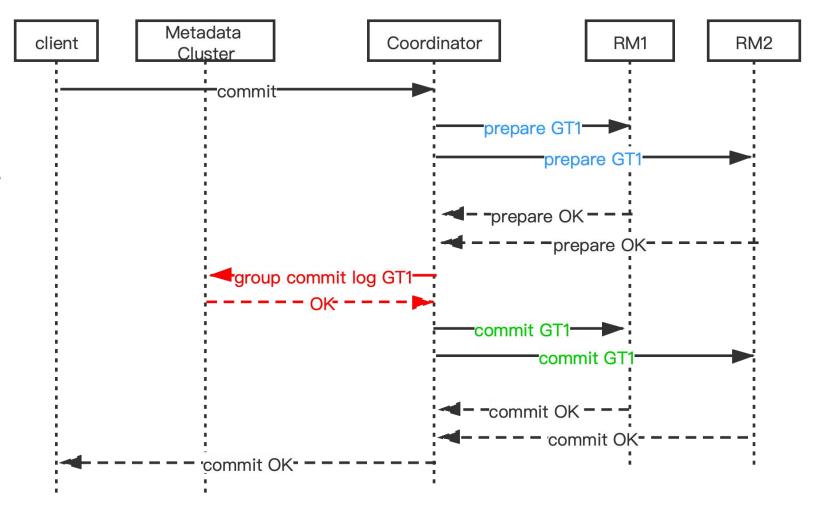
- Computing nodes
 - accept & validate user connections
 - accept & process user queries
 - parse -> optimize -> execute(send SQL -> receive & assemble)
 - \succ executes DDLs and DMLs
 - ➢ doesn't store user data, only store metadata locally
 - takes trivial storage/memory space
 - ➤ stores user data in storage shards
 - > can be built from metadata in metadata shard, equivalent to stateless proxies
 - needs no HA measures, no burden for DBAs
 - > add/remove nodes on demand, nodes independent from each other
- Support most PostgreSQL query processing features
 - Most DDLs and DMLs grammars
 - > indexes, views, materialized views, sequences; prepared statement, jit
 - cross shard join, subqueries, OLAP queries, stored procedures (on-going)
 - > All basic data types (numeric, string/text, date/time/timestamp, enum), json&spatial in future
- > Will support mysql client protocol and common MySQL private DML grammar(pending)

Kunlun Basics

- Storage shards
 - Stores application(user) data in standalone tables
 - PG single tables
 - PG table partitions
 - execute mostly single table queries
 - ➢ in a global transaction's local transaction branch
 - ➢ no table partitioning
 - Currently uses MySQL group replication(MGR) single primary mode for shard HA
 - ➤ primary election
 - robust consistency guarantees
 - ➢ Will support more types of HA (*) solutions
 - strongly consistent row based binlog replication
 - ➤ shared storage
 - Uses innodb only so far
 - never untransactional engines(myisam, etc)
 - > maybe myrocks in future
 - Require kunlun-percona-MySQL-8.0.18-9
 - developed based on percona-MySQL-8.0.18-9
 - > contains critical bug fixes & supporting features & performance enhancements
 - > will keep up with upper stream

Kunlun Distributed Transaction Processing

- Global Transaction Coordinator
 - component of a computing node
 - Tracks txn branch states
 - 2PC for multi-shard writers
 - IPC for single shard writers
 - IPC for read-only txn branches
 - before 2nd phase, group log commit decisions in commit-log
 - Abort txn on error during txn commit



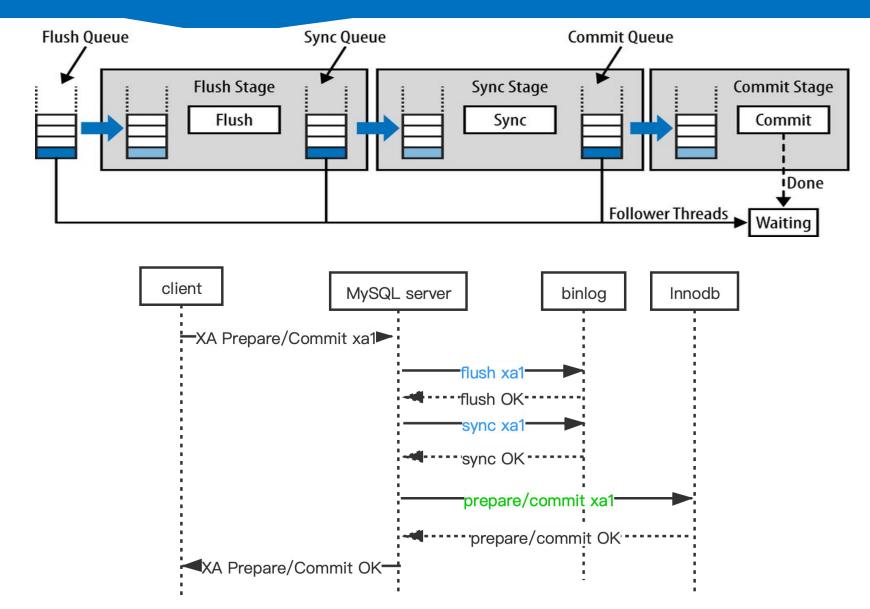
Potential Failures

- Node failures
 - Computing nodes (CNs)
 - Primary nodes of storage shards
 - Replica nodes of storage shards
 - Primary nodes of metadata shards
- Network failures between
 - CNs & Meta shard primary
 - CNs & Storage shard primary
 - > A primary MySQL node and its replicas
- ➢ Fail while many global txns are in 2PC commit phases
 - > partial commits
- Fail during binlog group commit
 - many regular txns doing internal 2PC (prepare->commit)
 - > many XA txns doing prepare
 - > many prepared XA txns doing commit

Requirement for crash-safety & consistency in MySQL

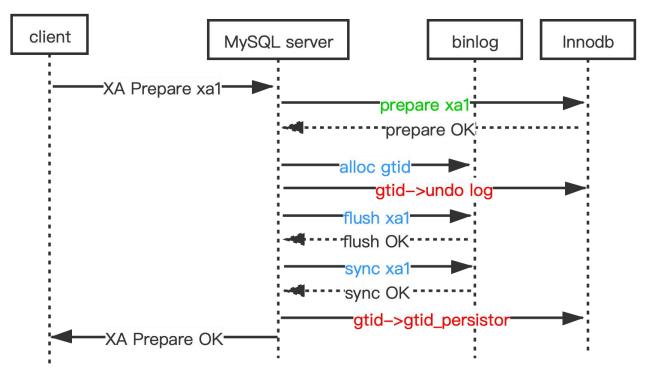
- Keep consistency between
 - \succ a primary and its replicas
 - ➤ replica local sync
 - ➤ same binlog event groups
 - ➤ innodb and binlog
 - ➢ for any txn T, T in innodb <=> T in binlog
 - ➢ key: no extra event groups in binlog
 - ➤ challenge: binlog is passive and inflexible
 - gtid set in innodb undo log, binlog and mysql.gtid_executed (MGE)
 - New requirement for MySQL-8.0 for clone
 - txn.gtid -> undo log -> mysql.gtid_executed (MGE)

MySQL Binlog Group Commit (BGC)



- > XA Prepare handling in community mysql
 - step1: binlog prepare GT
 - binlog group commit flush&sync phases
 - > wait for its reception by replicas before flush(MGR) or after sync (semisync)
 - step 2: engine(innodb, myrocks, etc) prepare GT
- Issue 1: primary may crash between steps #1 and #2
 - > Some replicas may have replicated binlogs of GT and replayed it, so GT is prepared in replica(s)
 - Crashed primary node has inconsistent binlogs&innodb
 - > after primary switch, old primary(now replica) can't do 'XA COMMIT GT', replication blocks
- > Solution: do binlog prepare after all engine prepare, so the new procedure is:
 - step1: engine(innodb, myrocks, etc) prepare GT
 - step2: binlog prepare GT
 - \succ the solution causes issue 2

- > XA Prepare Issue 2: dilemma caused by MySQL-8.0 clone
 - > no gtid at step #1, only has gtid at beginning of step #2
- Solution: interleave operations
 - engine prepare -> generate gtid -> write gtid to undo log -> flush binlog
 - innodb recovery: abort prepared txns without gtid in its undo log
 - gtid -> gtid_persistor -> MGE after binlog sync



- Binlog Recovery
 - ➢ engine recovery first
 - scans last binlog file to form internal XA txn ids S
 - > why only last binlog ?
 - commit/abort each recovered prepared innodb txn if it is/isn't in S
 - Issue 3: external XA txns are ignored in binlog recovery
 - Solution: for each prepared external XA txn xi in every storage engine
 - > commit xi if 'XA COMMIT xi' is found in last binlog file
 - > otherwise abort xi if 'XA ROLLBACK xi' is found in last binlog file
 - > otherwise leave xi prepared if 'XA PREAPRE xi' is found in S
 - > otherwise abort xi if 'XA PREPARE xi' isn't found in S
 - ➤ special handling for 'XA COMMIT ONE PHASE'

- ➤ Where are prepared txns?
 - > prepared XA txns can live across binlog file rotations
 - Issue 4: not sufficient to scan last binlog file for prepared external XA txns
 - > XA PREPARE & XA COMMIT seperated apart
- Solution: record all current prepared external XA txns at creation of each new binlog file
 - kunlun-percona: pack into Previous_gtids event
 - ➢ loss: ecosystem compatibility
 - ➤ gains: distributed txn crash safety

mysql> xa start 'xa2'; Query OK, 0 rows affected (0.00 sec)	<pre>mysql> show binlog events in 'binlog.000011';</pre>						
	Log_name	Pos	; Event_type	Server_	id End_log_p	oos	
Database changed mysql> insert into t1 values(7,8); Query OK, 1 row affected (0.01 sec) mysql> xa end 'xa2'; Duery OK, 0 rows affected (0.01 sec)	binlog.000011 binlog.000011	4 124	Format_desc Previous_gtid	295 s 295	11 1 11 2	124 223	Server ver: 8.0.18-9-kunlun-percona-mysql-debug, Binlog ver: 4 51078c3a-547e-11ea-9780-981fd1bd410d:1-9252:1000608-1000667 XA-PREPARED: xa1
	++ 2 rows in set (0.00 sec)						
	mysql> show binlog events in 'binlog.000012';						
			Event_type				٥
	binlog.000012 binlog.000012	4 124	Format_desc Previous_gtids	29511 29511	124 227	Ser 510	ver ver: 8.0.18-9-kunlun-percona-mysql-debug, Binlog ver: 4 78c3a-547e-11ea-9780-981fd1bd410d:1-9253:1000608-1000667 XA-PREPARED: xa1 xa2
	2 rows in set (0.00 sec)						
	mysql> flush logs; Query OK, 0 rows affected (0.04 sec)						
	mysql> show binlog events in 'binlog.000013'; +						
	+	+	Event_type	+	+	+	
			Format_desc Previous_gtids				ver ver: 8.0.18-9-kunlun-percona-mysql-debug, Binlog ver: 4 78c3a-547e-11ea-9780-981fd1bd410d:1-9254:1000608-1000667 XA-PREPARED: xa3 xa1 xa2
	t						
mysql> xa start 'xa3';							
Query OK, 0 rows affected (0.01 sec)							
<pre>mysql> insert into t1 values(9,10);</pre>							
Query OK, 1 row affected (0.00 sec)							
mysql> xa end 'xa3';							
Query OK, 0 rows affected (0.00 sec)							
mysql> xa prepare 'xa3';							
Query OK, 0 rows affected (0.00 sec)							

Thank you!