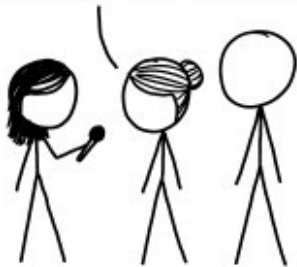


Threat modelling for developers

Arne Padmos

ASKING AIRCRAFT DESIGNERS
ABOUT AIRPLANE SAFETY:

NOTHING IS EVER FOOLPROOF,
BUT MODERN AIRLINERS ARE
INCREDIBLY RESILIENT. FLYING IS
THE SAFEST WAY TO TRAVEL.



ASKING BUILDING ENGINEERS
ABOUT ELEVATOR SAFETY:

ELEVATORS ARE PROTECTED BY
MULTIPLE TRIED-AND-TESTED
FAILSAFE MECHANISMS. THEY'RE
NEARLY INCAPABLE OF FALLING.



ASKING SOFTWARE
ENGINEERS ABOUT
COMPUTERIZED VOTING:

THAT'S TERRIFYING.



WAIT, REALLY?

DON'T TRUST VOTING SOFTWARE AND DON'T
LISTEN TO ANYONE WHO TELLS YOU IT'S SAFE.

WHY?

I DON'T QUITE KNOW HOW TO PUT THIS, BUT
OUR ENTIRE FIELD IS BAD AT WHAT WE DO,
AND IF YOU RELY ON US, EVERYONE WILL DIE.



THEY SAY THEY'VE FIXED IT WITH
SOMETHING CALLED "BLOCKCHAIN."

AAAAA!!!

WHATEVER THEY SOLD
YOU, DON'T TOUCH IT.
BURY IT IN THE DESERT.
WEAR GLOVES.



Safety

vs

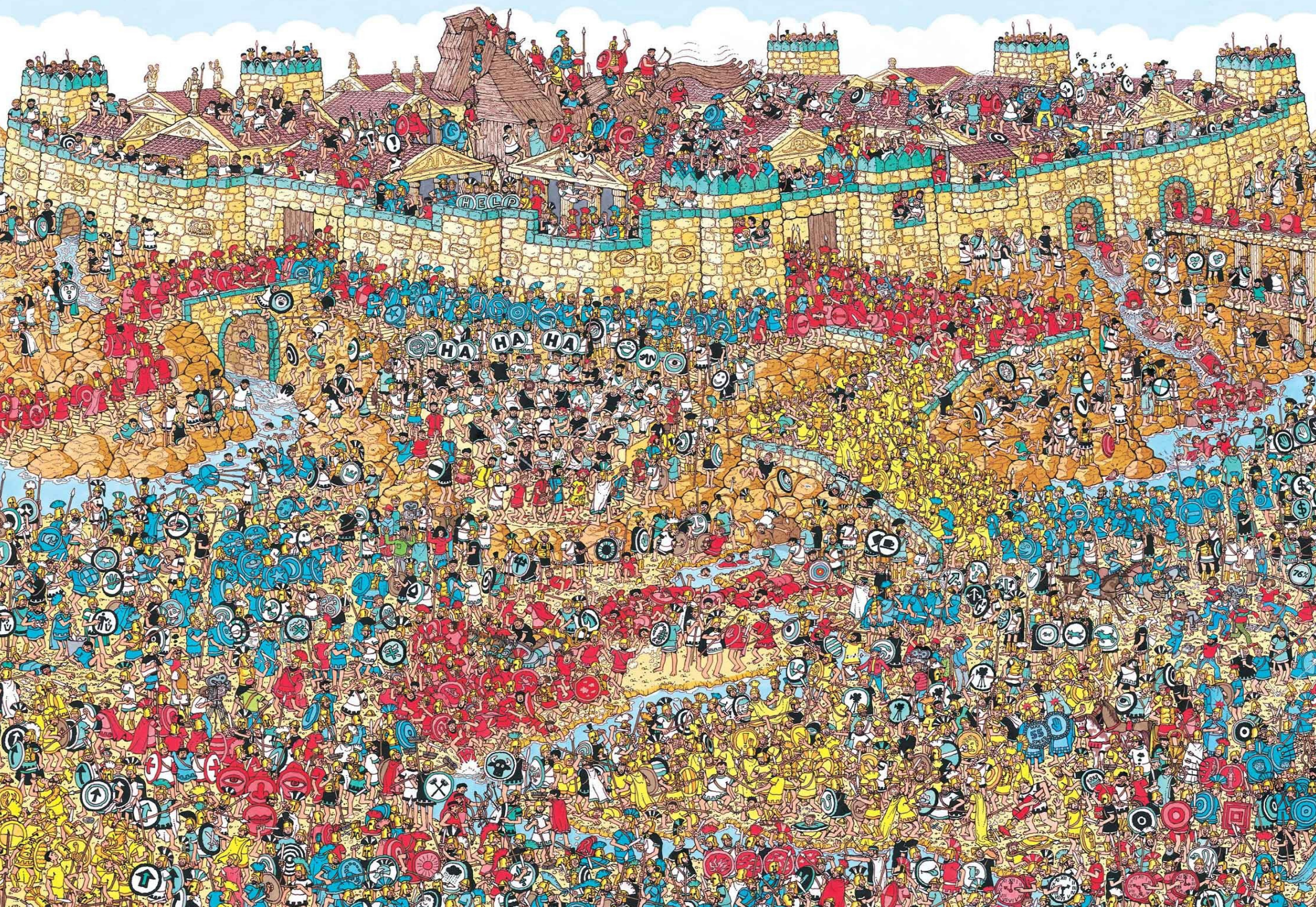
Security







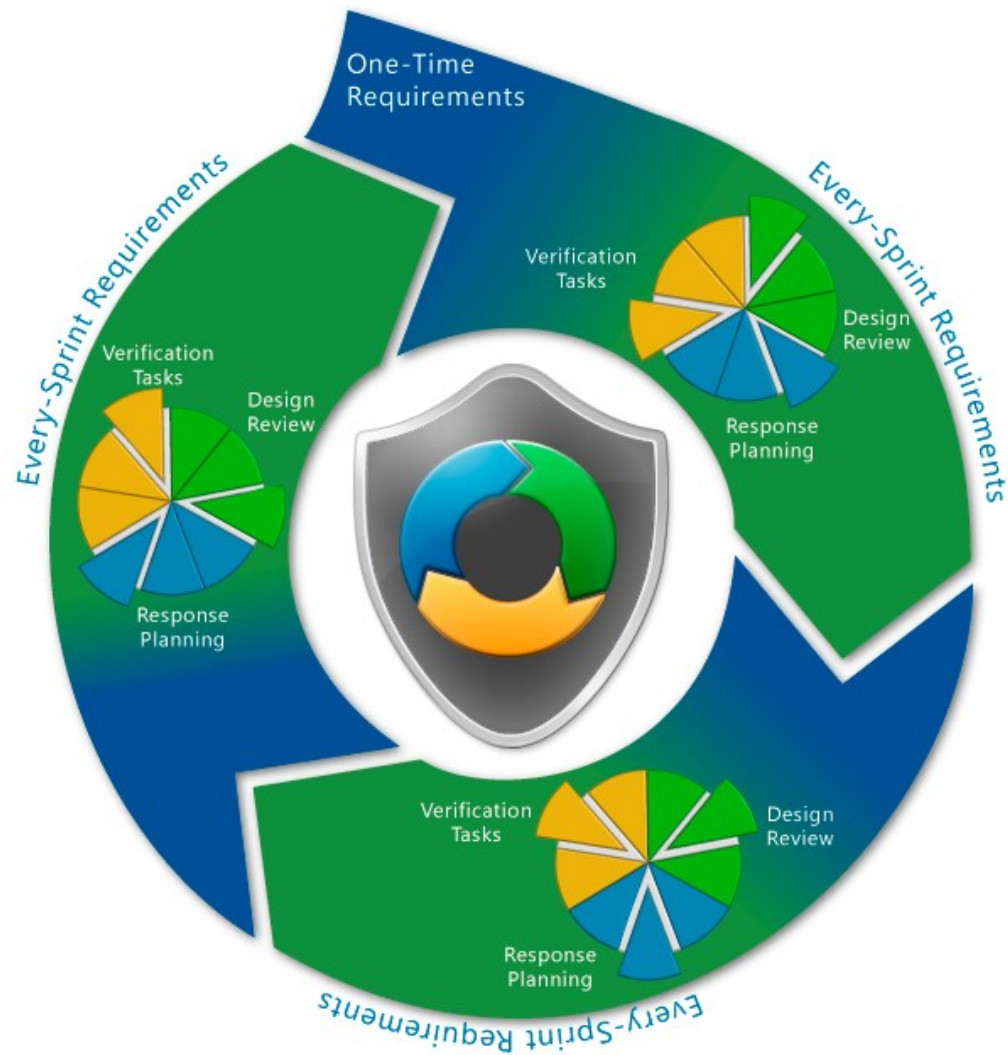
Are we doomed?

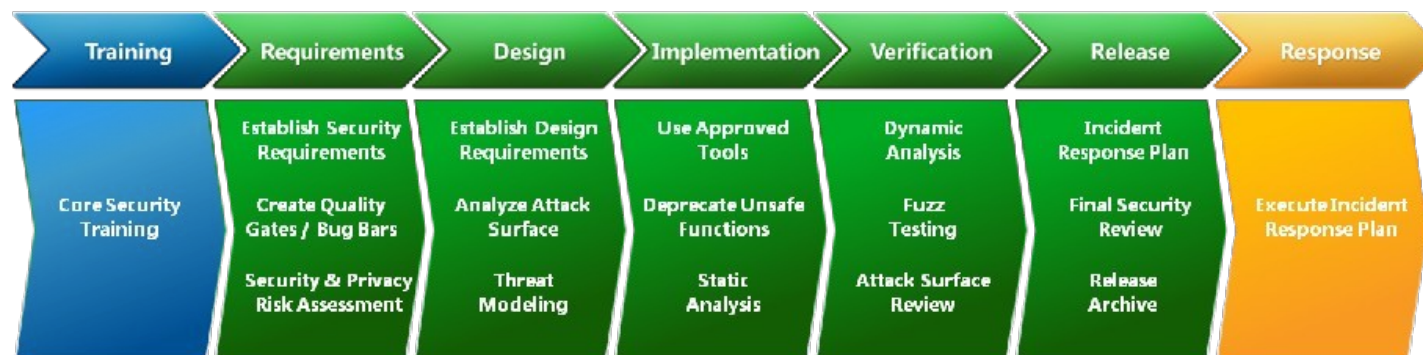
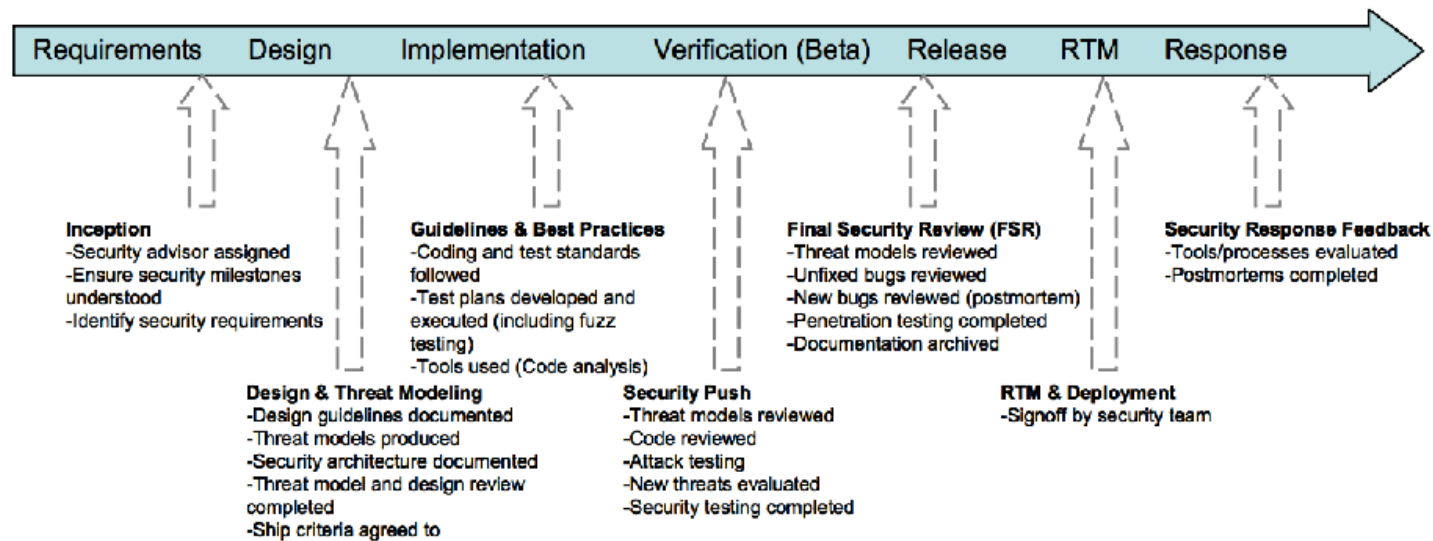


“ Building security in ”

“ Security by design ”

“ Shifting security left ”





*“If we ... could do only one thing
to improve software security ...
we would do threat modelling
every day of the week.”*

— Howard & Lipner

*“If we ... could do only one thing
to improve software security ...
we would do **threat modelling**
every day of the week.”*

— Howard & Lipner

Requirements engineering & Architectural analysis

What's your threat model?
(security assumptions)

On the Security of Public Key Protocols

DANNY DOLEV AND ANDREW C. YAO, MEMBER, IEEE

Abstract—Recently the use of public key encryption to provide secure network communication has received considerable attention. Such public key systems are usually effective against passive eavesdroppers, who merely tap the lines and try to decipher the message. It has been pointed out, however, that an improperly designed protocol could be vulnerable to an active saboteur, one who may impersonate another user or alter the message being transmitted. Several models are formulated in which the security of protocols can be discussed precisely. Algorithms and characterizations that can be used to determine protocol security in these models are given.

I. INTRODUCTION

THE USE of public key encryption [1], [11] to provide secure network communication has received considerable attention [2], [7], [8], [10]. Such public key systems are usually very effective against a “passive” eavesdropper, namely, one who merely taps the communication line and tries to decipher the intercepted message. However, as pointed out in Needham and Schroeder [8], an improperly designed protocol could be vulnerable to an “active” saboteur, one who may impersonate another user and may alter or replay the message. As a protocol might be compromised in a complex way, informal arguments that assert the security for a protocol are prone to errors. It is thus desirable to have a formal model in which the security

issues can be discussed precisely. The models we introduce will enable us to study the security problem for families of protocols, with very few assumptions on the behavior of the saboteur.

We briefly recall the essence of public key encryption (see [1], [11] for more information). In a public key system, every user X has an *encryption function* E_x and a *decryption function* D_x , both are mappings from $\{0, 1\}^*$ (the set of all finite binary sequences) into $\{0, 1\}^*$. A secure public directory contains all the (X, E_x) pairs, while the decryption function D_x is known only to user X . The main requirements on E_x, D_x are:

- 1) $E_x D_x = D_x E_x = 1$, and
- 2) knowing $E_x(M)$ and the public directory does not reveal anything about the value M .

Thus everyone can send X a message $E_x(M)$, X will be able to decode it by forming $D_x(E_x(M)) = M$, but nobody other than X will be able to find M even if $E_x(M)$ is available to them.

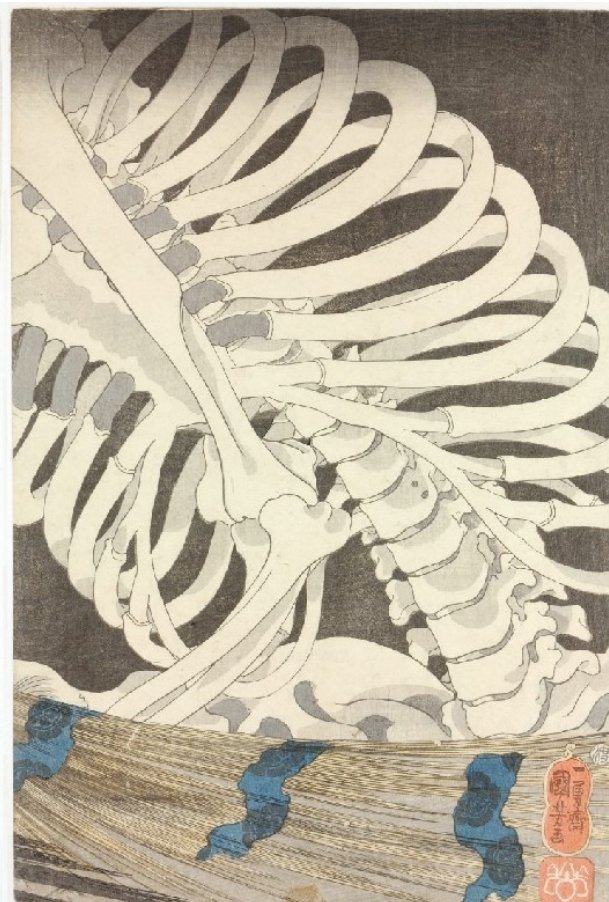
We will be interested mainly in protocols for transmitting a secret plaintext M between two users. To give an idea of the way a saboteur may break a system, we consider a few examples. A message sent between parties in the network consists of three fields: the sender's name, the receiver's name, and the text. The text is the encrypted part of the message. We will write a message in the format: sender's name, text, receiver's name.

Example 1: Consider the following protocol for sending

Manuscript received July 15, 1981; revised August 8, 1982. This work was supported in part by ARPA under Grant MDA-903-80-C-102 and by National Science Foundation under Grant MCS-77-05313-A01. This paper was partially presented at the 22nd Annual IEEE Symposium on Foundations of Computer Science, Nashville, TN, October 28-30, 1981.

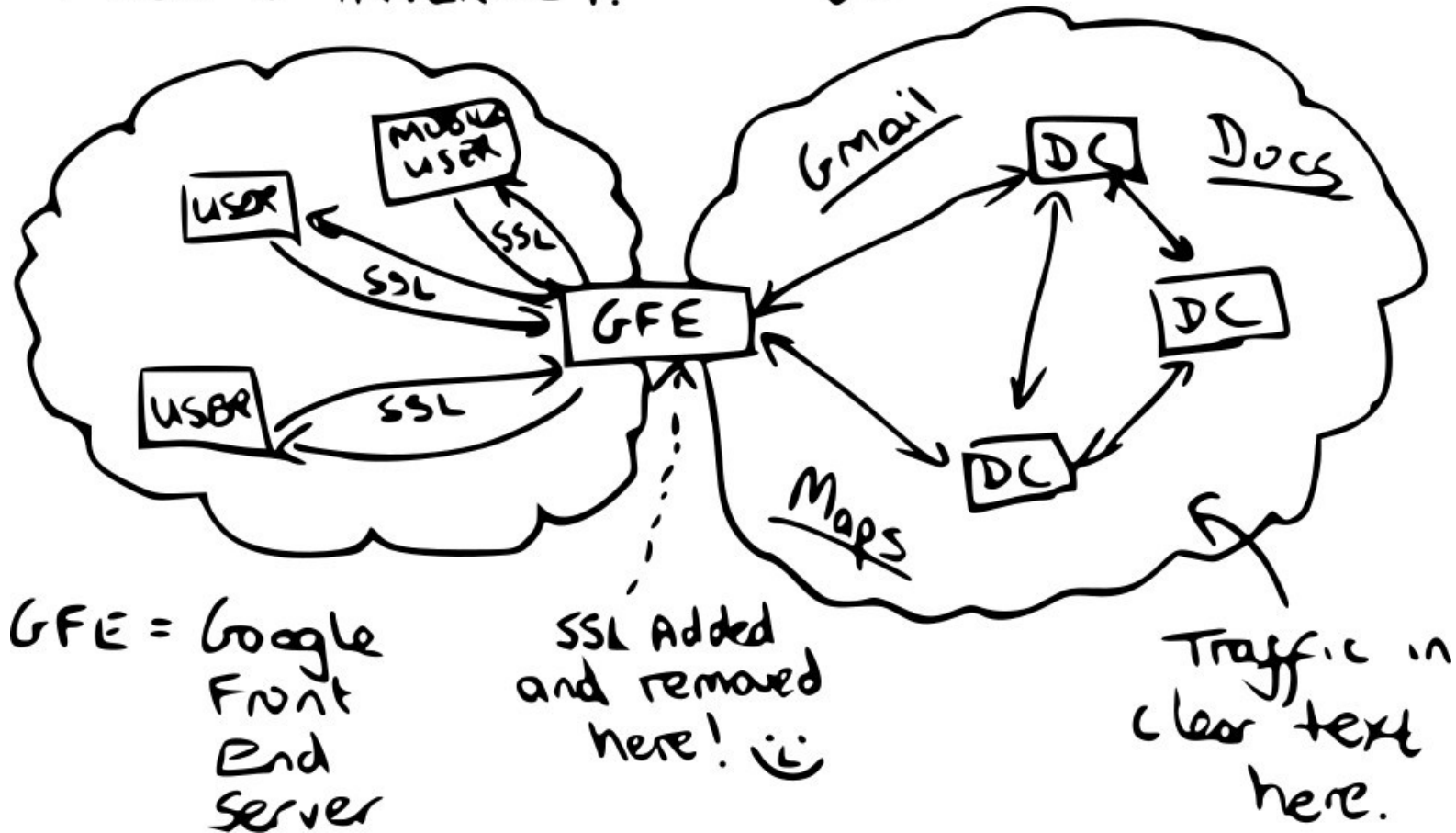
“ More precisely, we will assume the following about a saboteur: ”

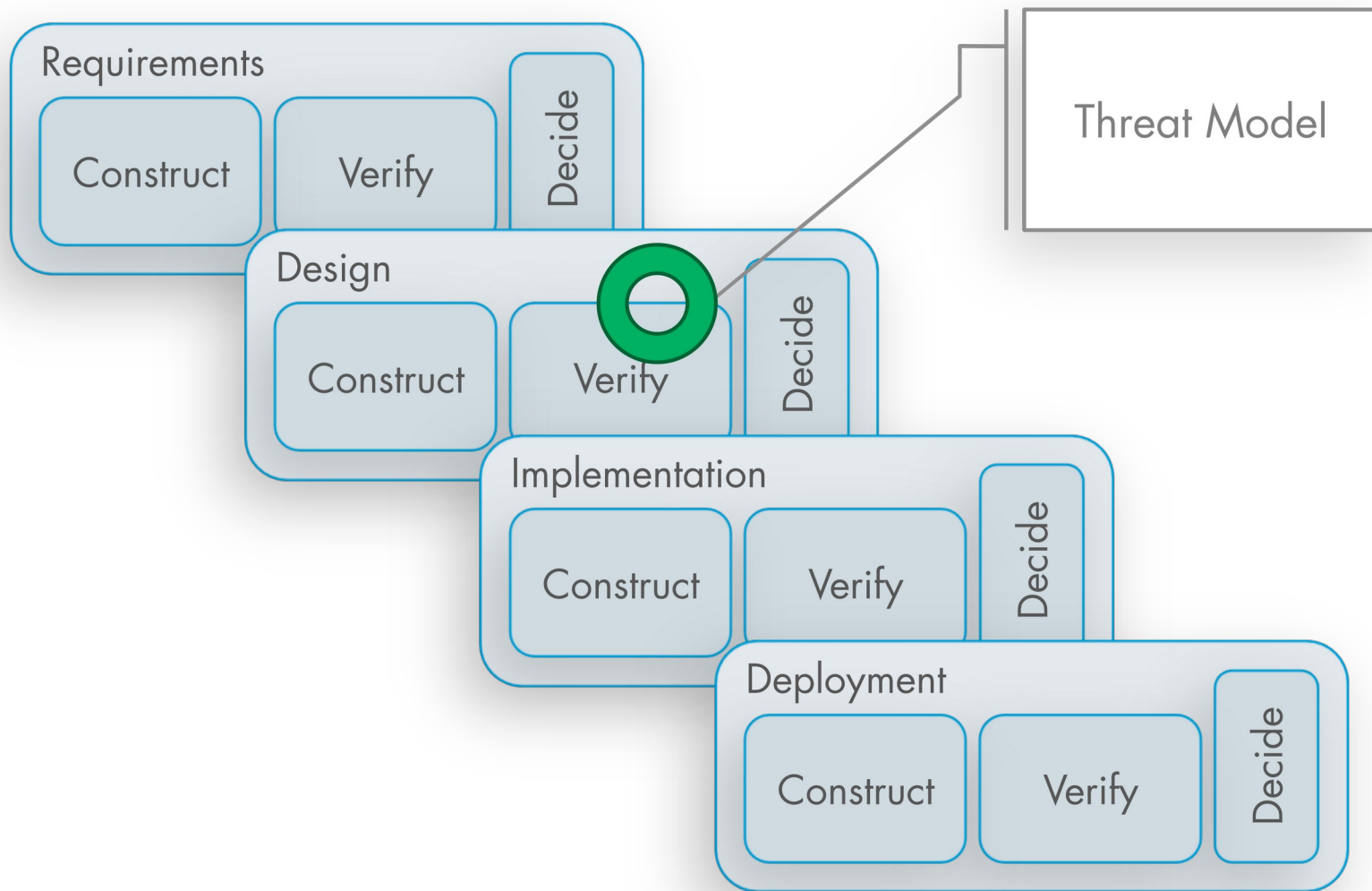
- obtain any message
- initiate any conversation
- be a receiver to any user



PUBLIC INTERNET.

GOOGLE CLOUD





**What could
possibly
go wrong?**

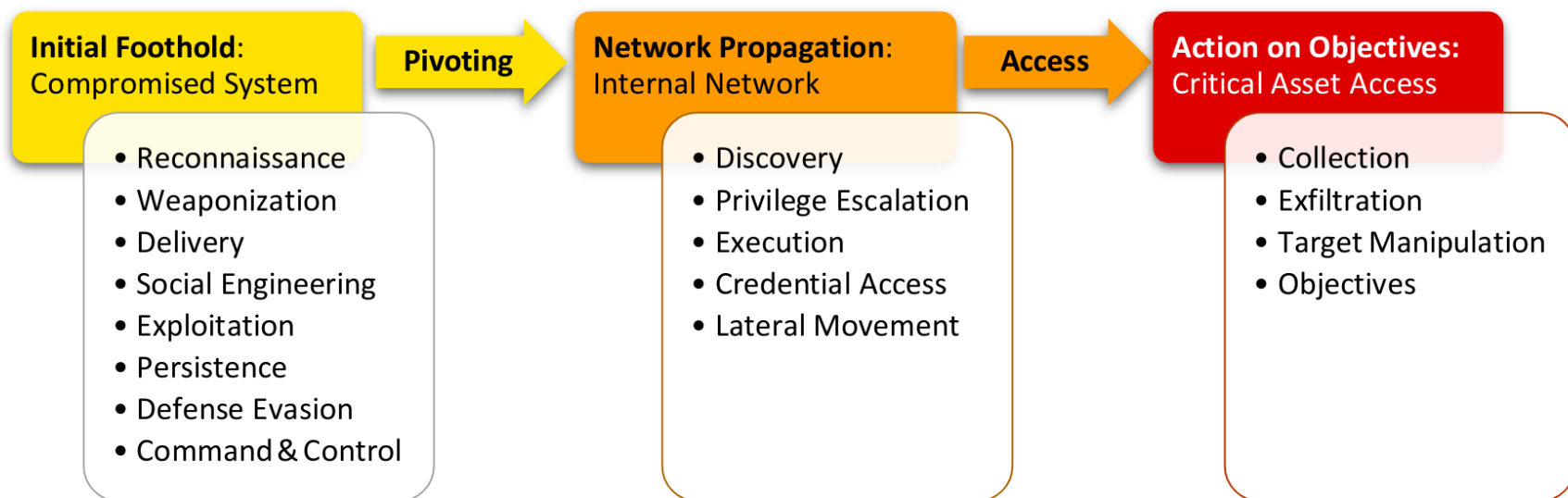
**What could
possibly
go wrong?**

& how

Types of threat modelling

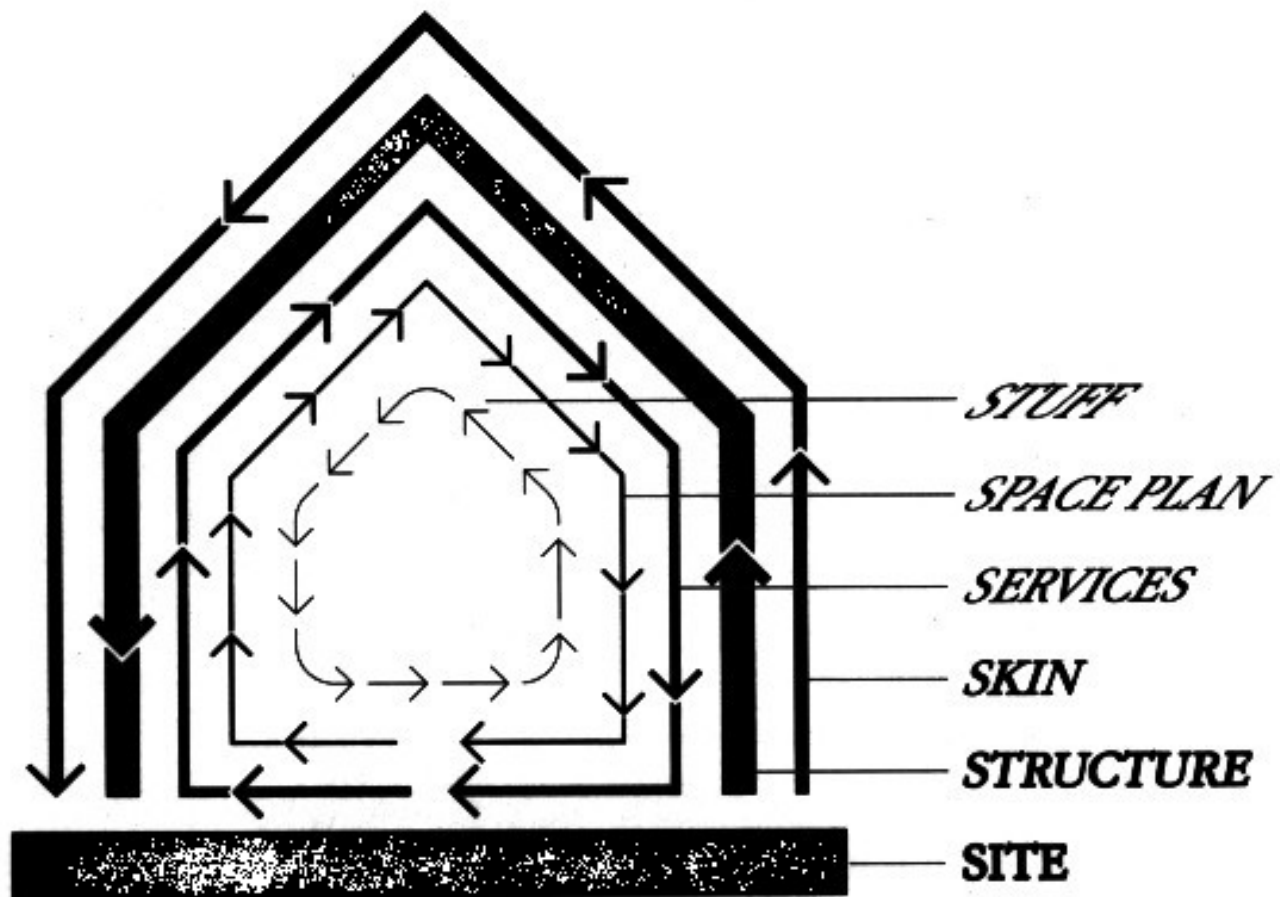
- Attacker-centric
- Asset-centric
- System-centric

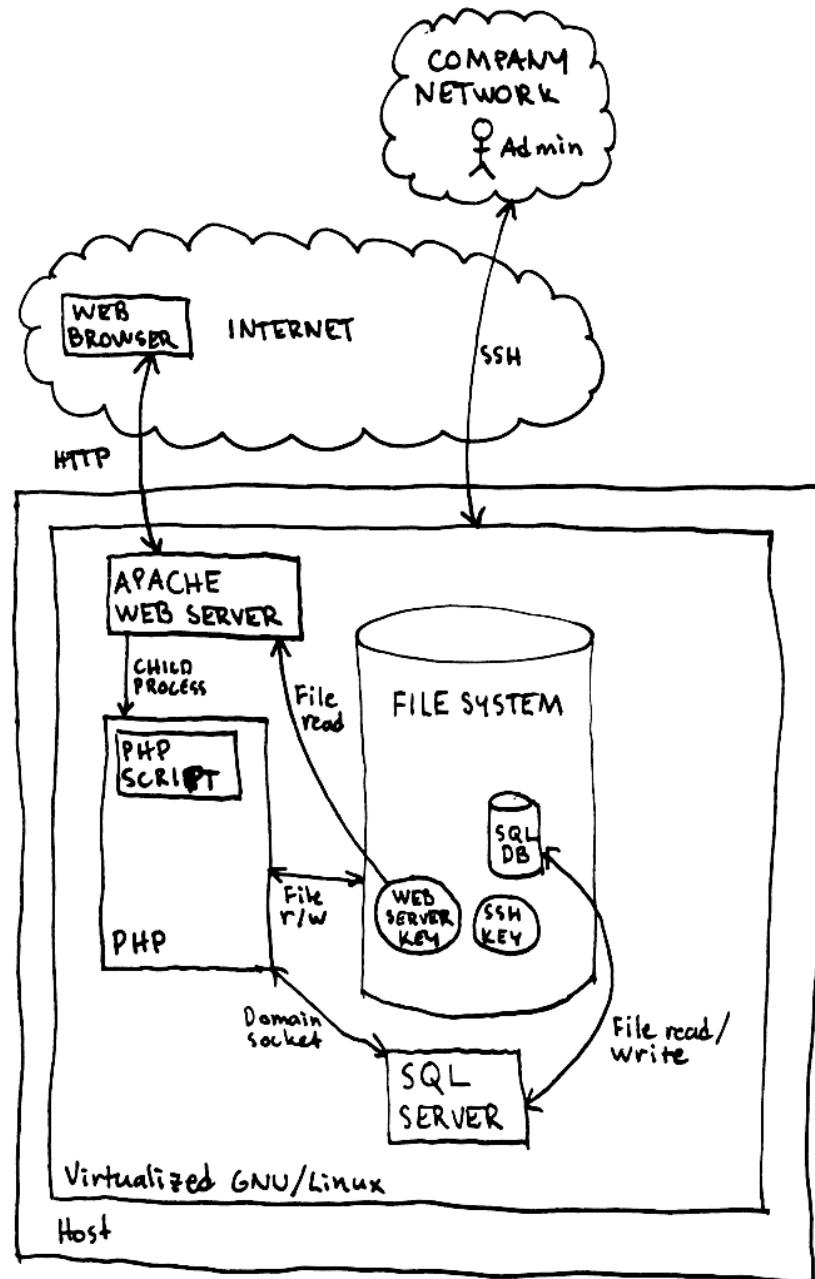






| | | | | | | | | | | | | | | |
|----------|---|---|--------|--|--|--------|--|--|--------|--|--|--|--|--|
| | This action does not apply to this asset, based on the asset's type in the Data Model tab. | | | | | | | | | | | | | |
| | (Never) The system should never let this actor take this action on this asset. | | | | | | | | | | | | | |
| | (Conditionally) The system should let this actor take this action on this asset when certain conditions (typically documented in the cell comment) are met. | | | | | | | | | | | | | |
| | (Always) The system should always let this actor take this action on this asset. | | | | | | | | | | | | | |
| Example: | | | | | | | | | | | | | | |
| C | R | X | Actor | | | | | | | | | | | |
| U | D | F | Author | | | Editor | | | Reader | | | | | |
| Asset | Blog | | | | | | | | | | | | | |
| | | | | | | | | | | | | | | |
| | Blog Post | | | | | | | | | | | | | |
| | | | | | | | | | | | | | | |





Popular approaches (*system-centric*)

- STRIDE
- Trike
- PASTA

Relevant questions

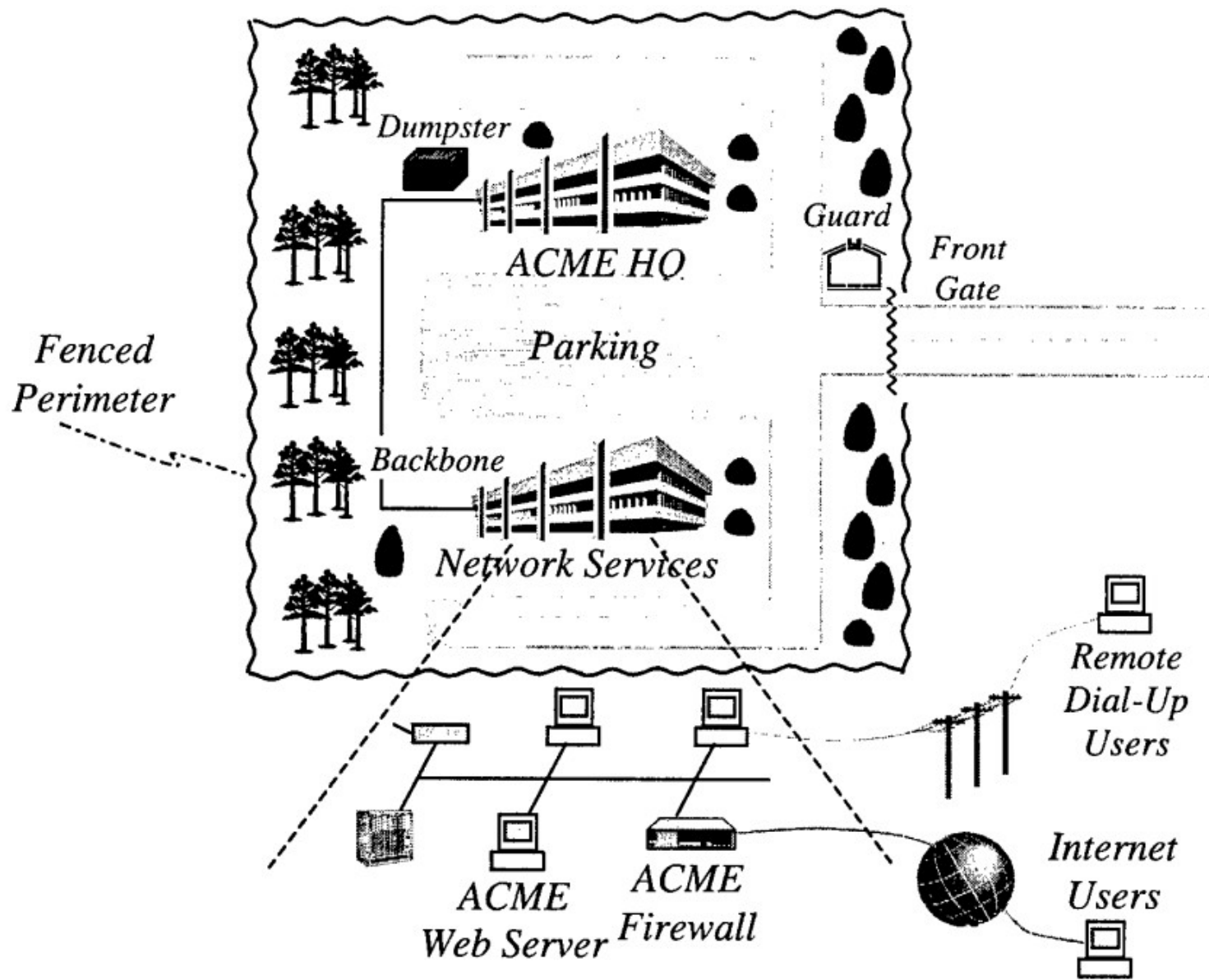
1. What are we working on?
2. What can go wrong?
3. What are we going to do?
4. Did we do a good job?

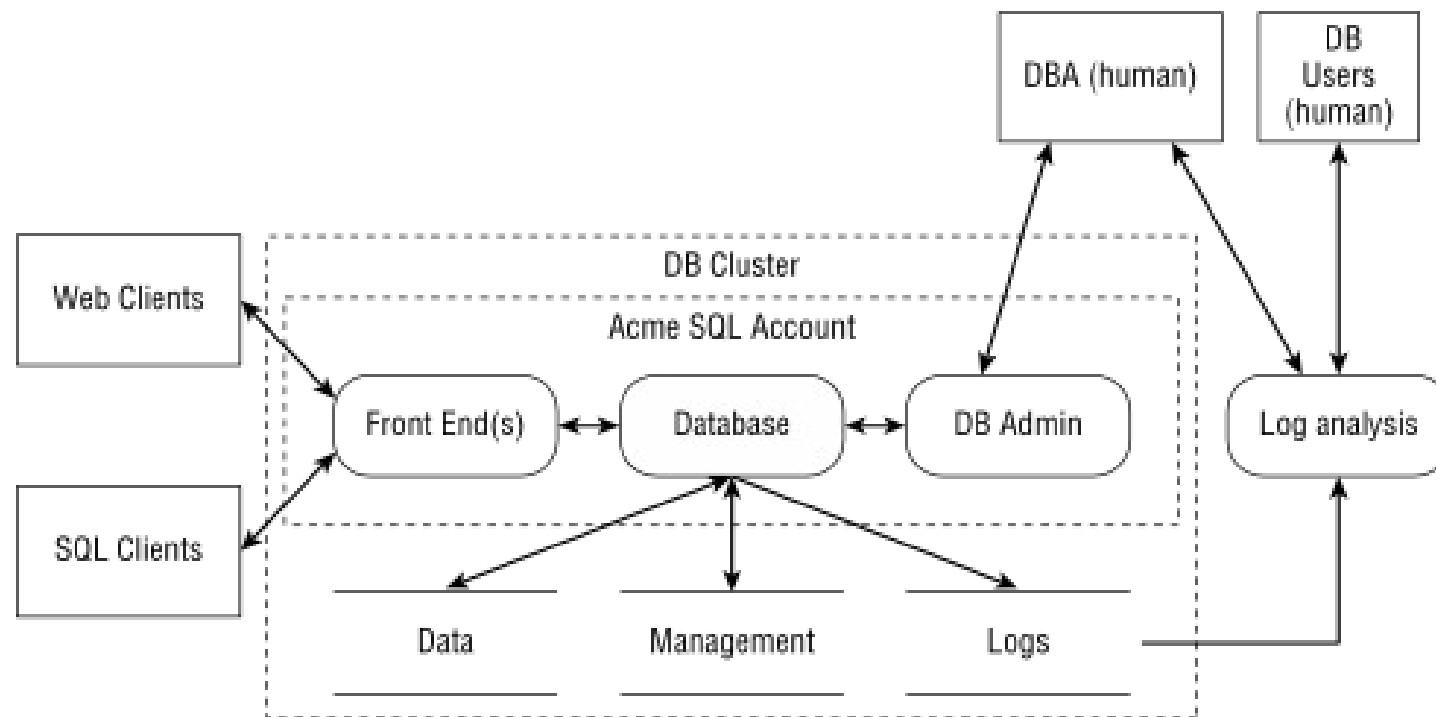
Lightweight methodology

1. Draw data flows
2. Elicit threats
3. Ranking + controls
4. Check your work

Lightweight methodology

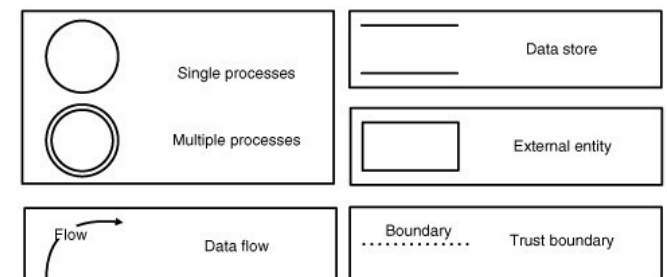
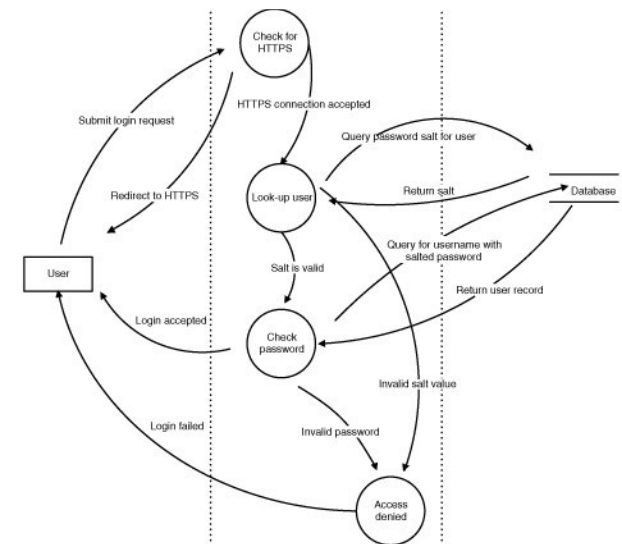
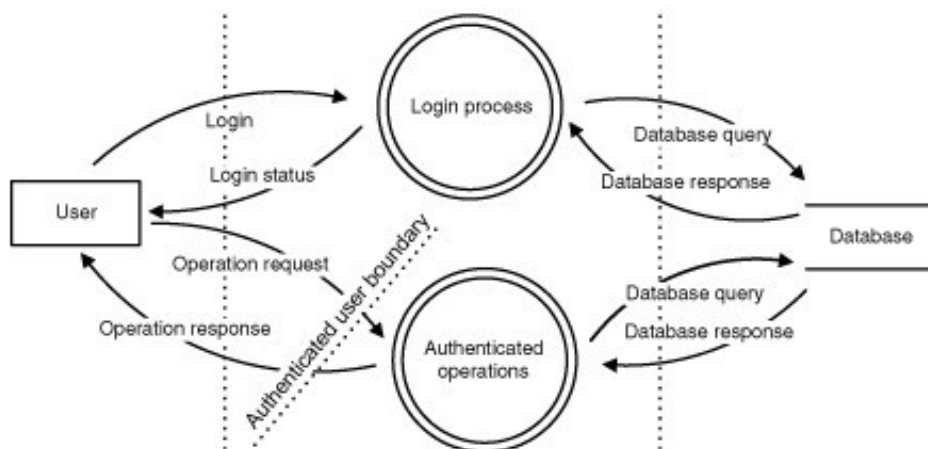
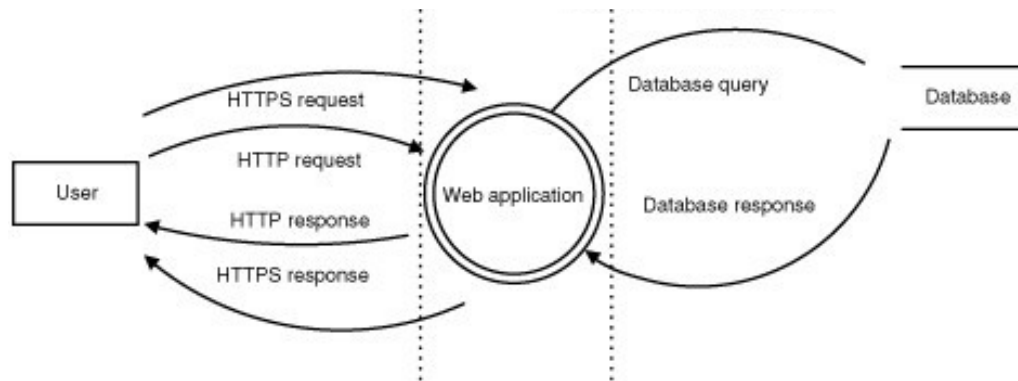
- 1. Draw data flows**
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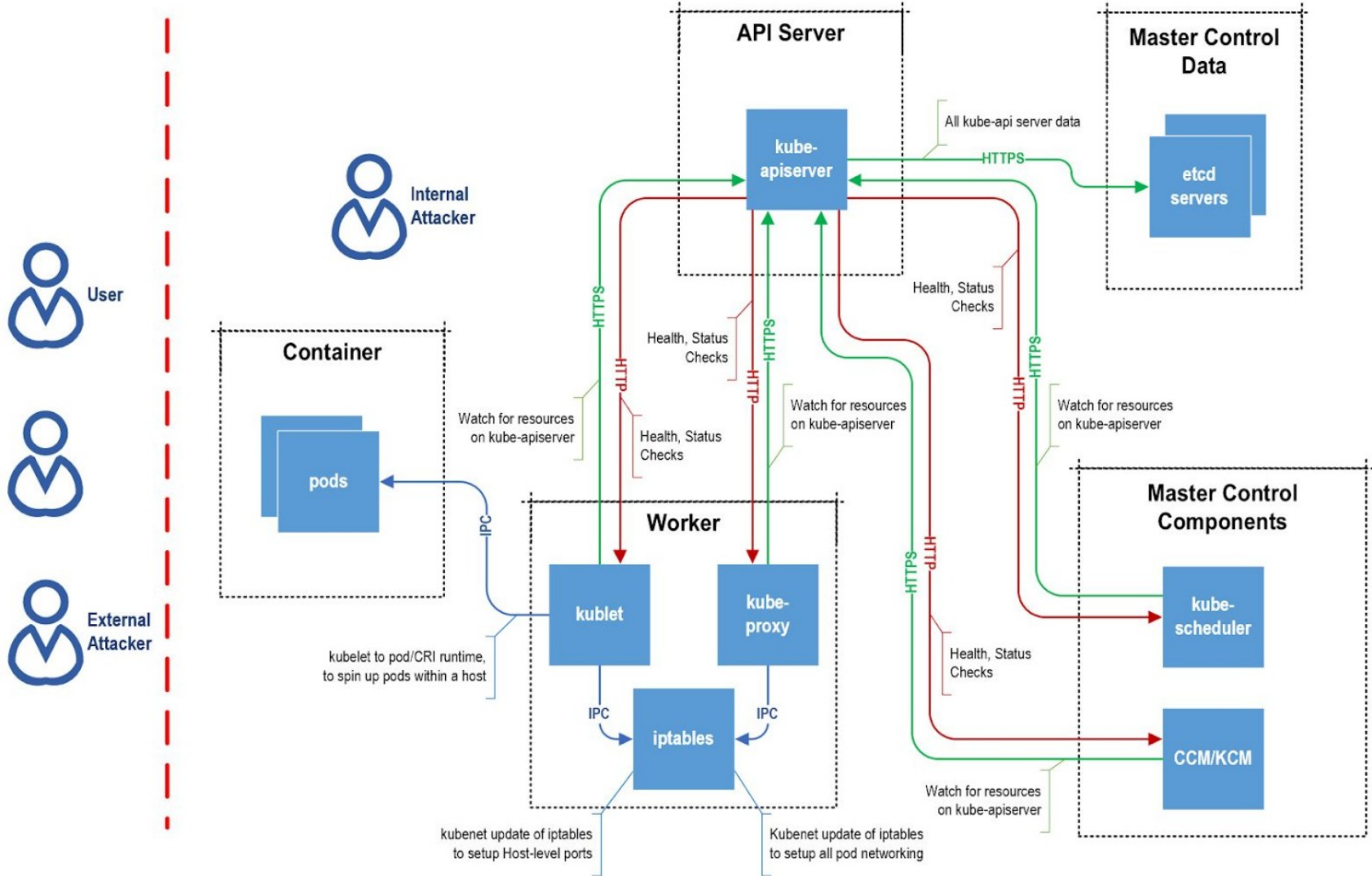


Key:





Internet



Lightweight methodology

1. Draw data flows
- 2. Elicit threats**
3. Ranking + controls
4. Check your work

Confidentiality

Integrity

Availability

Authentication

Authorisation

Accountability

Information disclosure

Tampering

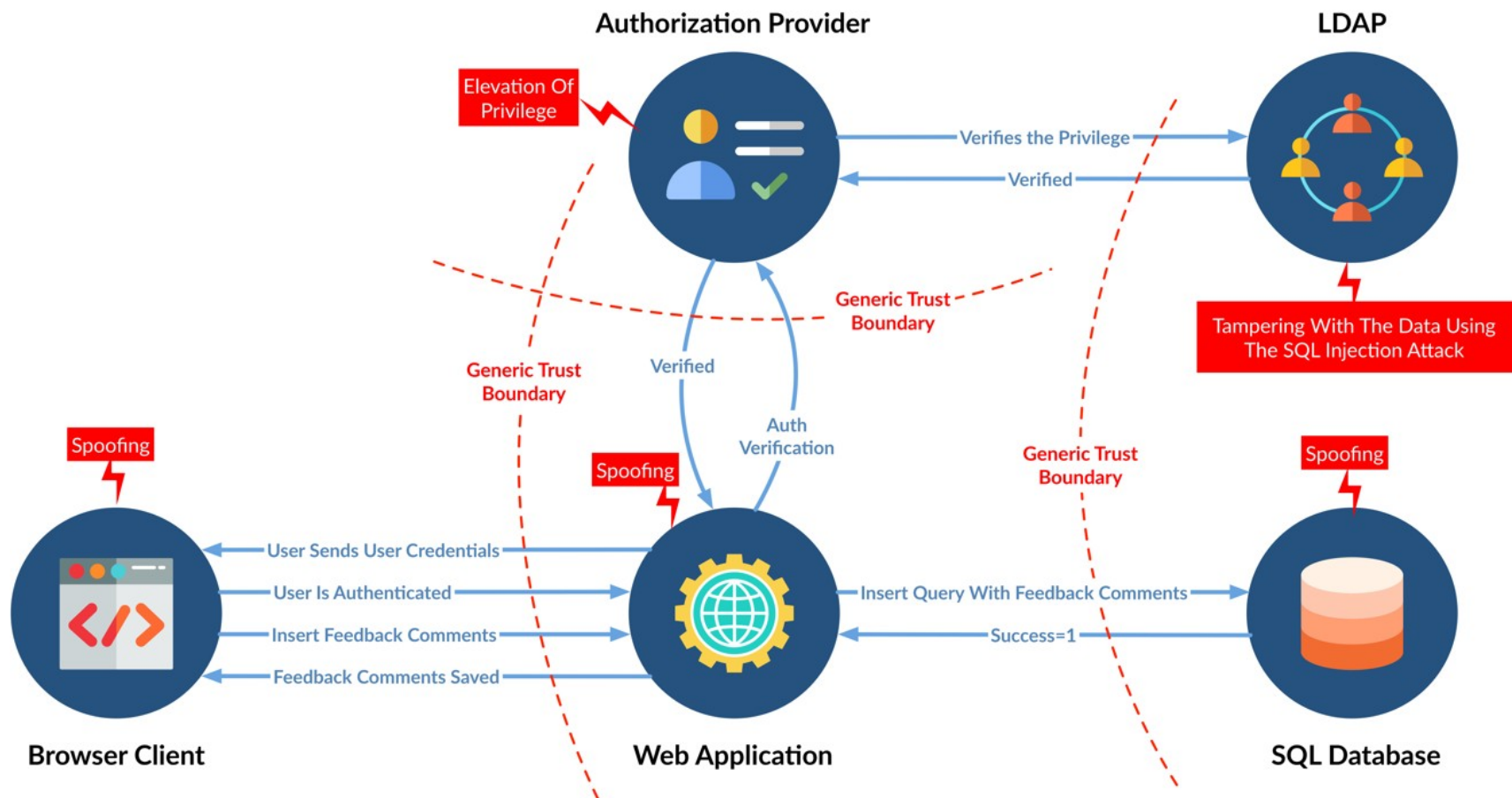
Denial of service

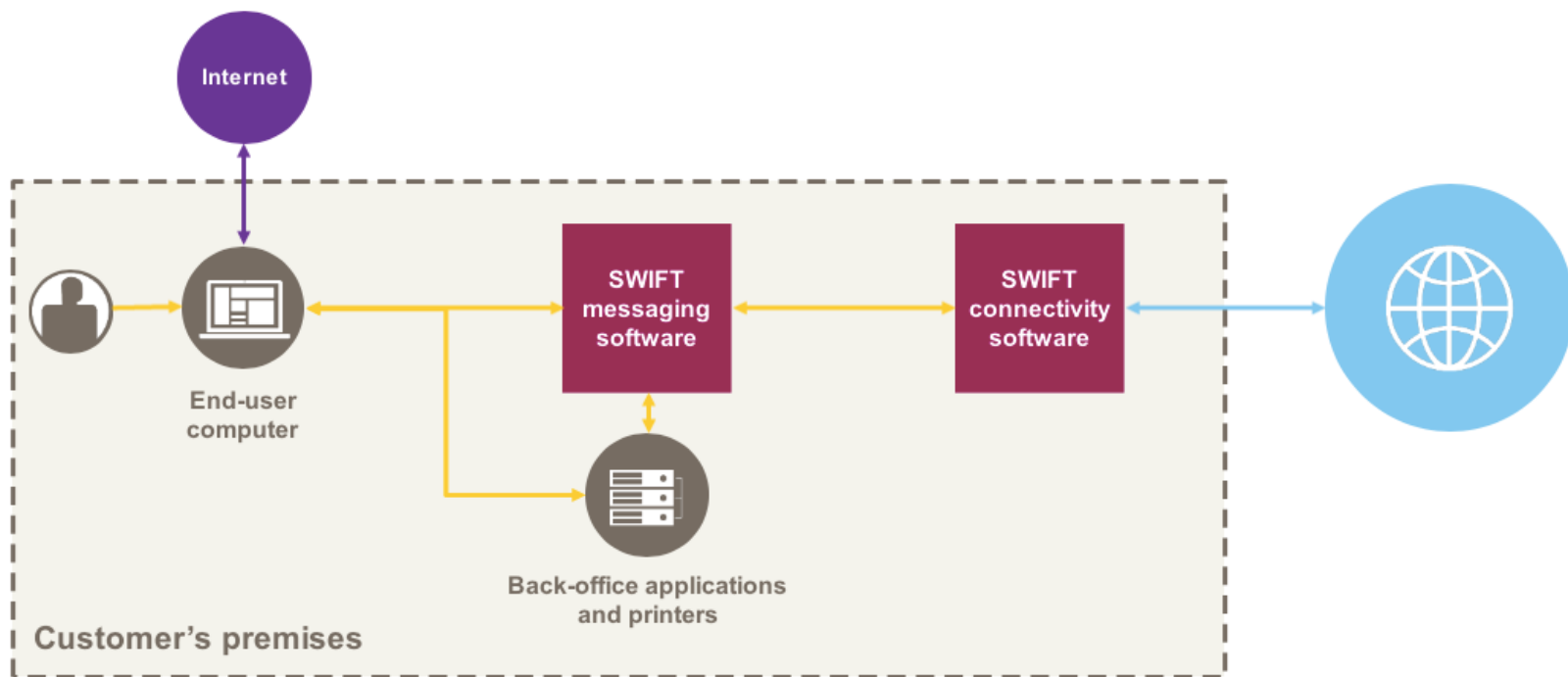
Spoofing

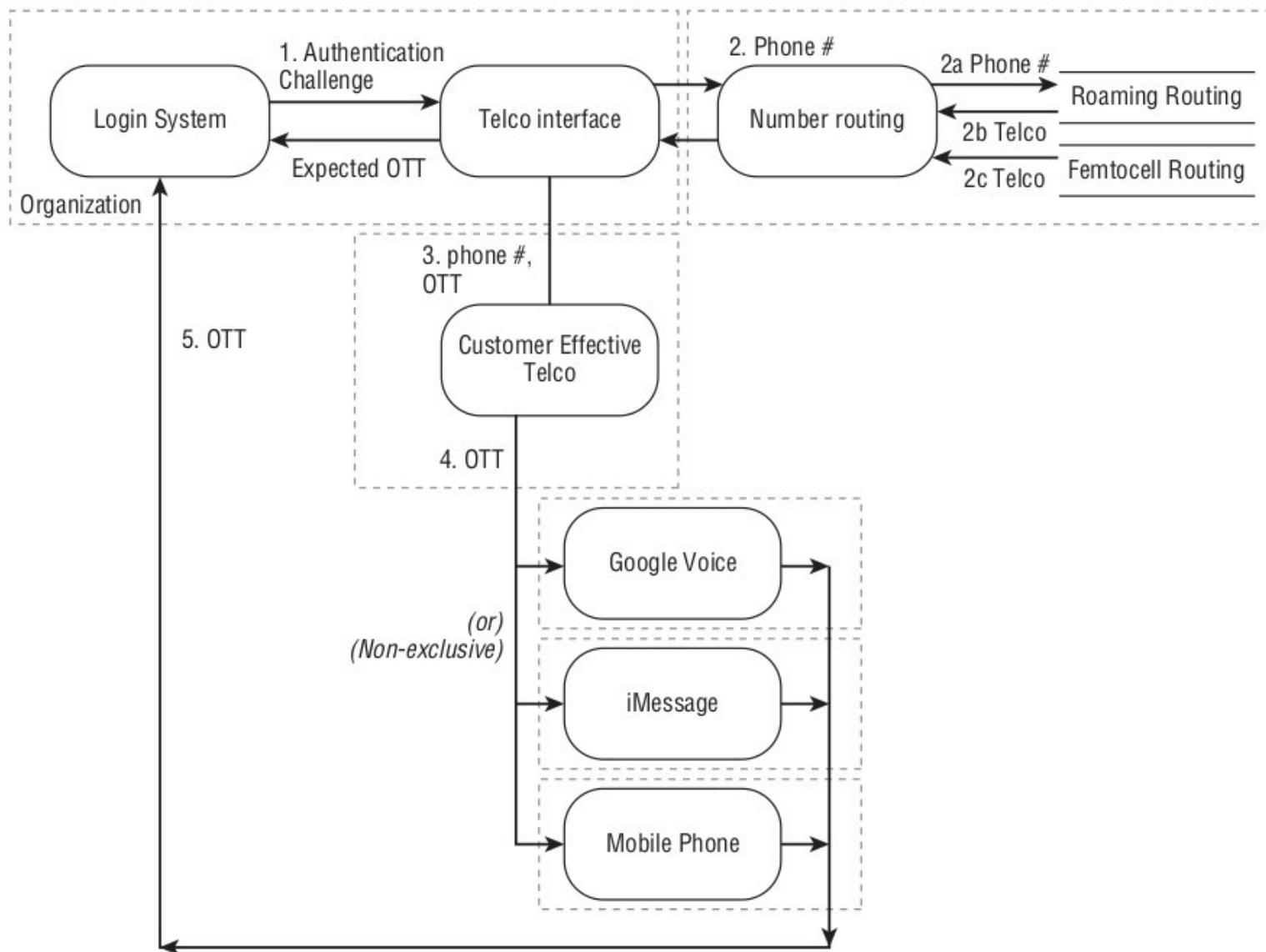
Elevation of privilege

Repudiation

“STRIDE”

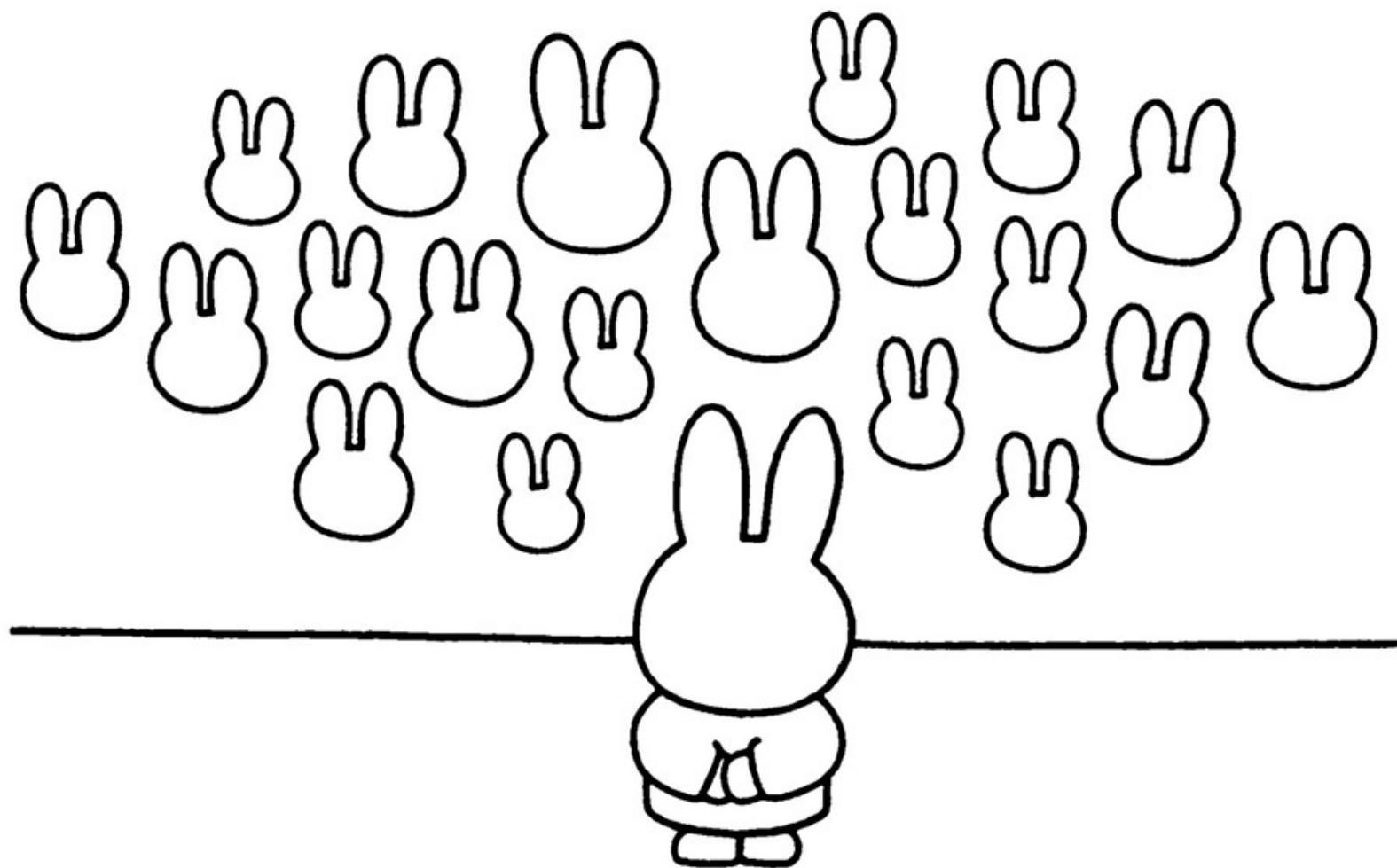






Lightweight methodology

1. Draw data flows
2. Elicit threats
- 3. Ranking + controls**
4. Check your work



RISK

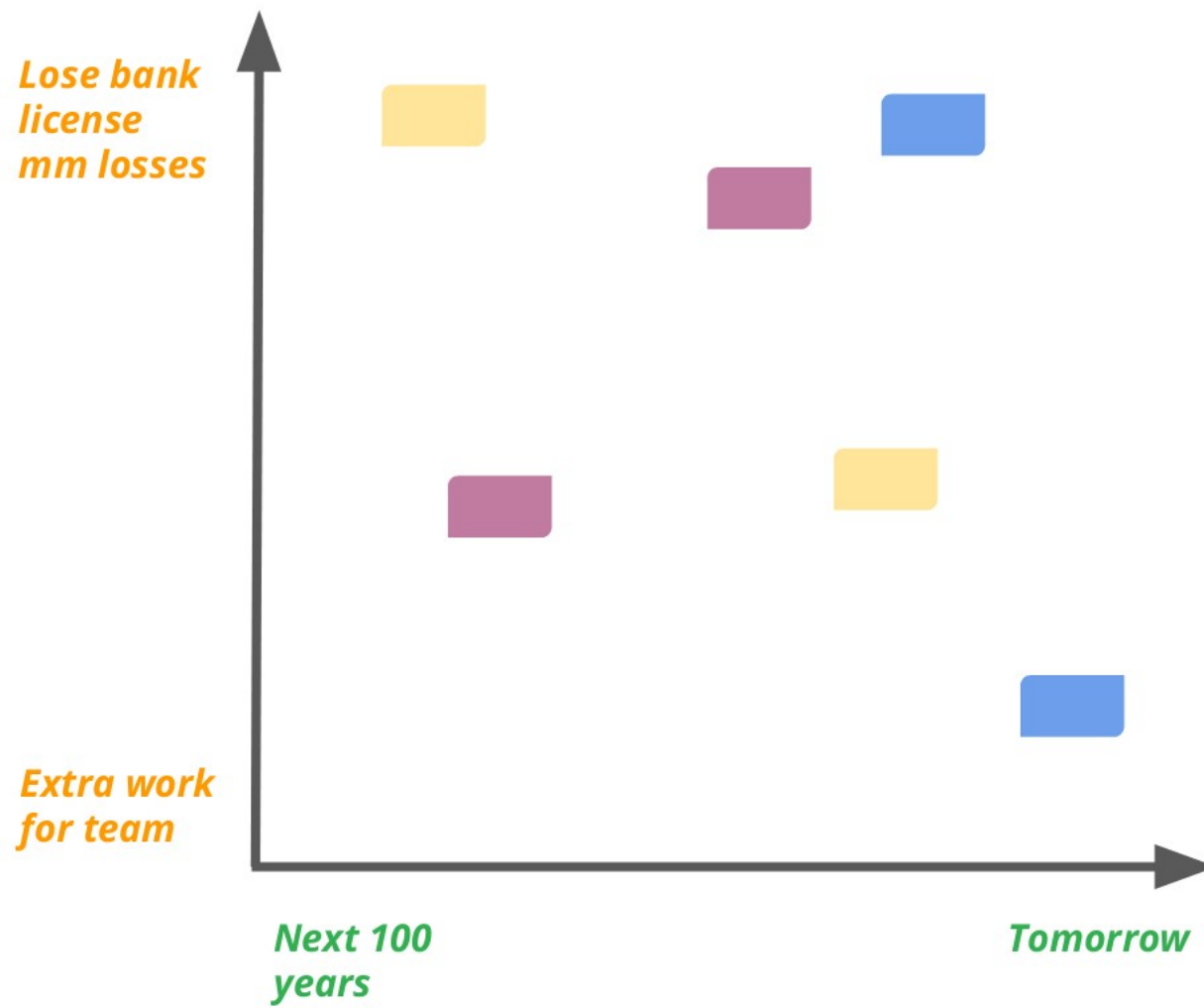
Parker Brothers' Trademark for its Continental Game Equipment.

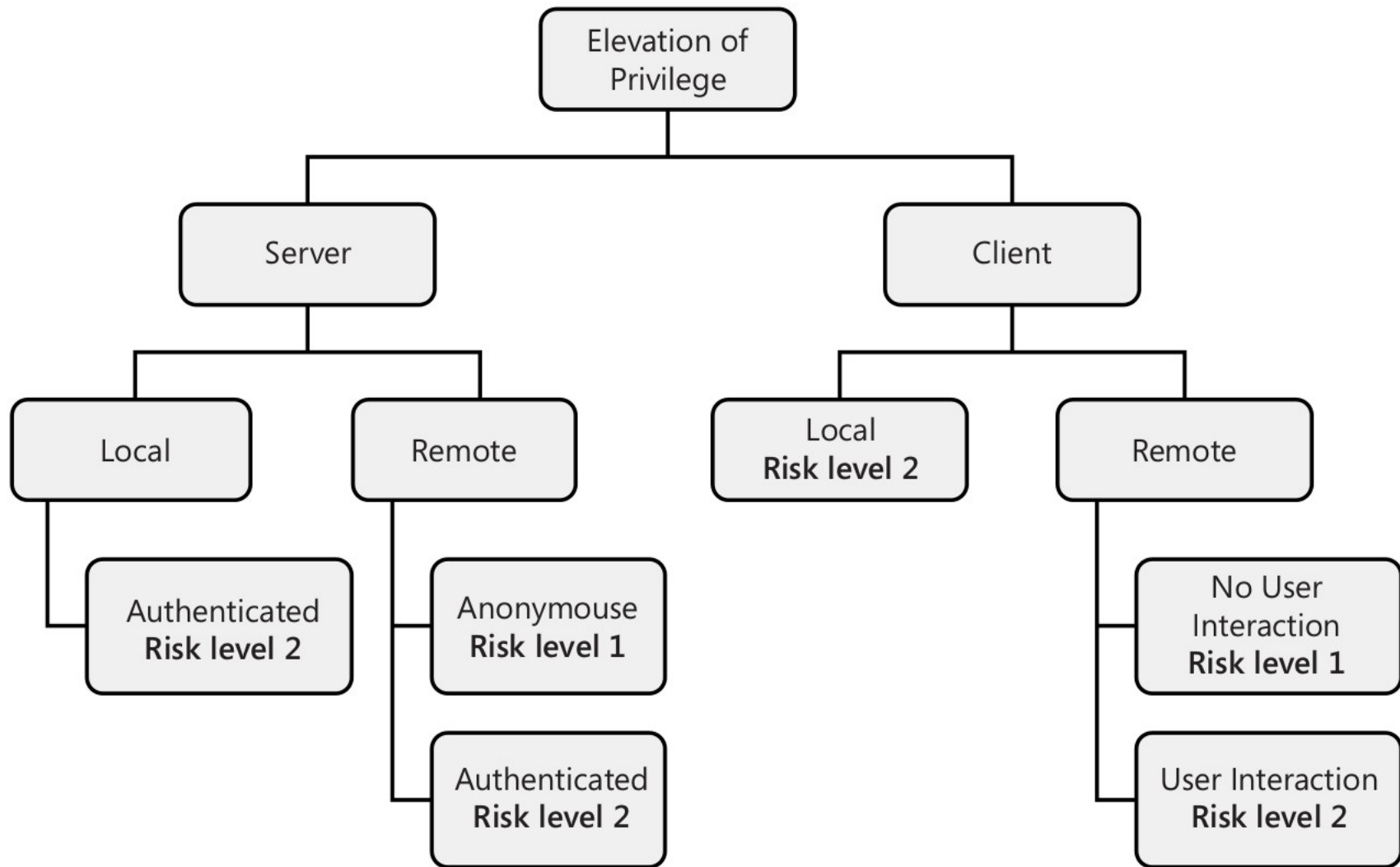
Ages 10 to adult For 2 to 6 players

**PARKER BROTHERS
CONTINENTAL GAME**

© 1988 Parker Brothers Inc. All Rights Reserved. Parker Brothers Inc., Franklin, Mass. 01890. Made in U.S.A.

$$\text{Risk} \approx \text{likelihood} \times \text{impact}$$



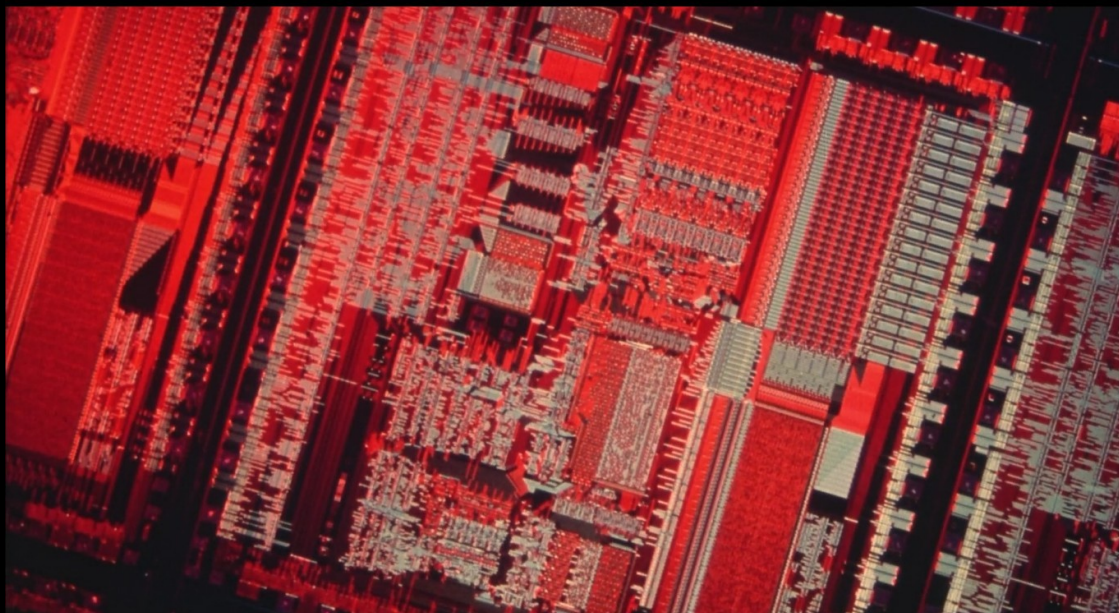


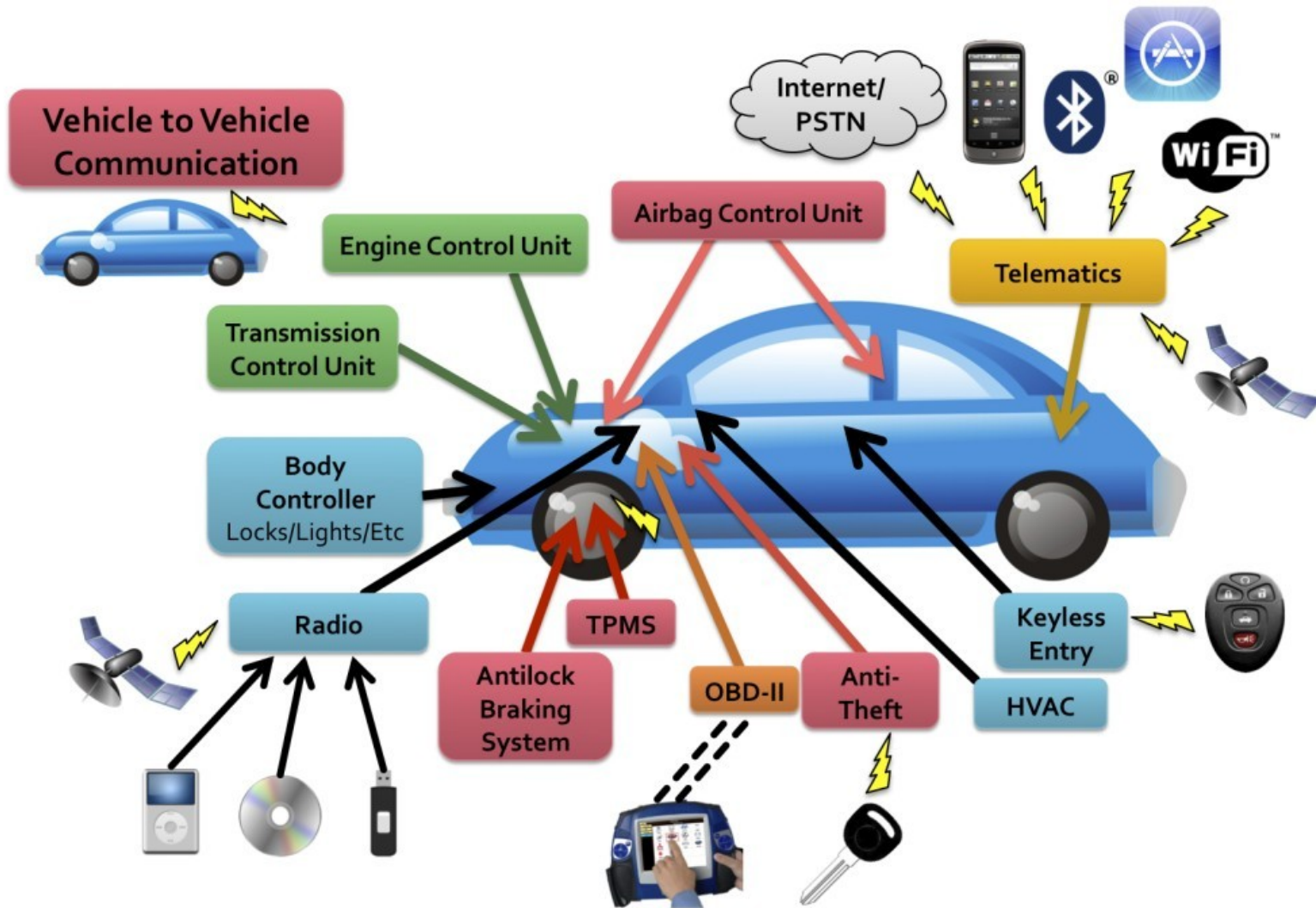
Lightweight methodology

1. Draw data flows
2. Elicit threats
3. Ranking + controls
4. **Check your work**

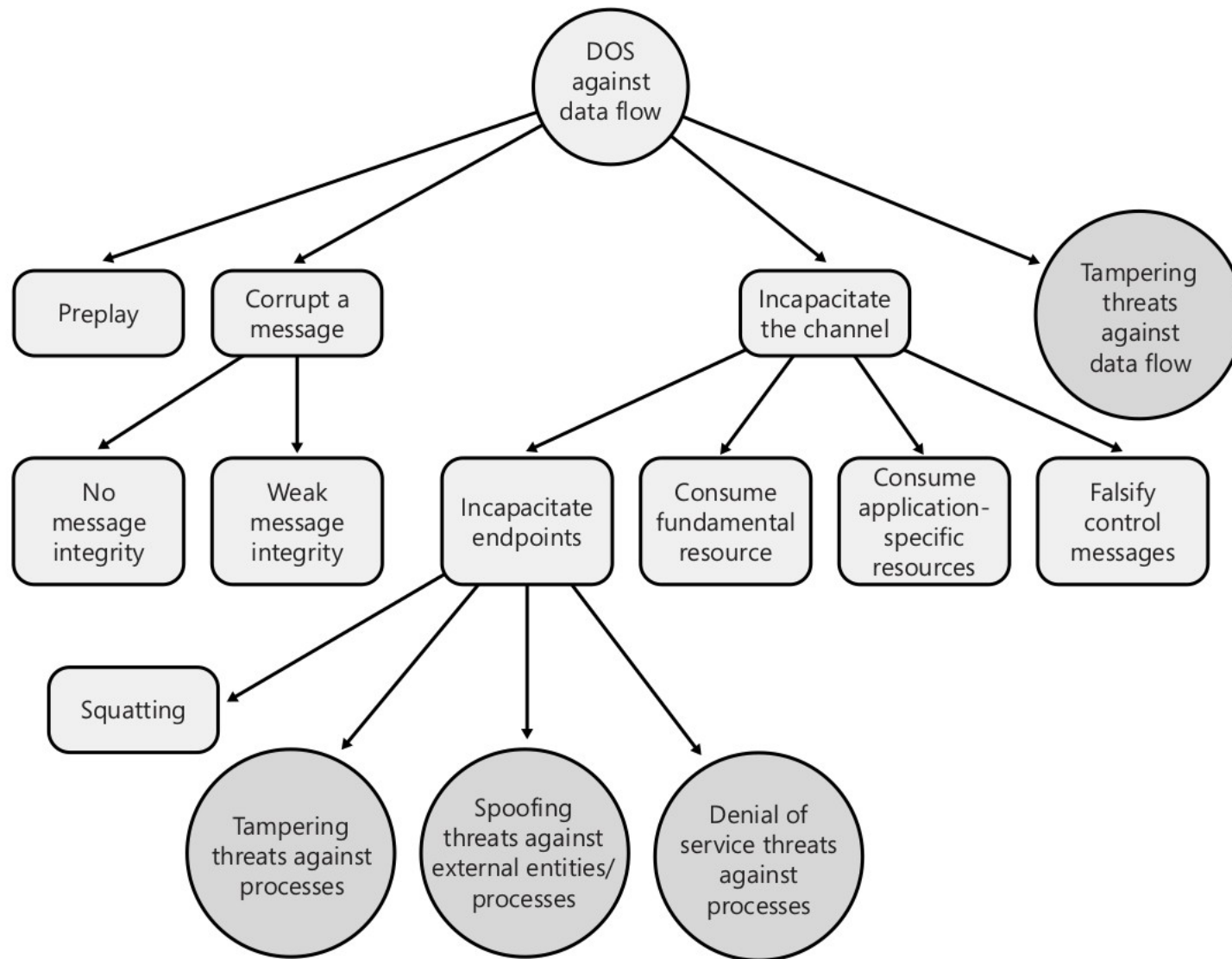
*“ All models are wrong,
some models are useful. ”*

— George Box





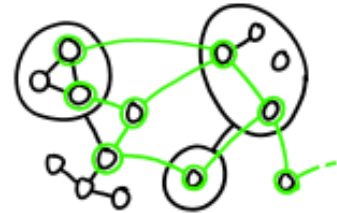
1. information disclosure
2. data interpreted as code
3. resource exhaustion/denial
4. race conditions
5. canonicalisation
6. insufficient access control
7. environment (mis) configuration
8. logic errors
9. predictability
10. poor usability



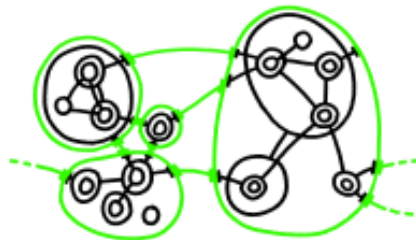
"I WISH THESE PARTS
COULD COMMUNICATE
MORE EASILY."



"OOH, THIS NEW TECHNOLOGY
MAKES IT EASY TO CREATE
ARBITRARY CONNECTIONS,
INTEGRATING EVERYTHING!"



"OOH, THIS NEW TECHNOLOGY
MAKES IT EASY TO ENCLOSE
ARBITRARY THINGS IN
SECURE SANDBOXES!"

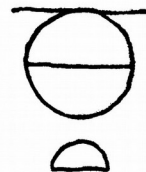
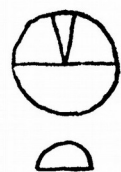
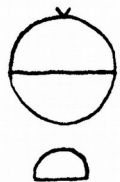
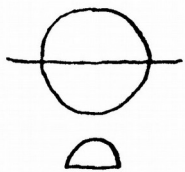
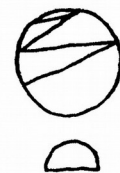
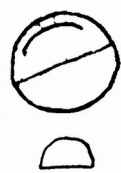
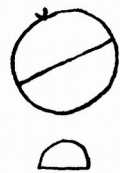
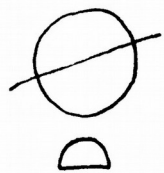


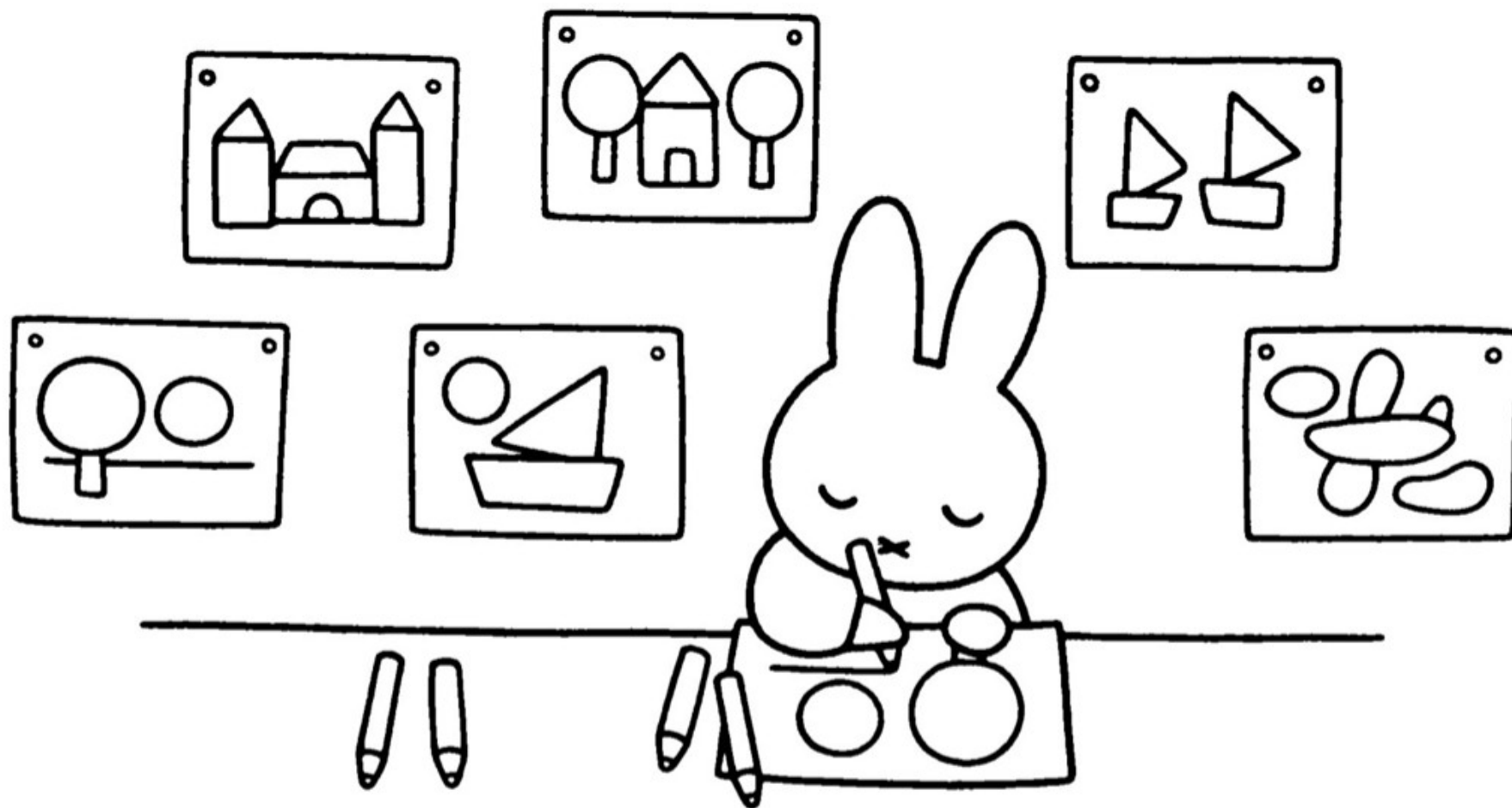
"UH-OH, THERE ARE
SO MANY CONNECTIONS
IT'S CREATING BUGS
AND SECURITY HOLES!"

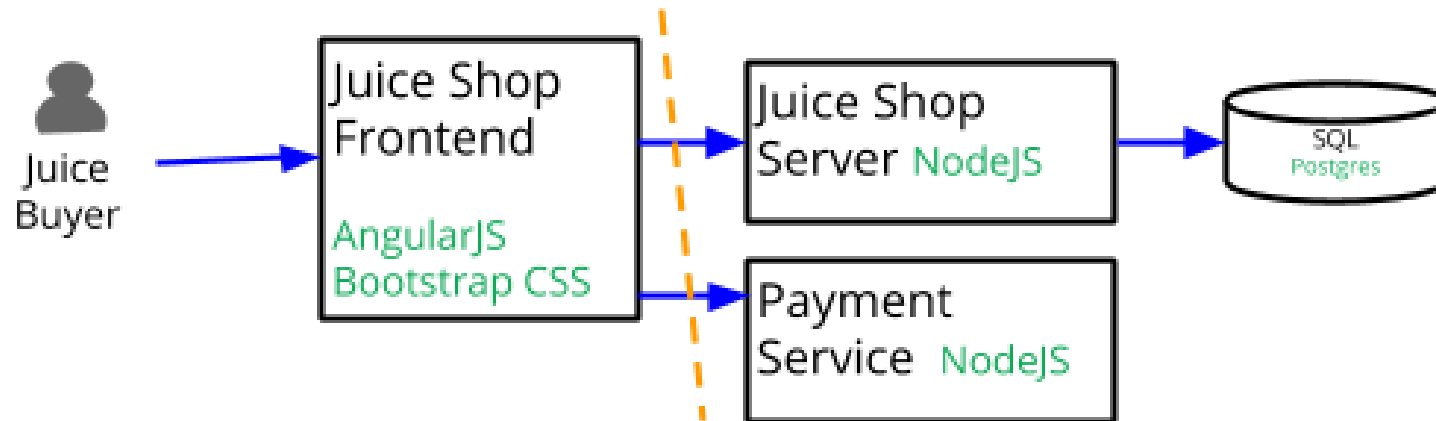


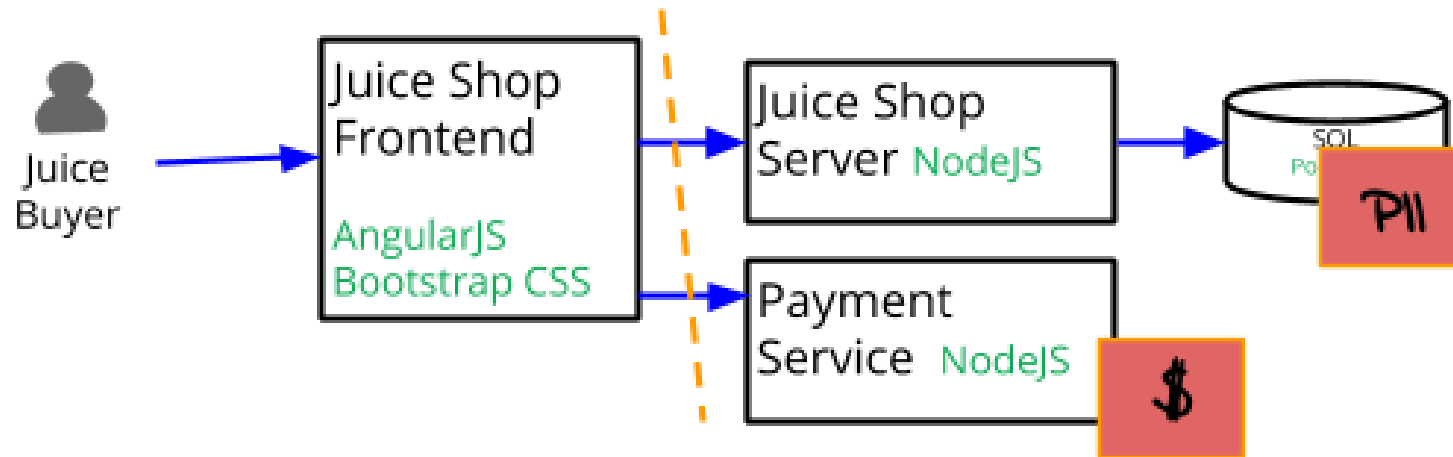
Lightweight methodology

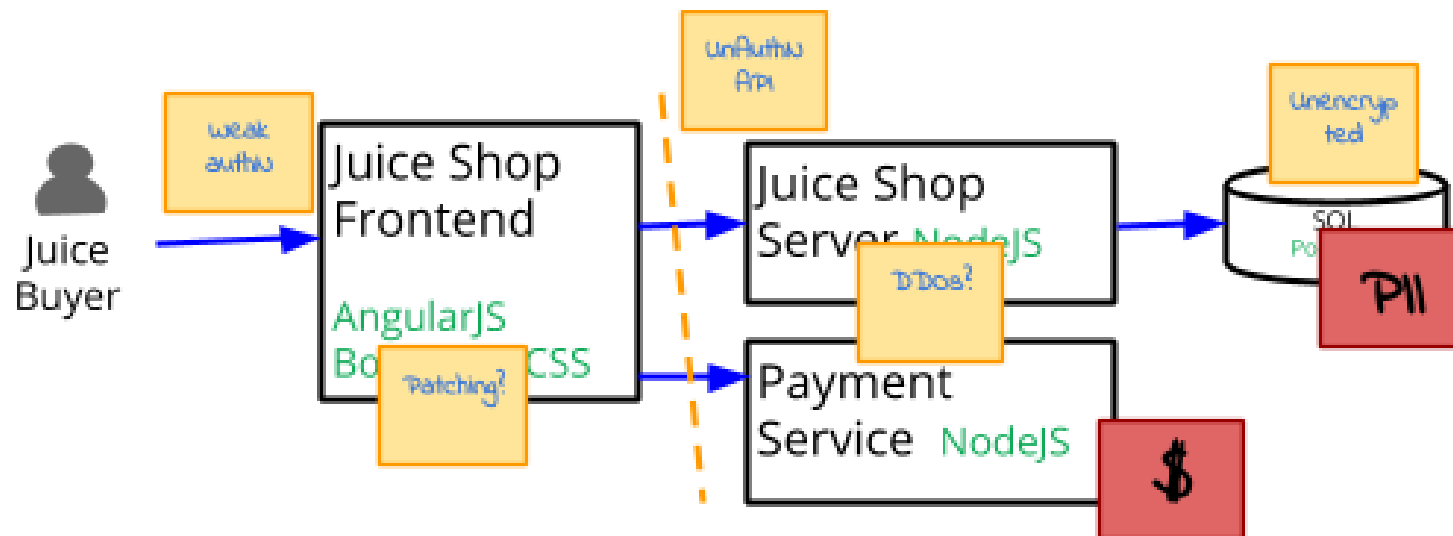
1. Draw data flows
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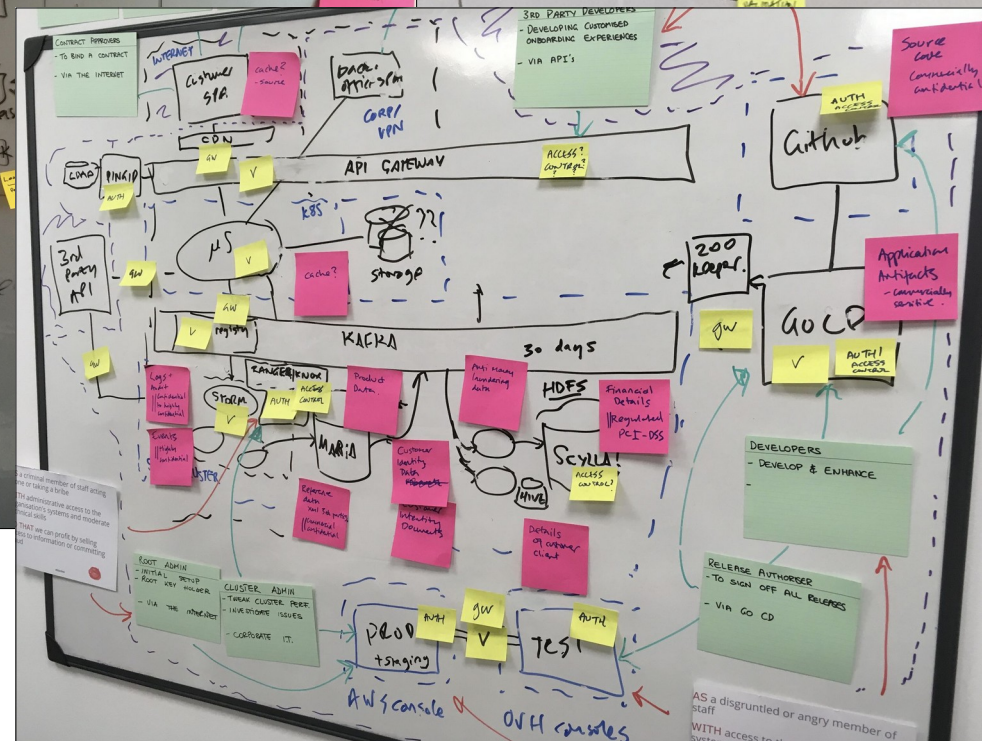














Microsoft

secure software
DEVELOPMENT SERIES

BEST PRACTICES

THE SECURITY DEVELOPMENT LIFECYCLE



*SDL: A Process for Developing Demonstrably
More Secure Software*

Michael Howard and Steve Lipner

Foreword by Jim Allchin

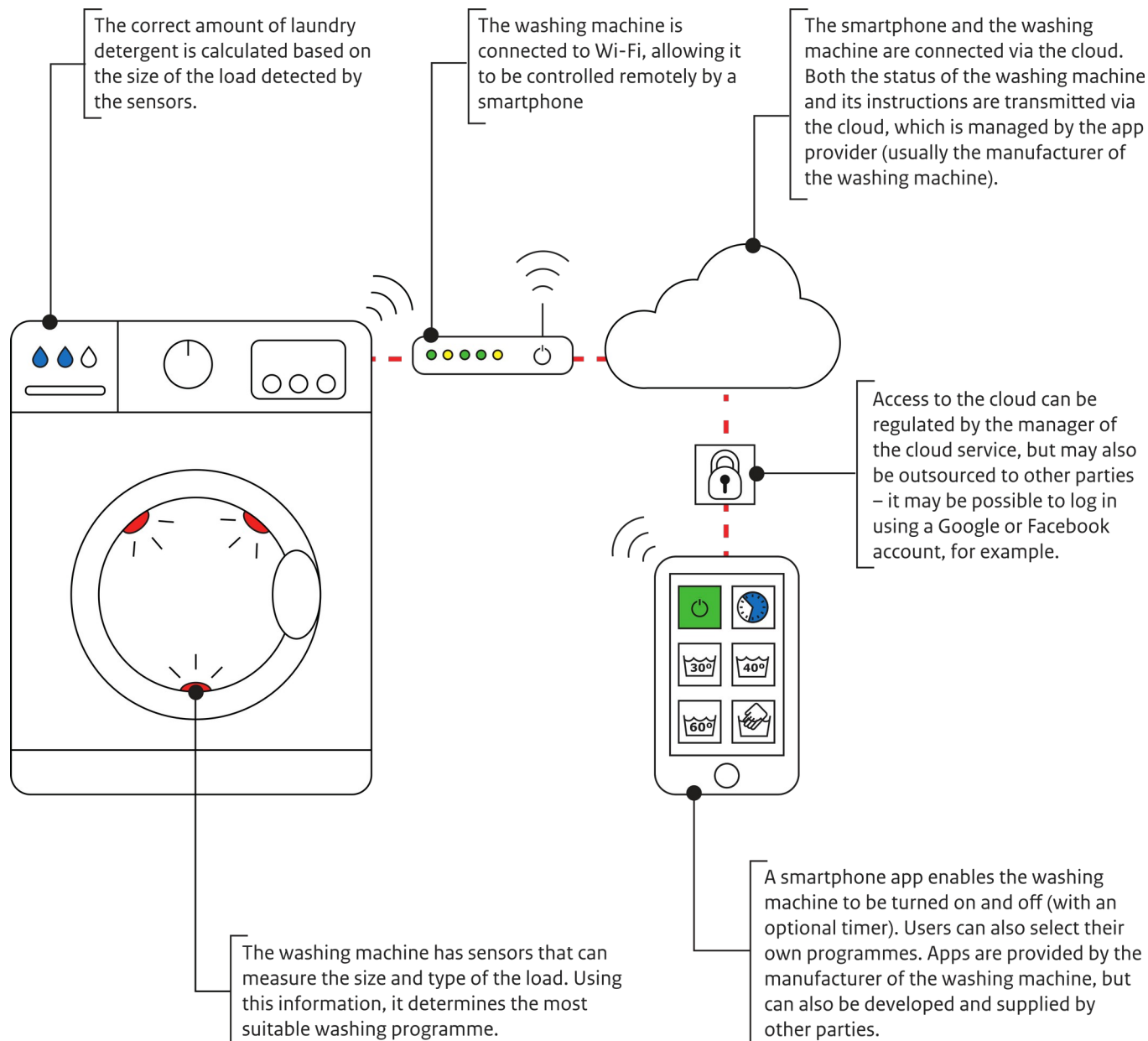
Co-President, Platforms & Services Division, Microsoft Corporation



CD Includes:

- A security training class video
 - Sample SDL documents
- See back cover





**What could
possibly
go wrong?**

Arne Padmos

hello@arnepadmos.com



**Hanzehogeschool
Groningen**

University of Applied Sciences



IT Academy
Noord-Nederland

github.com
/arnepadmos/resources

my “toy collection”